

A History of Approximate Nearest Neighbor Search from an Applications Perspective

Workshop of Approximate Nearest Neighbor Search: Bridging Theoretical
Foundations and Industrial Frontiers at FOCS 2025

Dec 14, 2025 @Sydney, Australia

Yusuke Matsui
The University of Tokyo



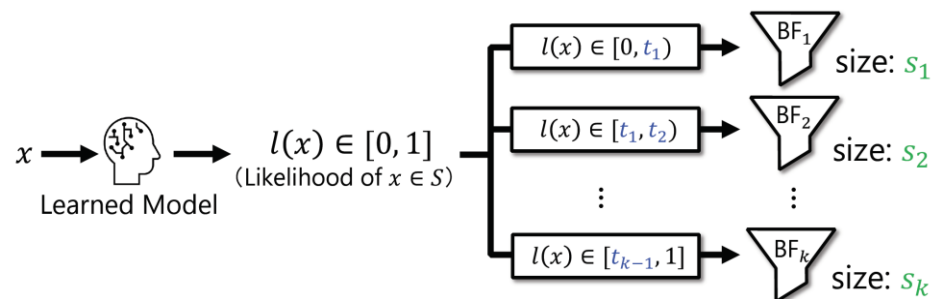


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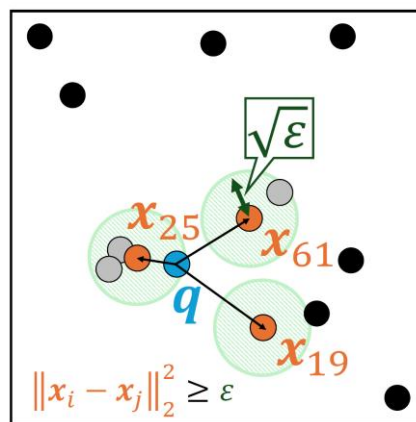
Lecturer (Assistant Professor), the University of Tokyo, Japan

🌐 <http://yusukematsui.me> 🐦 @utokyo_bunny 🐙 @matsui528

- ✓ Computer vision
- ✓ Data structure + Machine Learning

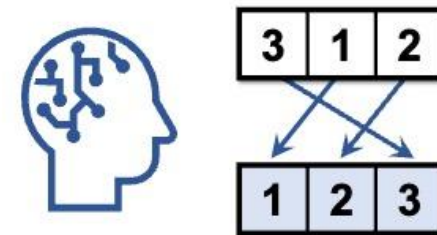


ML-enhanced Bloom Filter
[Sato & Matsui, NeurIPS 23]



Diverse nearest neighbor search
[Matsui, CVPR 25]

**$O(n \log \log n)$ Sort
with Machine Learning**



ML-enhanced Sorting
[Sato & Matsui, TMLR 25]

The 1st Workshop on Vector Databases (VecDB)

VecDB@[ICML2025](#), 18 July 2025, Vancouver, Canada

Speakers



[Matthijs Douze](#) (Meta)

Faiss



[Hiroyuki Ootomo](#) (NVIDIA)

CAGRA



[Prashant Pandey](#) (Northeastern University)

Organizers



[Yusuke Matsui](#) (The University of Tokyo)



[Martin Aumüller](#) (IT University of Copenhagen)



[Magdalen Dobson Manohar](#) (Carnegie Mellon University)



[Harsha Vardhan Simhadri](#) (Microsoft)

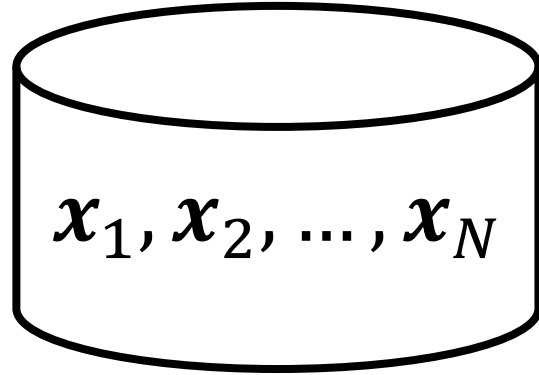
DiskANN

- Organized the 1st workshop on Vector Databases at ICML 2025
- **Planning the 2nd edition in 2026**
- Forum for NN researchers from various fields
- Welcome your submissions!



The best paper was a theory paper!
H. Xu, P. Indyk, S. Silwal, "Bi-metric Framework for Efficient Nearest Neighbor Search"

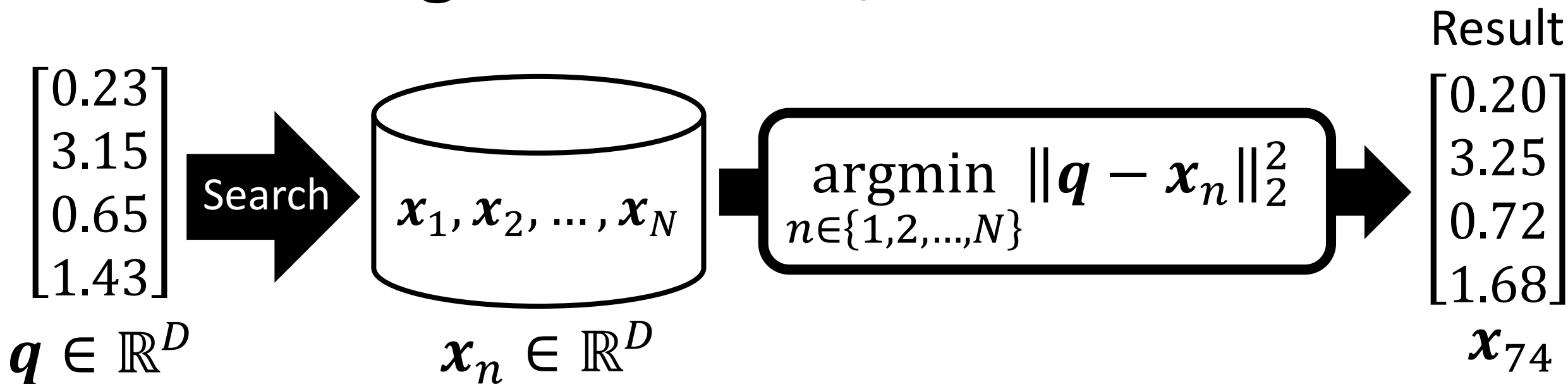
Nearest Neighbor Search; NN



$$\mathbf{x}_n \in \mathbb{R}^D$$

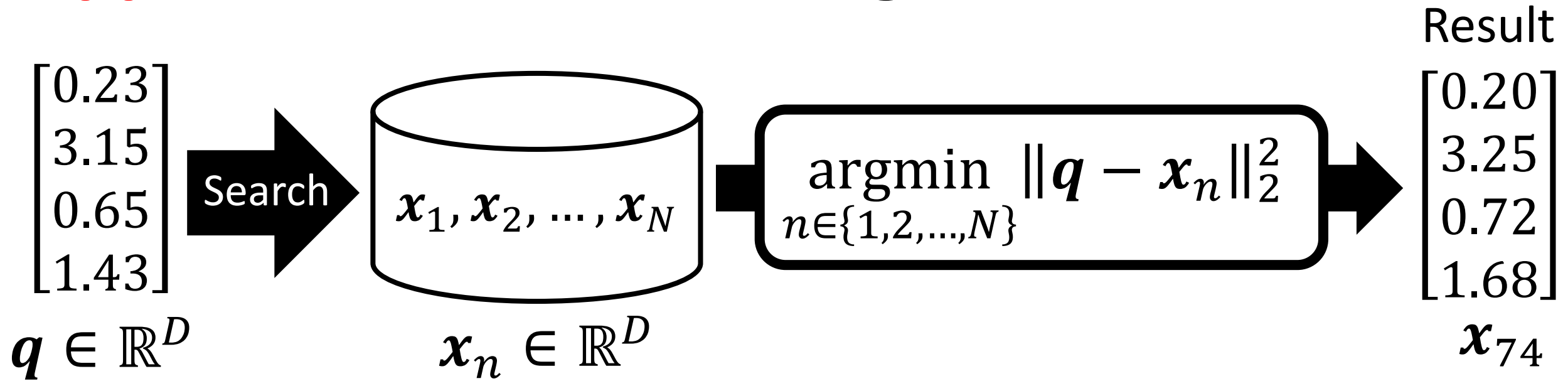
➤ N D -dim database vectors: $\{\mathbf{x}_n\}_{n=1}^N$

Nearest Neighbor Search; NN



- N D -dim database vectors: $\{x_n\}_{n=1}^N$
- Given a query q , find the closest vector from the database
- One of the fundamental problems in computer science
- Solution: linear scan, $O(ND)$, slow 😞

Approximate Nearest Neighbor Search; ANN



- Faster search
- Don't necessarily have to be exact neighbors
- Trade off: runtime, accuracy, and memory-consumption

Why NN/ANN?

- NN/ANN is an interesting research area because:
 - ✓ it's a **pure theory** (yes, this is FOCS WS!)
 - ✓ at the same time, it is **directly used in applications**, e.g., Vector DBs

- Several research areas (CV, NLP, DB, ...)

- Because of RAG and Vector DB, ANN has become popular more and more

Jie+, "Survey of Vector Database Management Systems", arXiv 2023

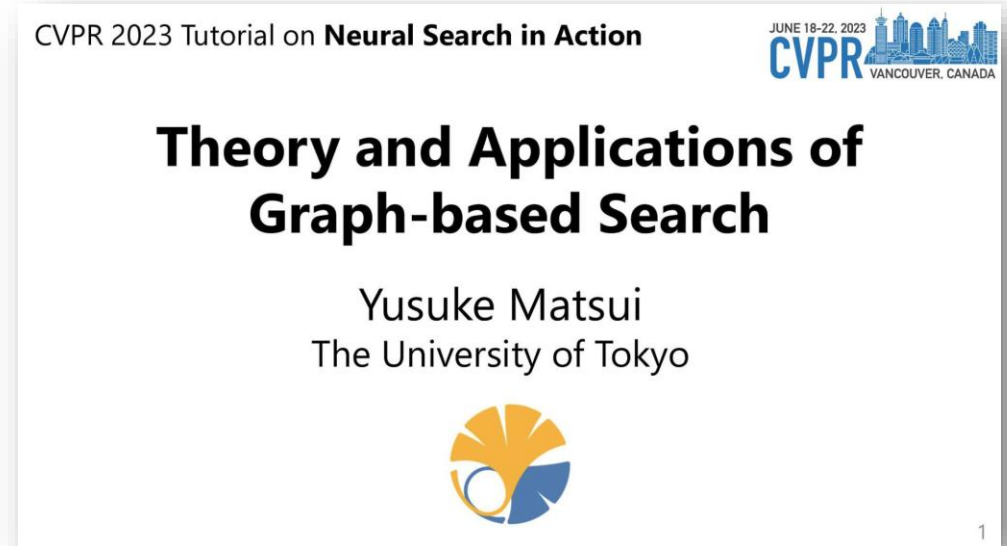
Name	Type	Sub-Type	Vector	
			Ex.	Ap
EuclidesDB (2018) [6]	Nat.	Vec.	✓	✓
Vearch (2018) [17,81]	Nat.	Vec.	✓	✓
Pinecone (2019) [9]	Nat.	Vec.	✓	✓
Vald (2020) [8]	Nat.	Vec.	✓	✓
Chroma (2022) [14]	Nat.	Vec.	✓	✓
Weaviate (2019) [5]	Nat.	Vec.	✓	✓
Milvus (2021) [12,125]	Nat.	Mix	✓	✓
NucliaDB (2021) [13]	Nat.	Mix	✓	✓
Qdrant (2021) [10]	Nat.	Mix	✓	✓
Manu (2022) [63]	Nat.	Mix	✓	✓
Marqo (2022) [11]	Nat.	Mix	✓	✓
Vespa (2020) [18]	Nat.	Mix	✓	✓
Cosmos DB (2023)	Ext.	NoSQL	✓	✓
MongoDB (2023)	Ext.	NoSQL	✓	✓
Neo4j (2023)	Ext.	NoSQL	✓	✓
Redis (2023)	Ext.	NoSQL	✓	✓
AnalyticDB-V (2020) [133]	Ext.	NoSQL	✓	✓
PASE+PG (2020) [139]	Ext.	Rel.	✓	✓
pgvector+PG (2021) [7]	Ext.	Rel.	✓	✓
SingleStoreDB (2022) [19,105]	Ext.	Rel.	✓	✓
ClickHouse (2023) [20]	Ext.	Rel.	✓	✓
MyScale (2023) [21]	Ext.	Rel.	✓	✓

This talk

- Reorganized and expanded version of my previous tutorials of CVPR2020 and CVPR2023
- See the original tutorials for more detailed contents



- CVPR 2020 Tutorial on Image Retrieval in the Wild
- Y. Matsui, "Billion-scale Approximate Nearest Neighbor Search"
- https://speakerdeck.com/matsui_528/cvpr20-tutorial-billion-scale-approximate-nearest-neighbor-search



- CVPR 2023 Tutorial on Neural Search in Action
- Y. Matsui, "Theory and Applications of Graph-based Search"
- https://speakerdeck.com/matsui_528/cvpr23-tutorial-theory-and-applications-of-graph-based-search

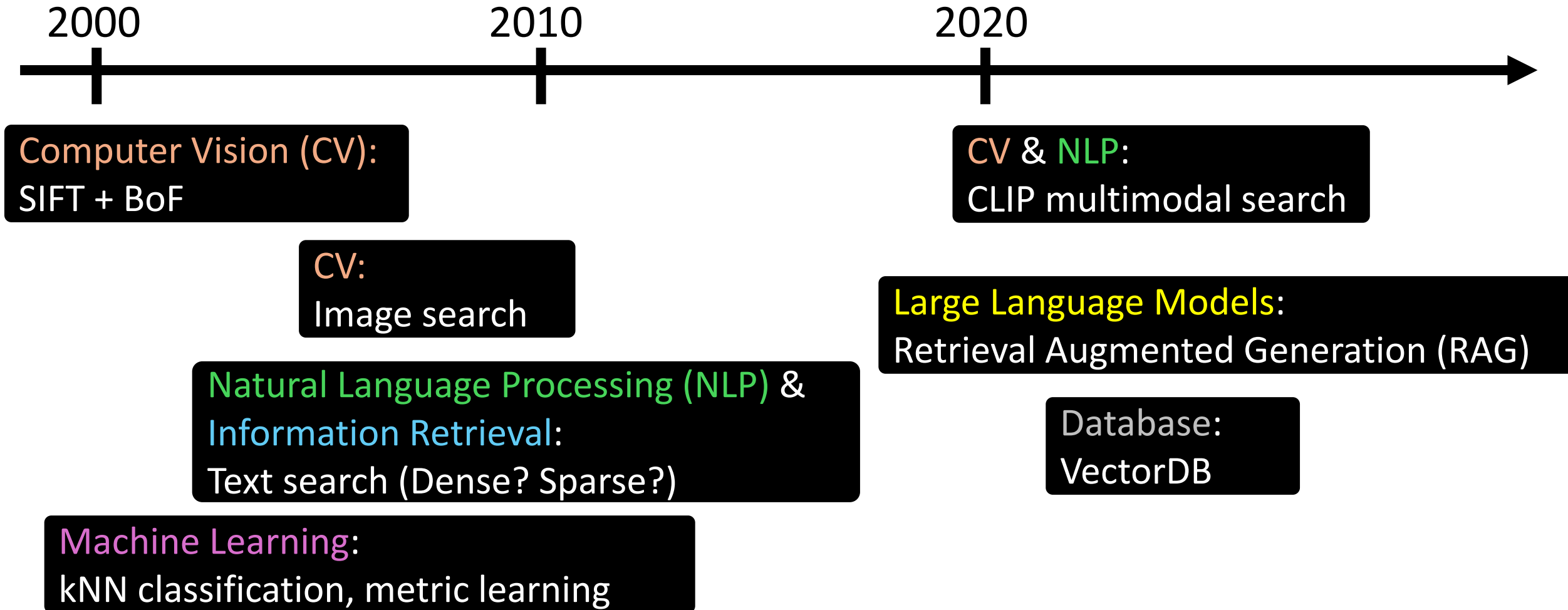
Outline

1. History from an applications perspective
2. Importance of implementation:
nearest neighbor search in faiss
3. Basics of modern baseline: graph-based search

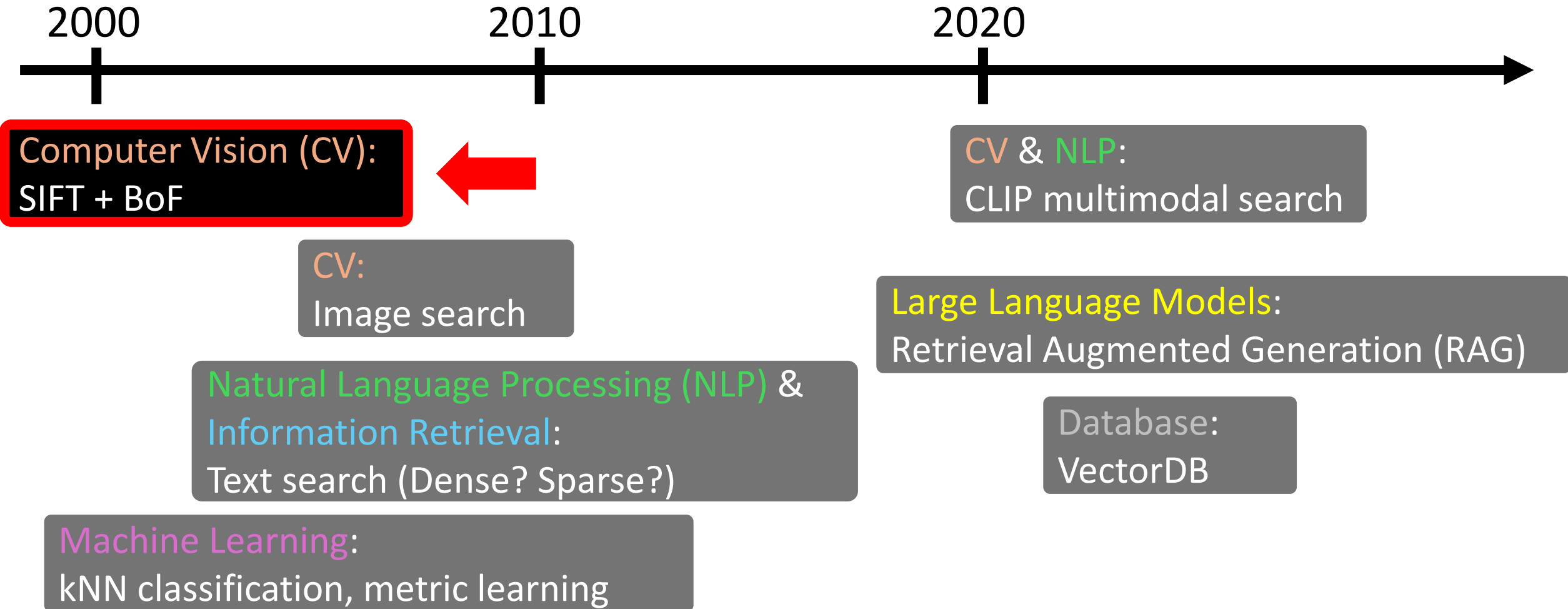
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What tasks has ANN been used for?

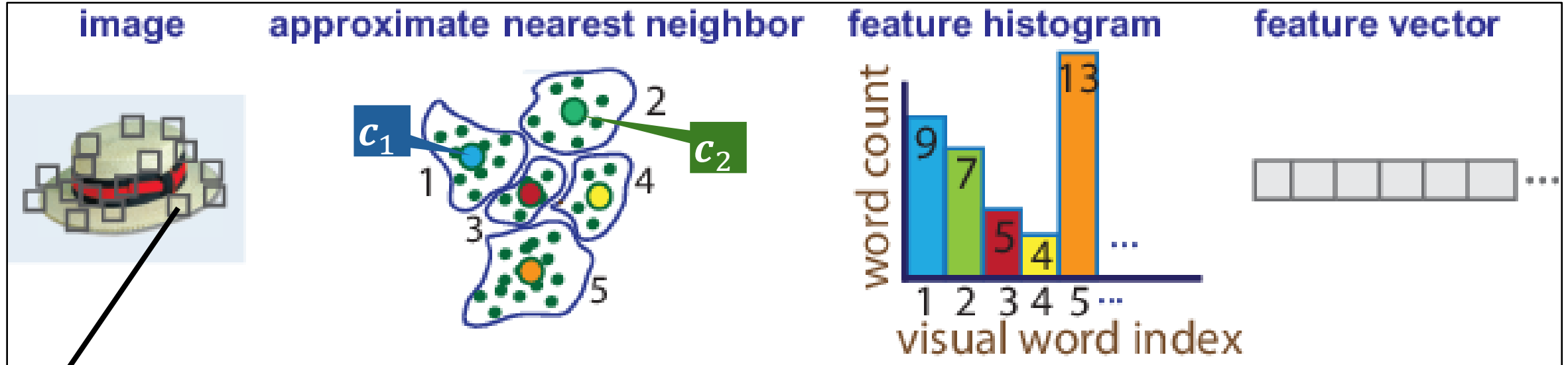


What tasks has ANN been used for?



SIFT (local feature) + BoF (Bag-of-features) + SVM

<https://jp.mathworks.com/help/vision/ug/image-classification-with-bag-of-visual-words.html>



$$\mathbf{x} \in \mathbb{R}^{128}$$

Extract a local patch

Given codewords $\{c_k\}$,
find the closest one

$$k^* = \underset{k \in \{1, \dots, 5\}}{\operatorname{argmin}} \|\mathbf{x} - c_k\|_2^2$$

Create a histogram, run SVM,
recognize an image...

- This is nearest neighbor search!
- K is 10^3 to 10^4
- Must be in memory

- To compute BoF fast, several practical ANN technologies have been invented in the CV area in the 2000s – 2010s... E.g.: Product Quantization
- Good old days.... 🤖

What tasks has ANN been used for?

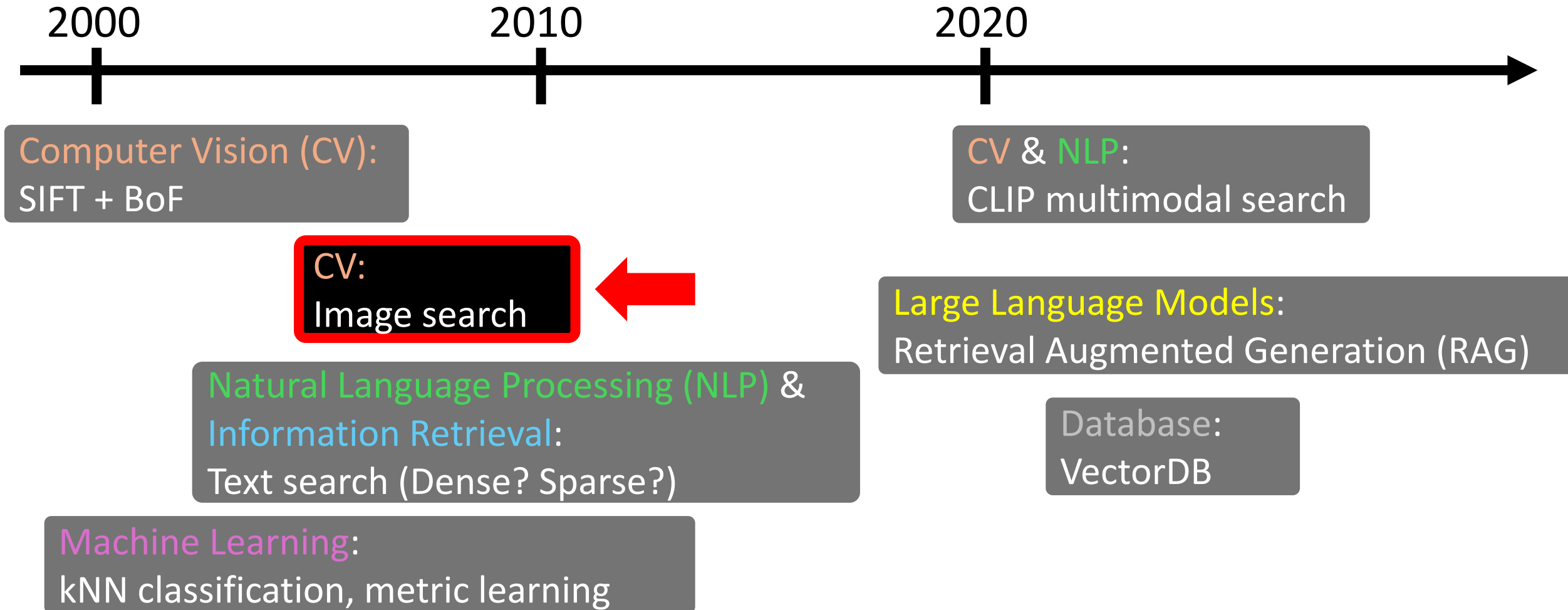


Image Search

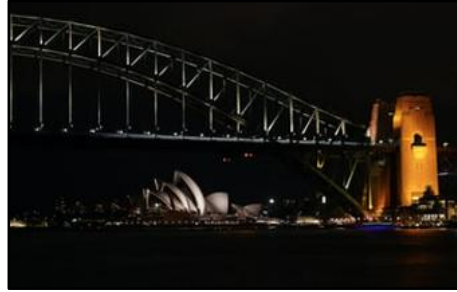


Image are from: <https://github.com/haltakov/natural-language-image-search>
Credit: Photos by [Genton Damian](#), [bruce mars](#), [Dalal Nizam](#), and [Richard Burlton](#) on [Unsplash](#)

Image Search

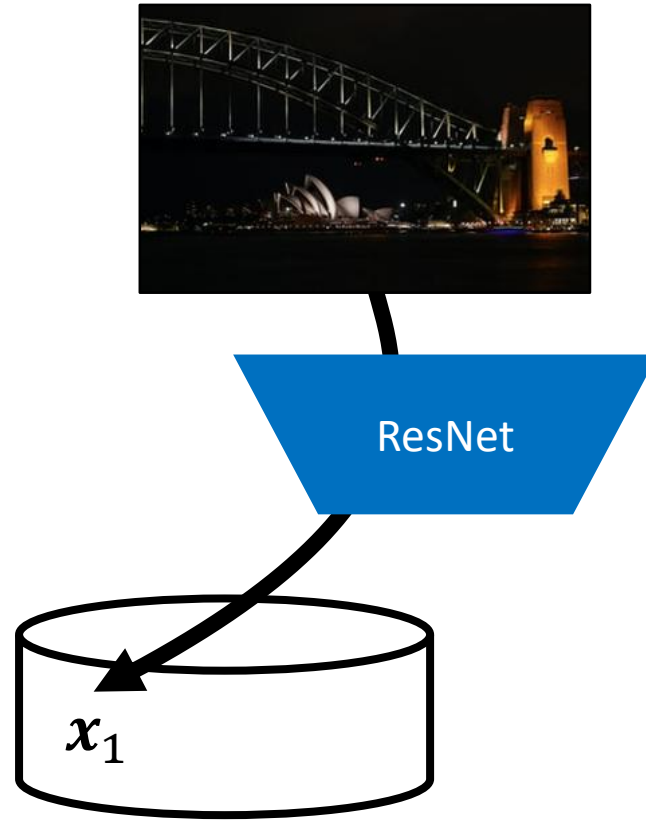


Image Search

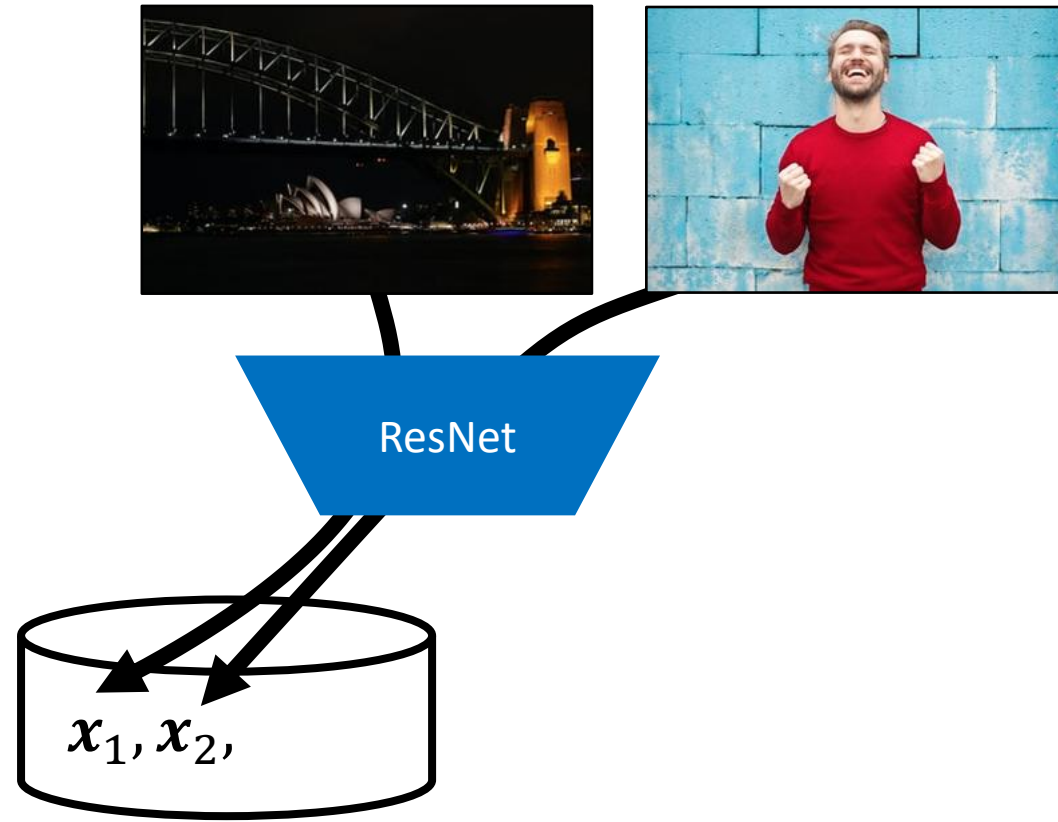


Image Search

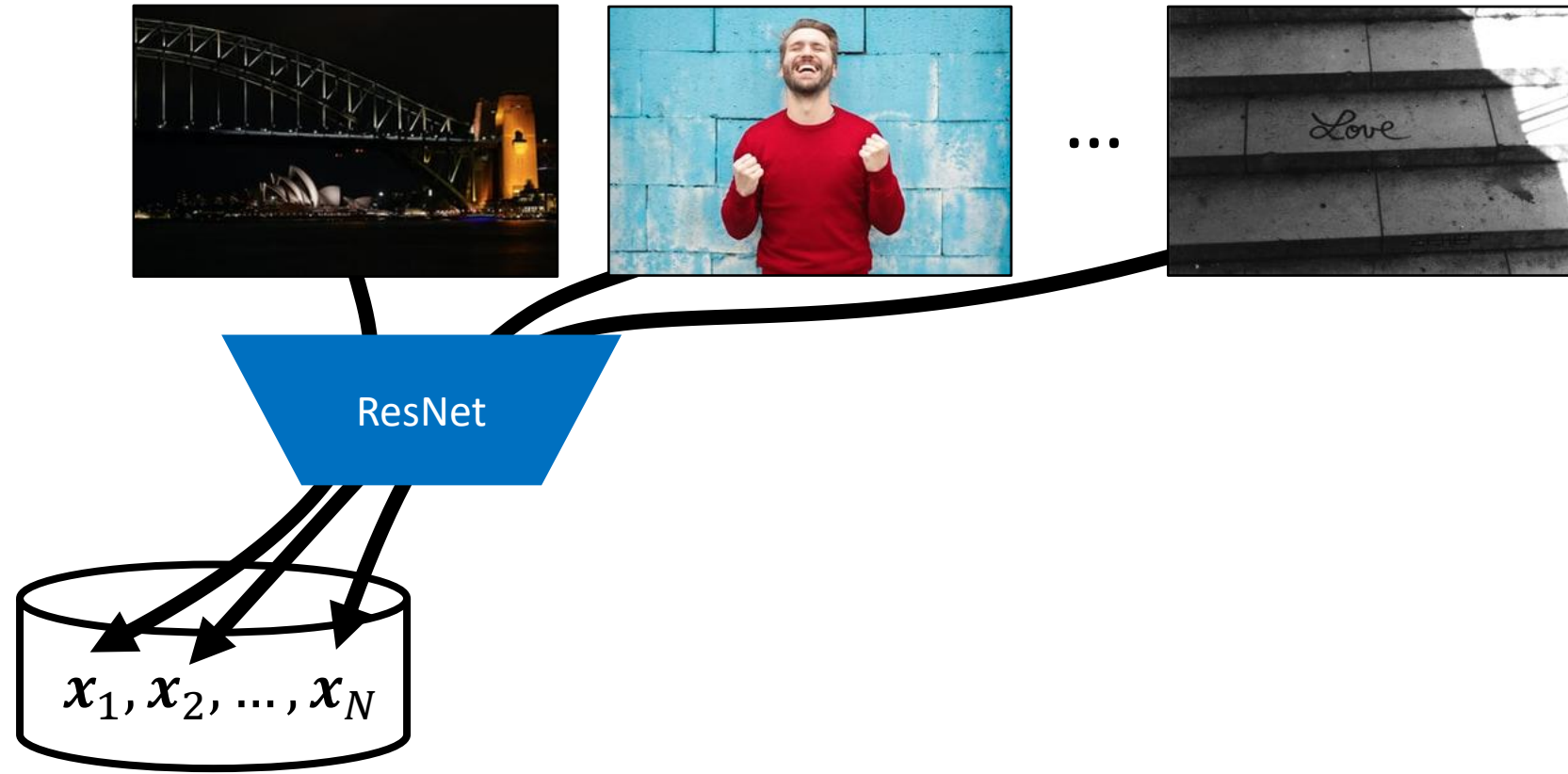


Image Search

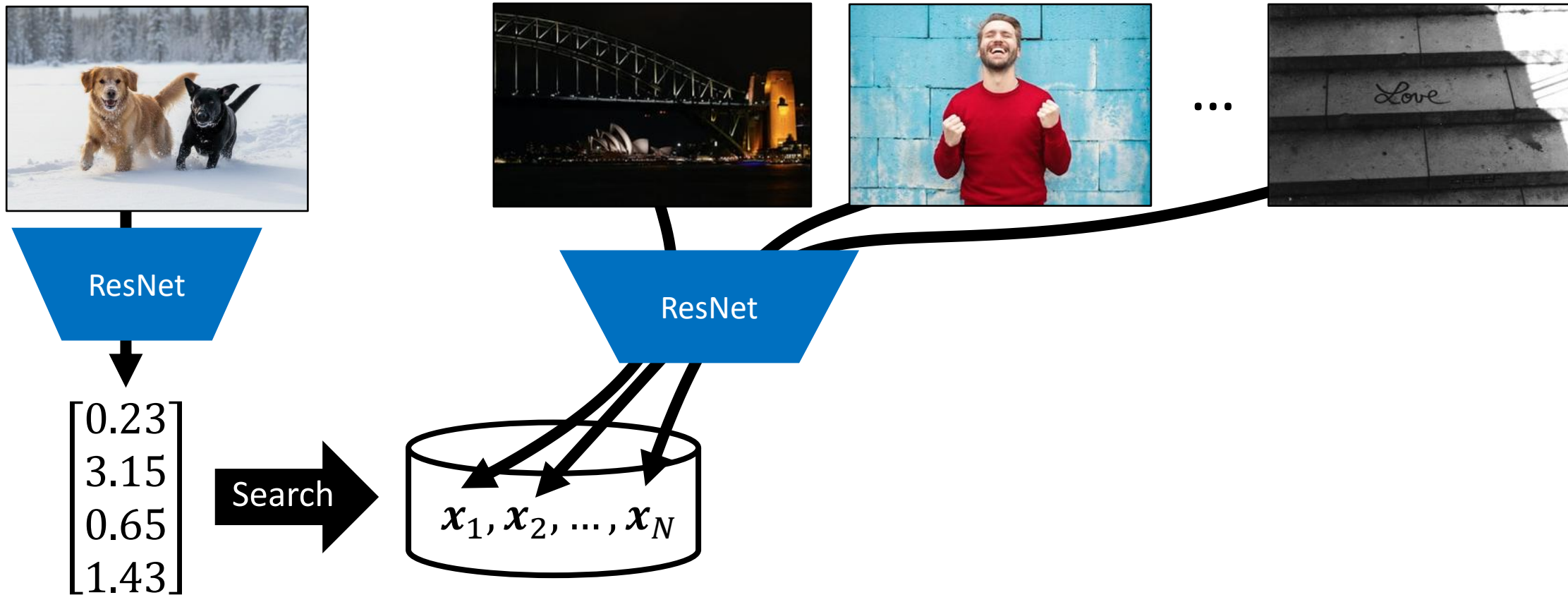


Image Search

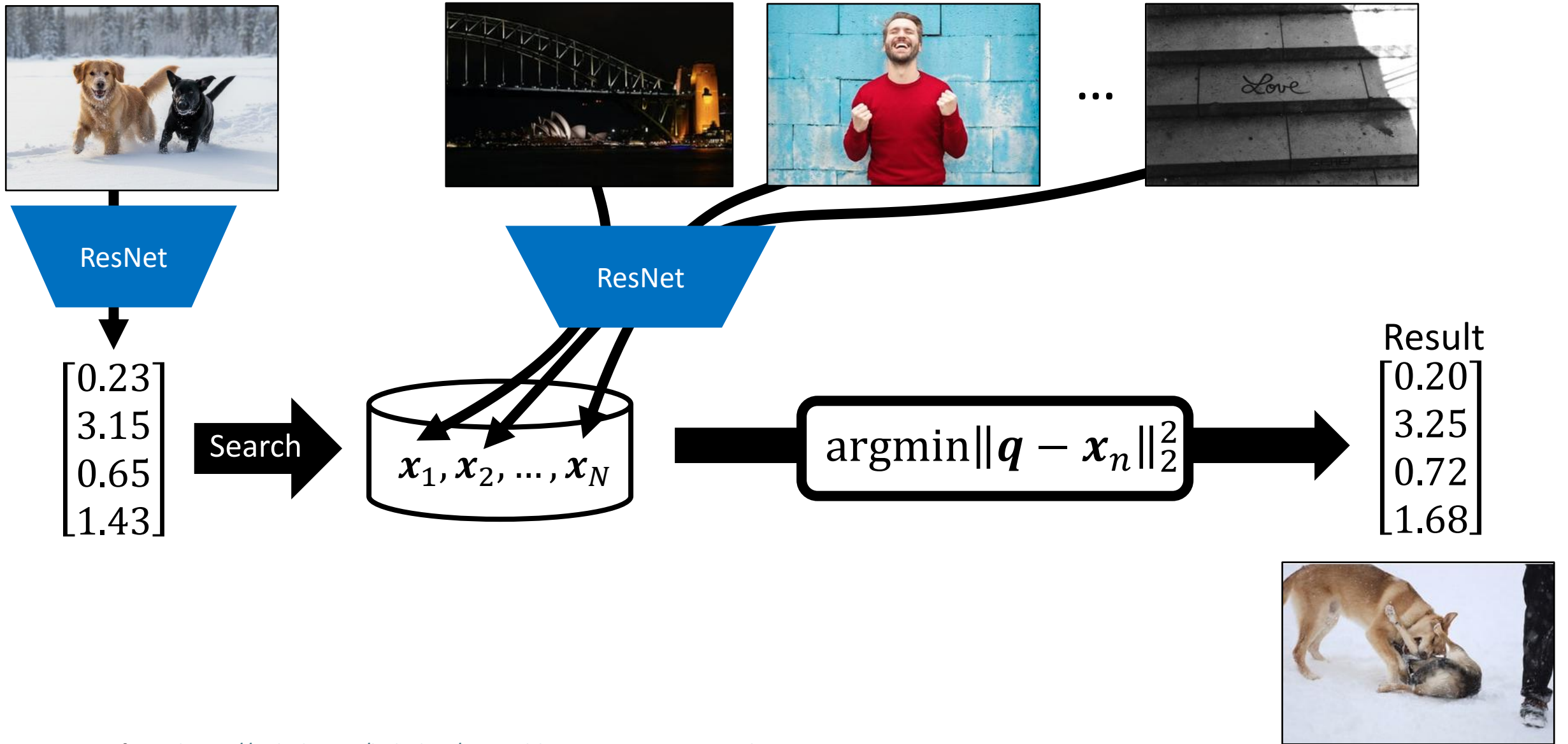
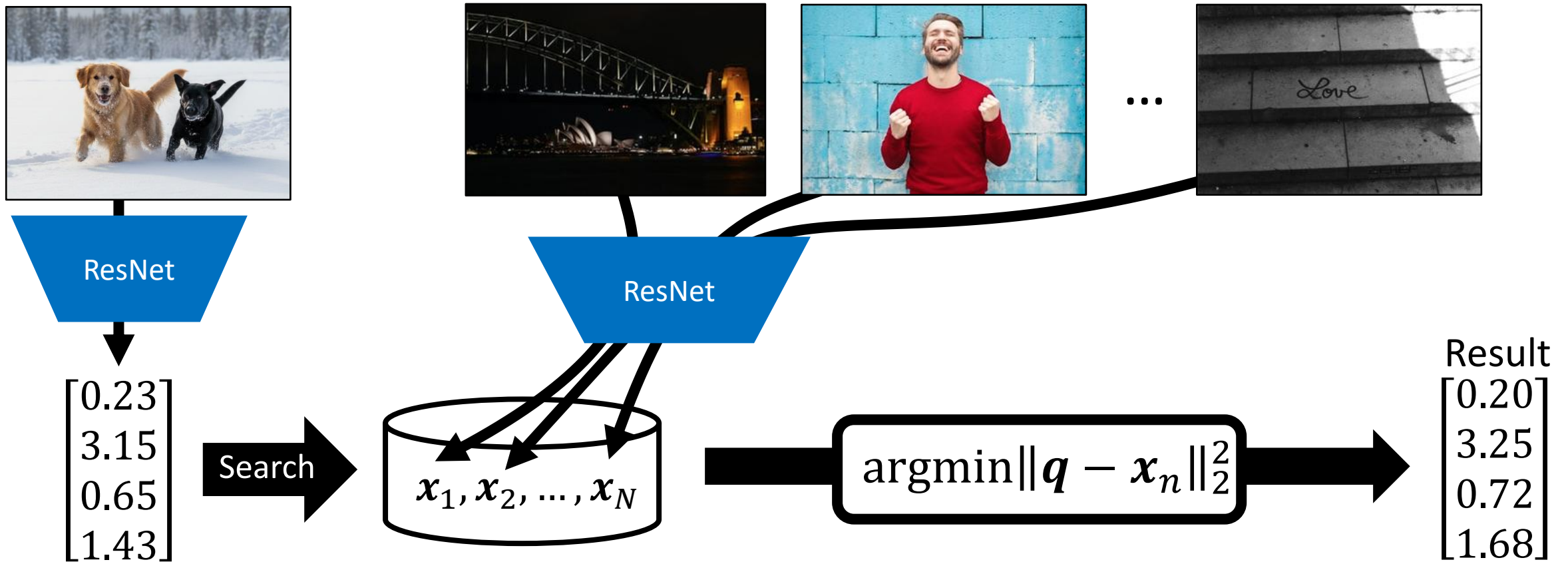


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Image Search



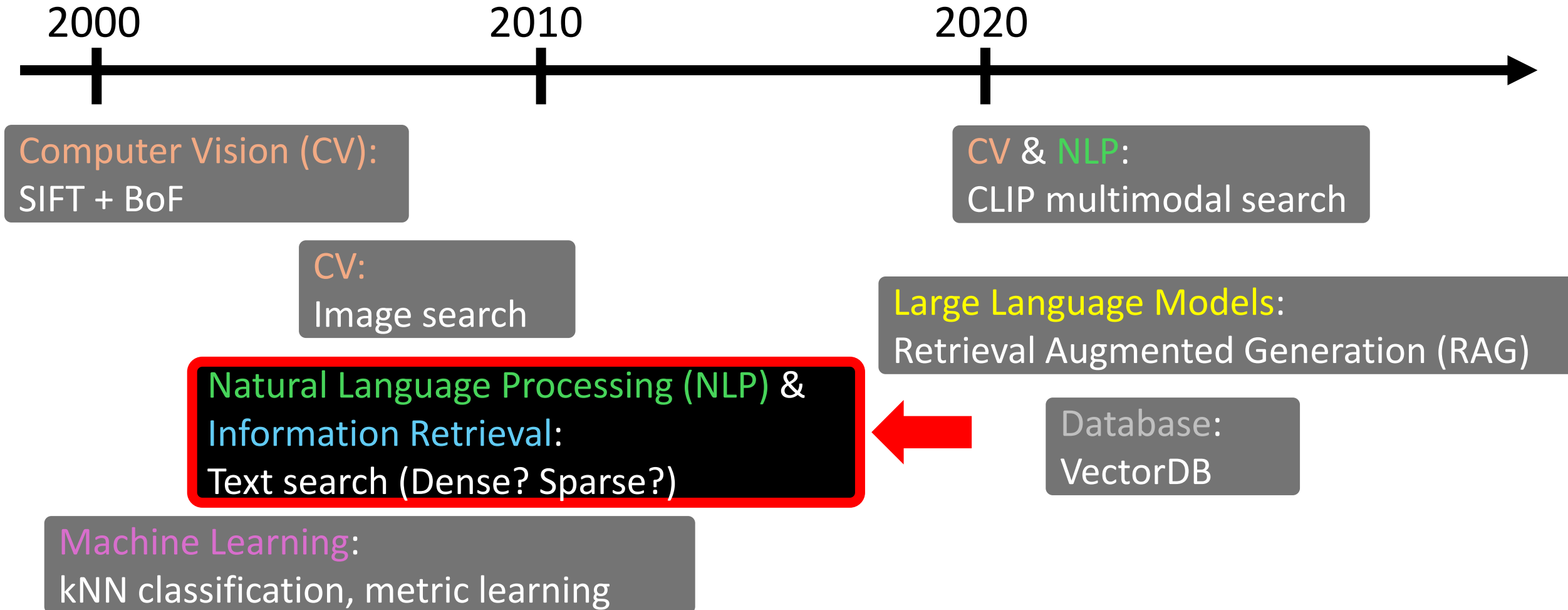
- Use a pre-trained CNN model (e.g., ResNet) as a feature extractor
- Represent an image as a high-dimensional vector
- Image-retrieval by nearest neighbor search

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What tasks has ANN been used for?



Text Search (Dense and/or Sparse)

Query: “fox”

search

Doc1: “The cat sleeps on chairs”

Doc2: “A quick blue fox runs”

Doc3: “The cat sits near window”

...

Dense search

q_{fox}

Search

$x_1, x_2, \dots, x_N$

$\operatorname{argmin} \|q_{\text{fox}} - x_n\|_2^2$

x_2

➤ How to design embedding? BERT and its successors...

Sparse search

“fox”

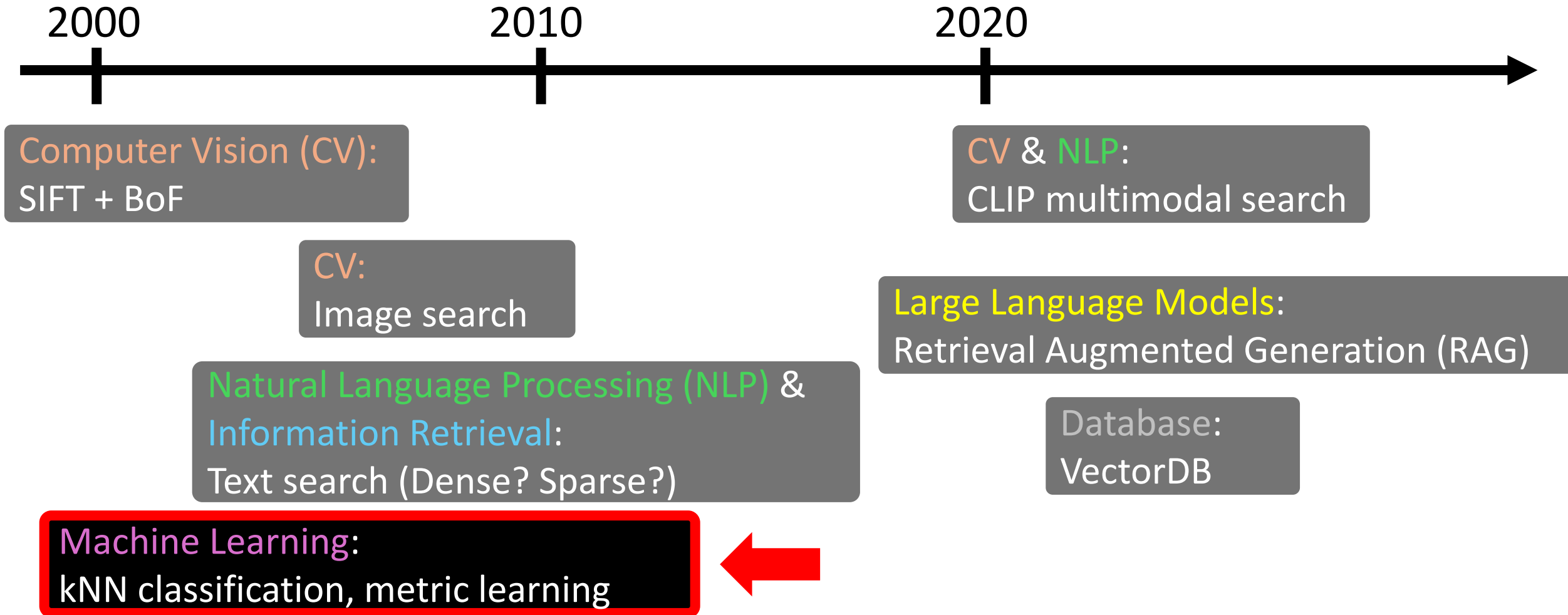
Term	ID
the	1, 3
fox	2
cat	2, 3
Blue	2
...	...

Doc2

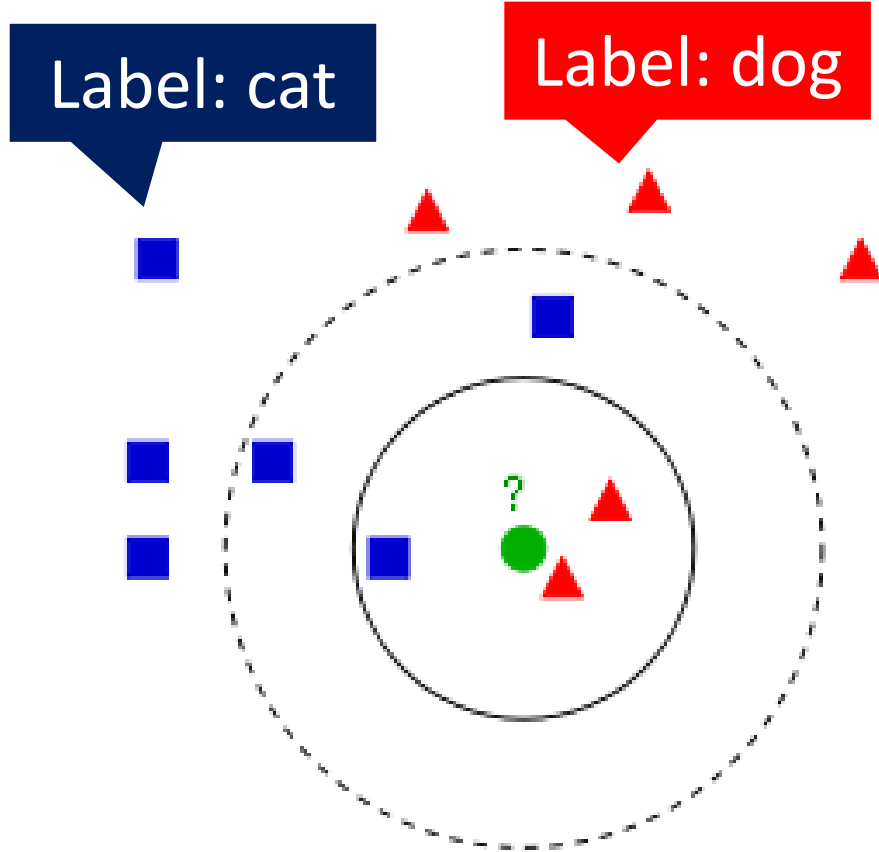
- Classical “matching” and its extensions
- TF-IDF, BM25, SPLADE...

- Dense search: accurate
- Sparse search: fast
- Combining two approaches is a hot topic in information retrieval

What tasks has ANN been used for?



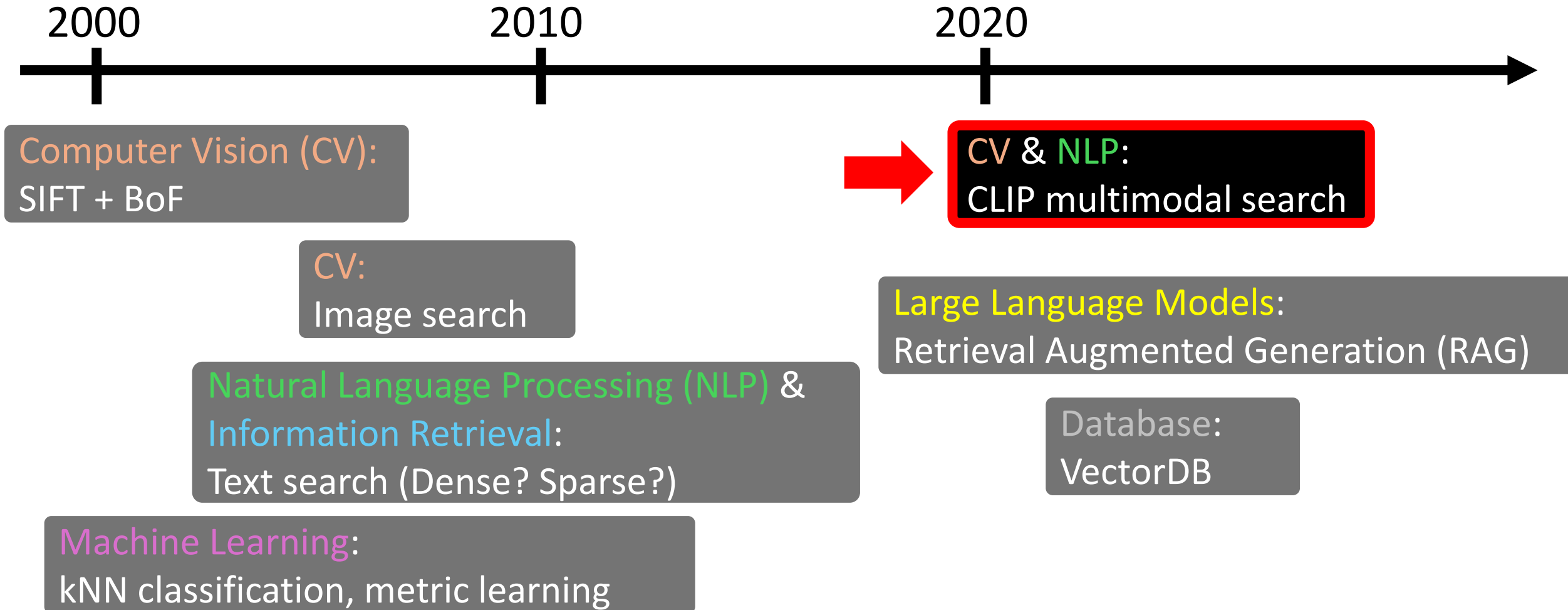
kNN classification



https://en.wikipedia.org/wiki/K-nearest_neighbors_algorithm

- The most straightforward approach for classification
- Given a query, find the closest sample from training data, and report its label
- Although super simple, it's actually effective if the embedding is good and #samples are large
- We can just run ANN search
- Complex ML -> Simple but large-scale search

What tasks has ANN been used for?



CLIP multimodal search

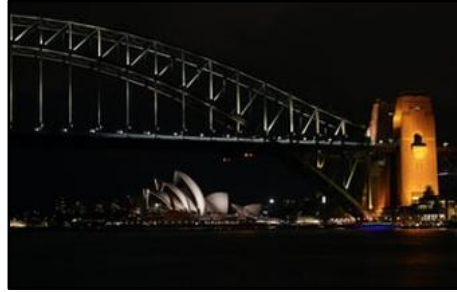
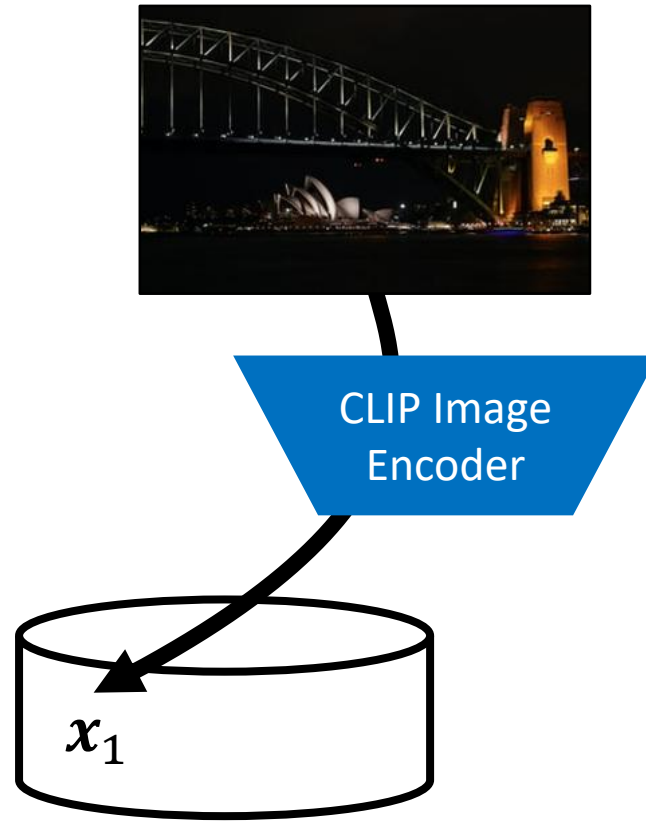
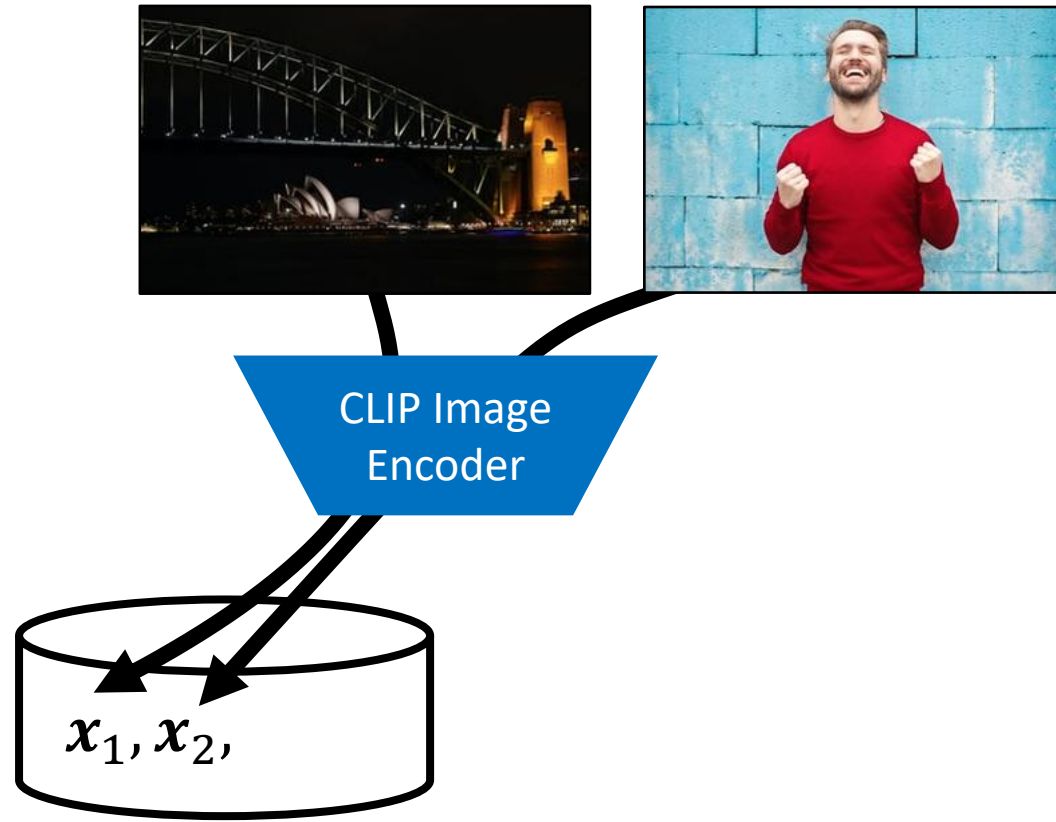


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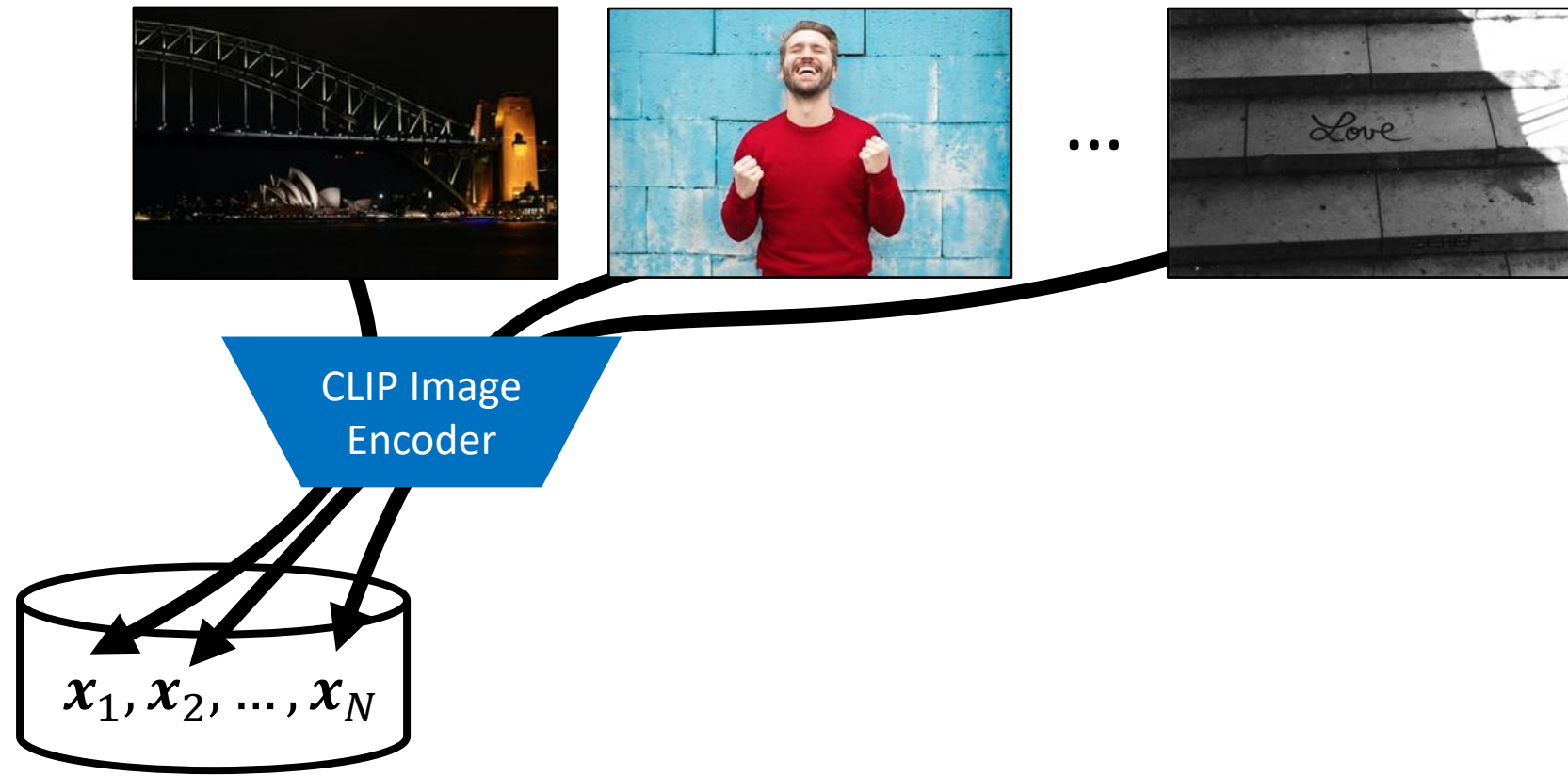
CLIP multimodal search



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CLIP multimodal search

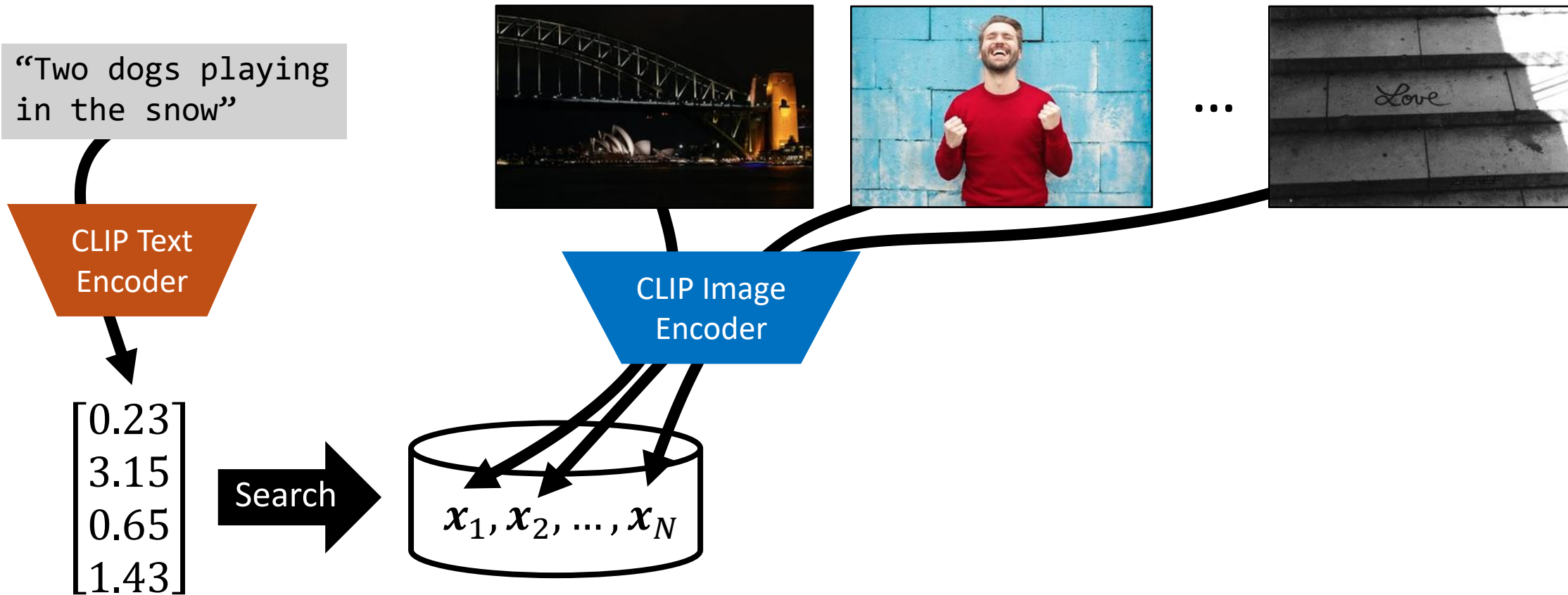


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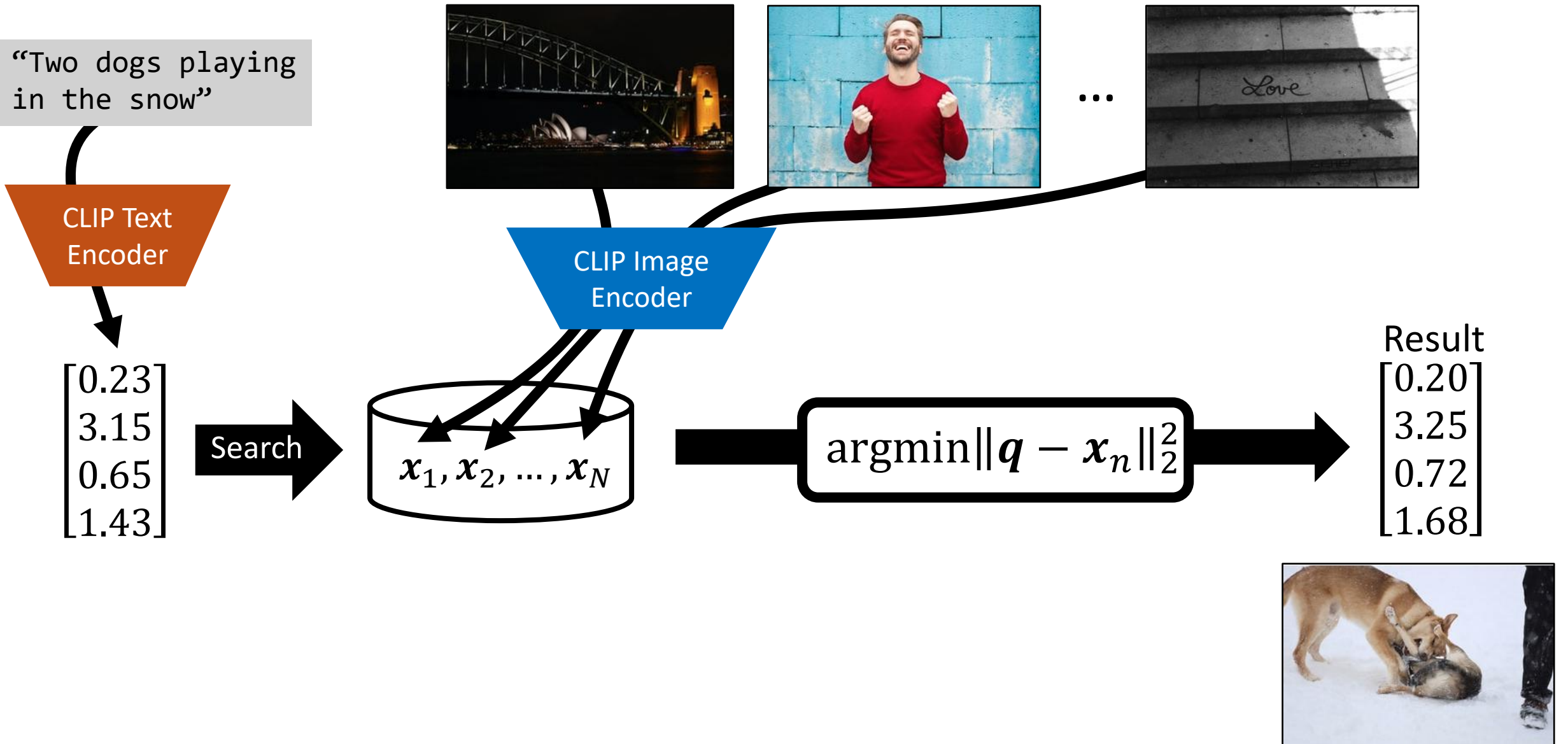
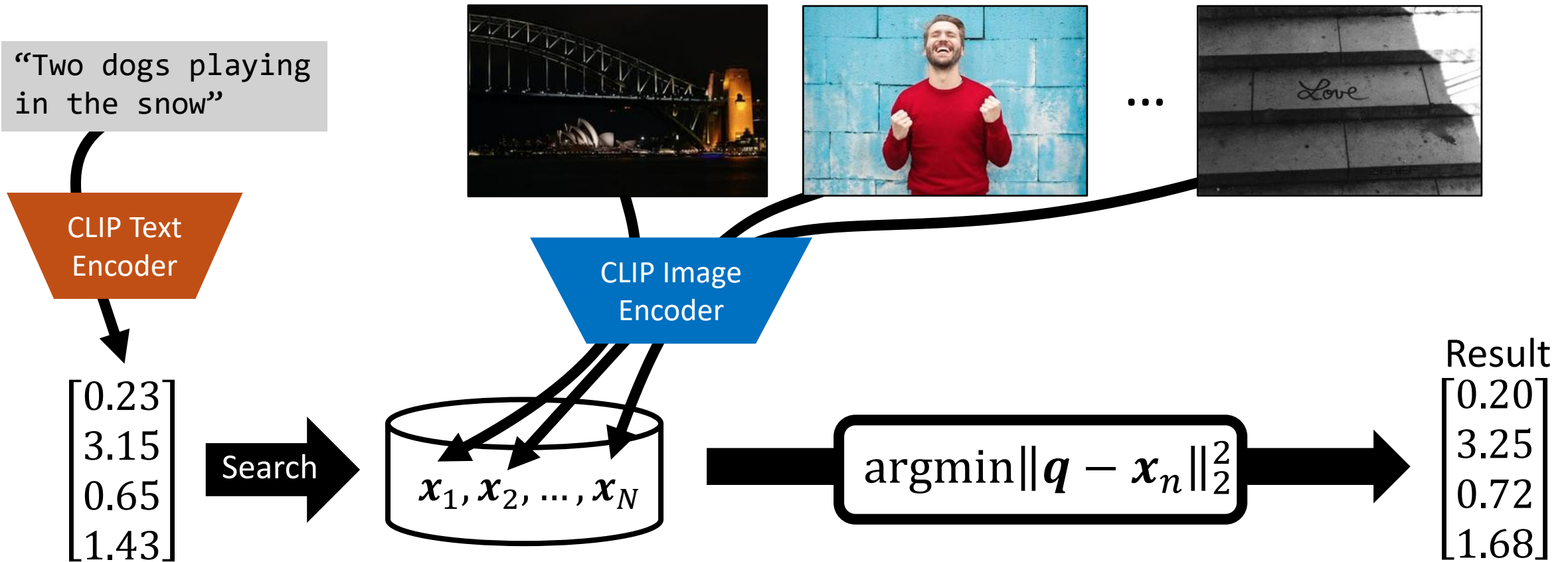


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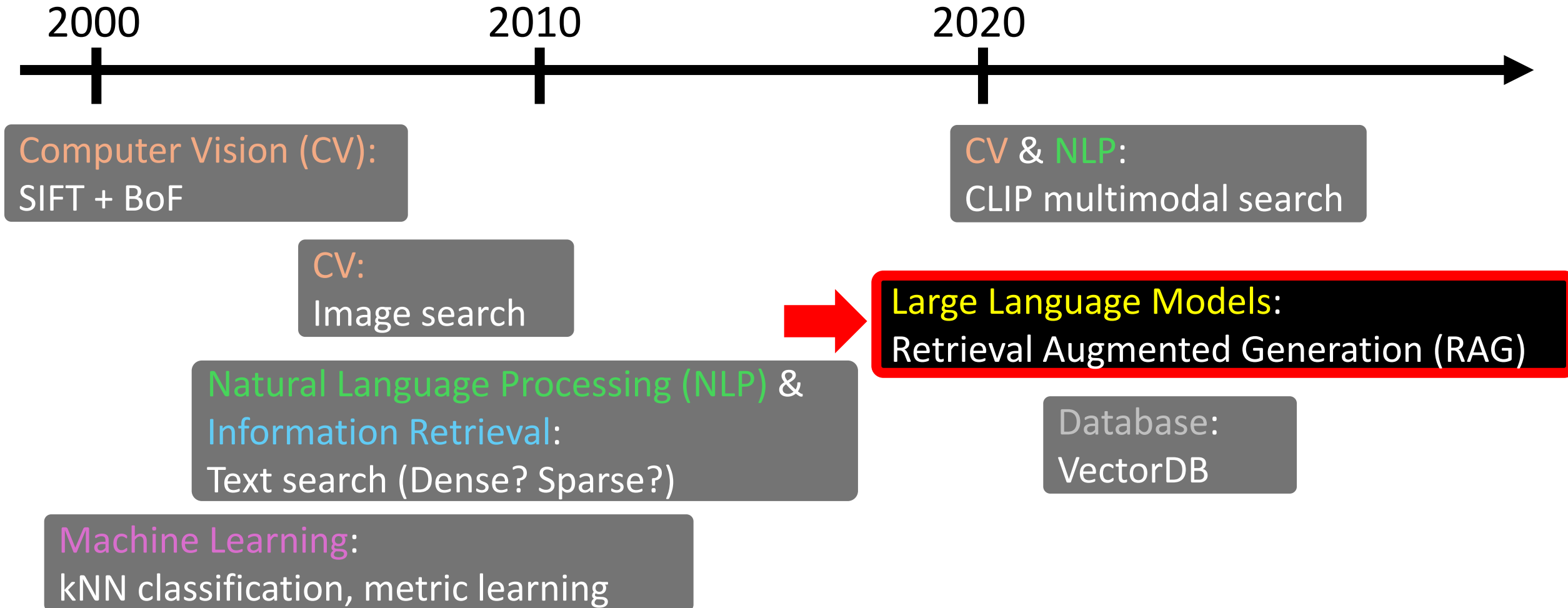
- CLIP enables us to compare images and texts
- Encoder determines **the upper bound** of the accuracy of the system
- ANN determines a **trade-off** between accuracy, runtime, and memory

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What tasks has ANN been used for?



RAG: LLM + embedding

"Who won curling
gold at the 2022
Winter Olympics?"



ChatGPT 3.5
(trained in 2021)

RAG: LLM + embedding

"Who won curling gold at the 2022 Winter Olympics?"

Ask



ChatGPT 3.5
(trained in 2021)

"I'm sorry, but as an AI language model, I don't have information about the future events."



RAG: LLM + embedding

"Who won curling
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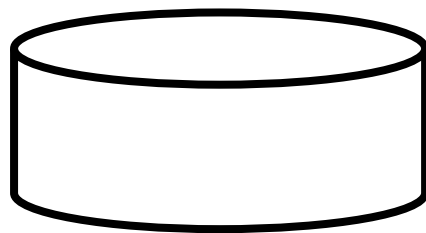
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"Chinami Yoshida\n\n=Personal..."

"Lviv bid for the 2022 Winter..."

⋮

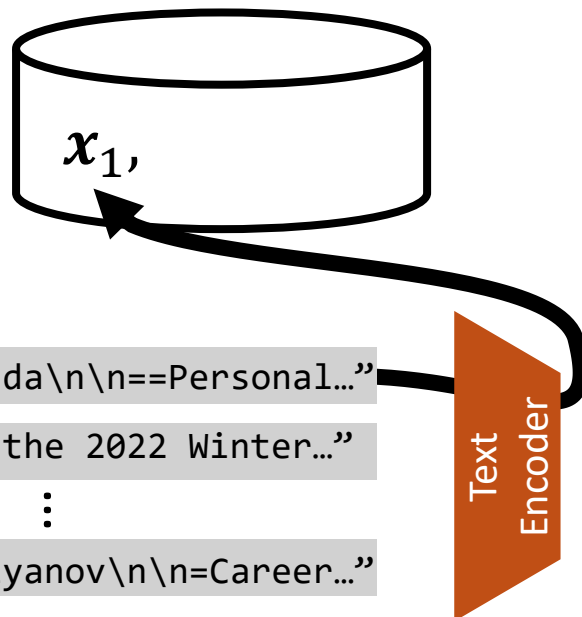
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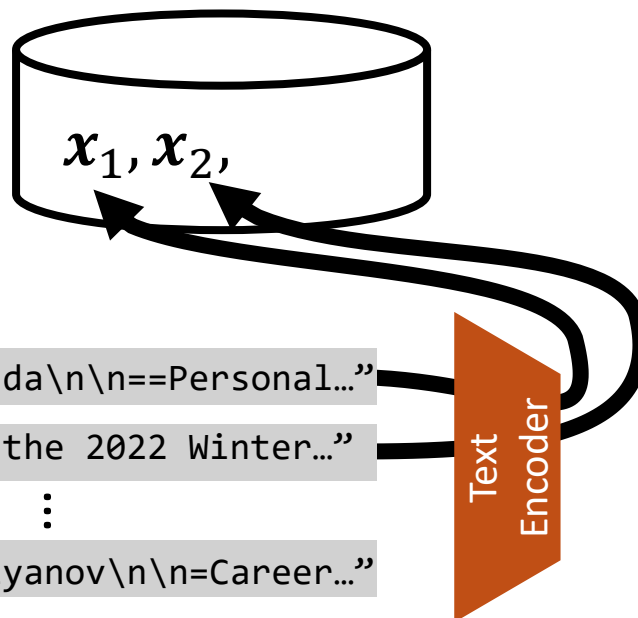


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Text
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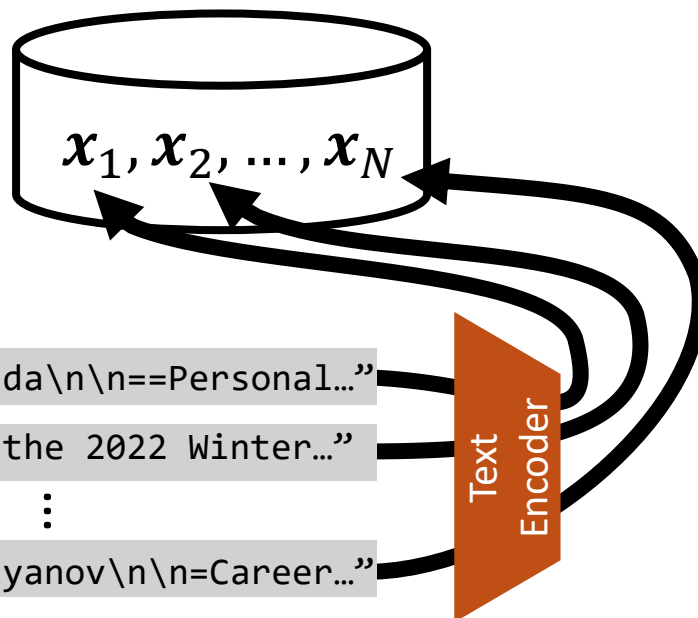


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Text
Encoder

$\begin{bmatrix} 0.23 \\ 3.15 \\ 0.65 \\ 1.43 \end{bmatrix}$



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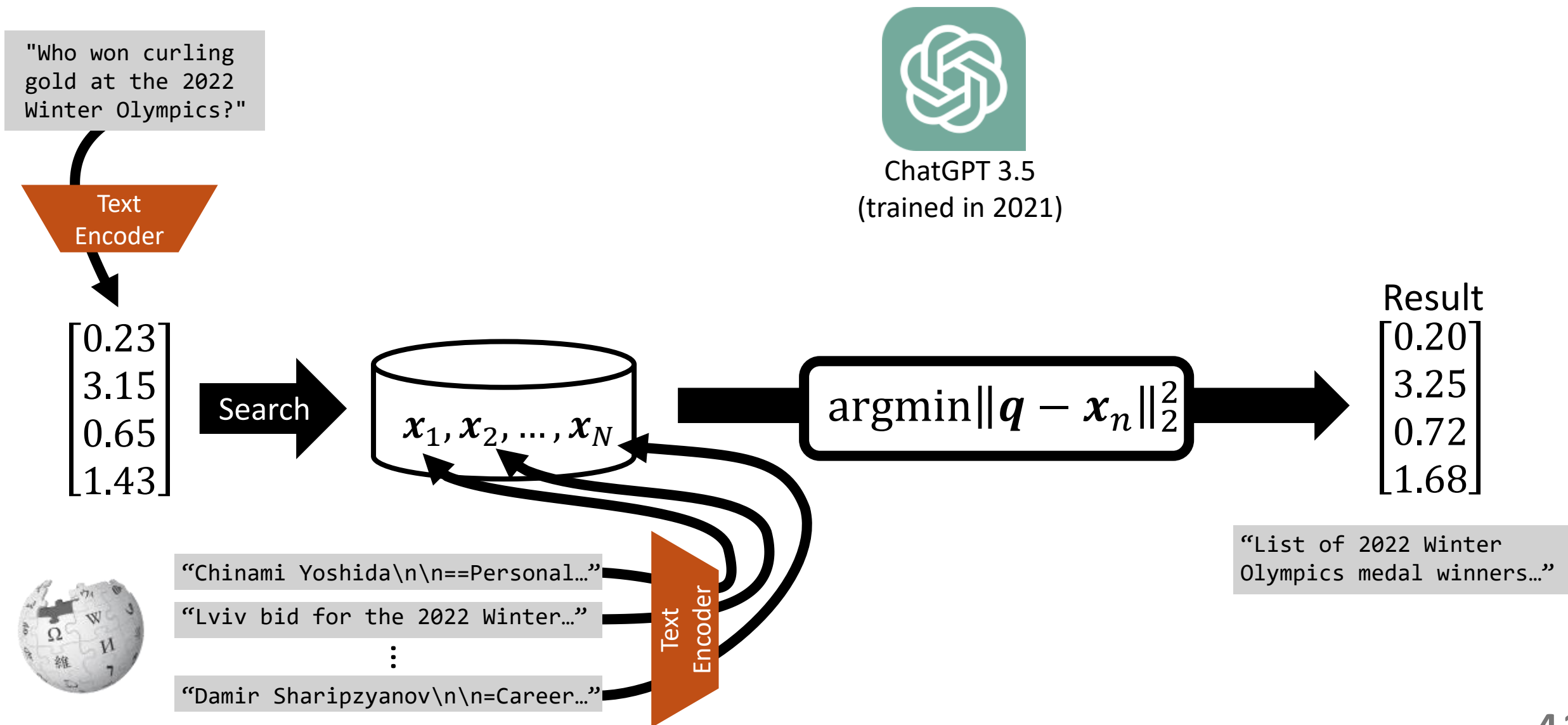
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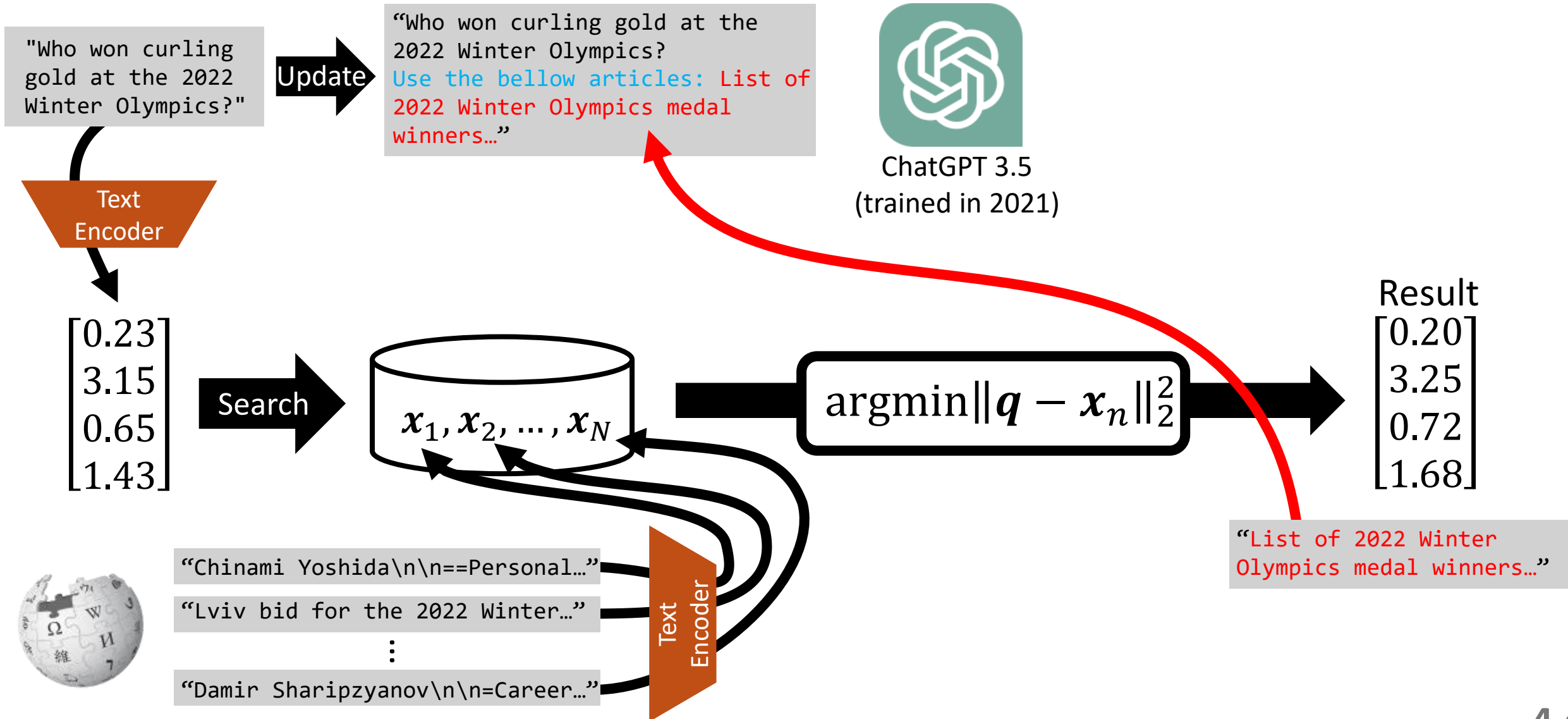


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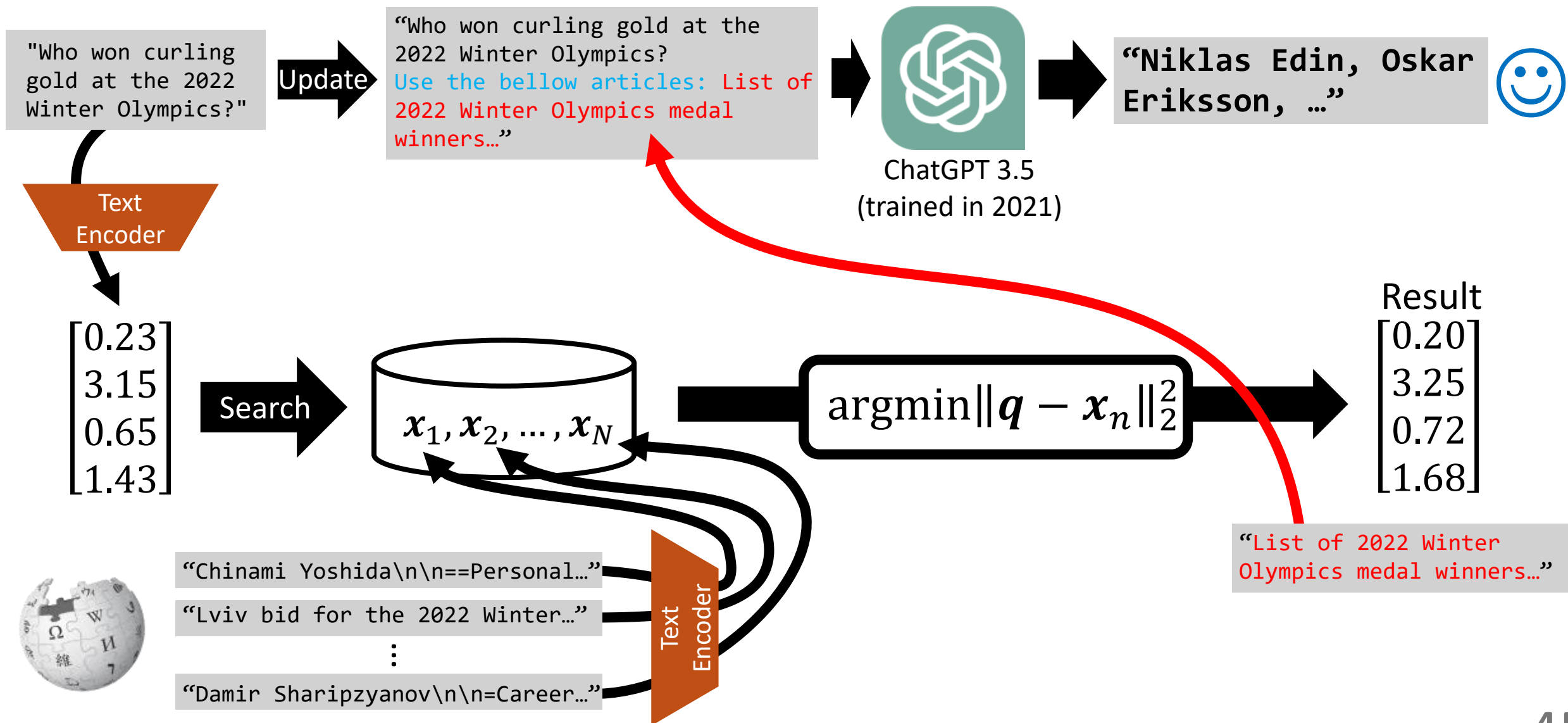
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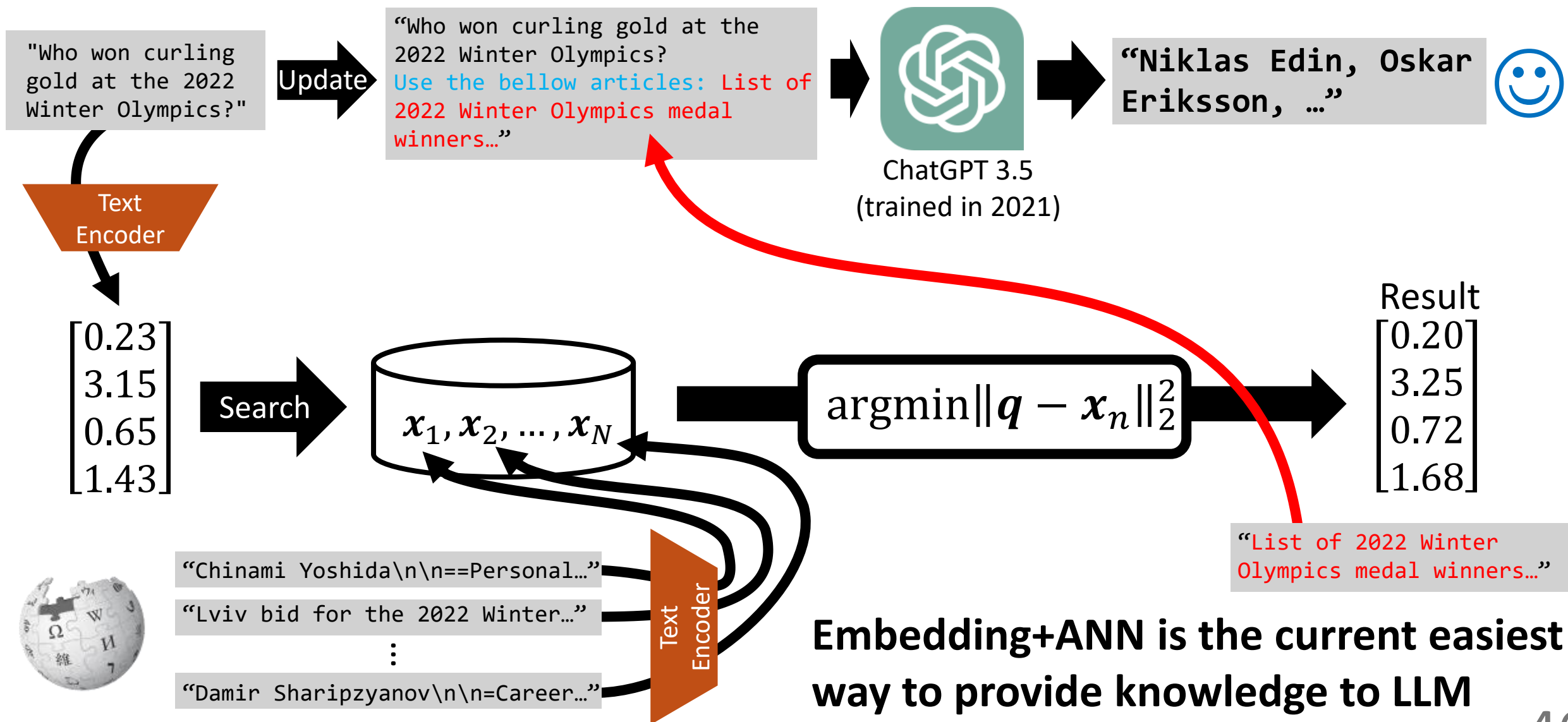
RAG: LLM + embedding



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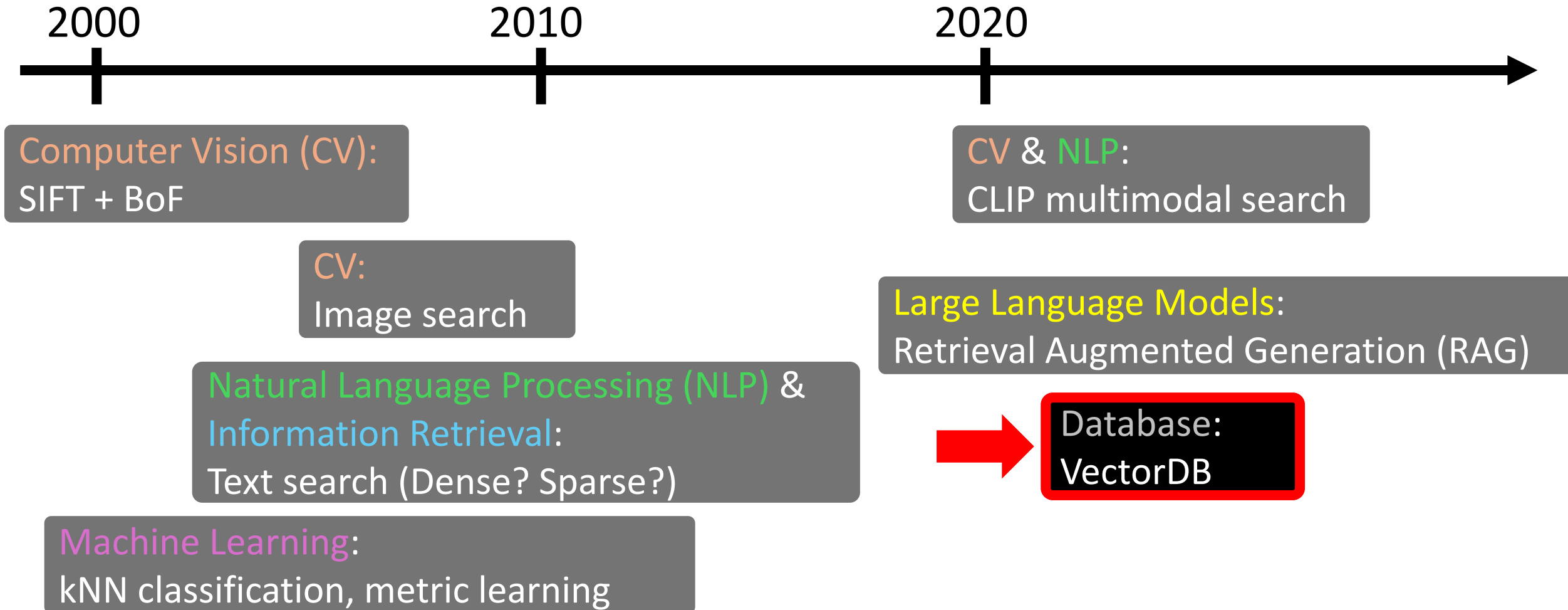


RAG: LLM + embedding



Embedding+ANN is the current easiest way to provide knowledge to LLM

What tasks has ANN been used for?



What is Vector DB?

Algorithm

- Scientific paper
- Math
- Often, by researchers

Product Quantization +
Inverted Index (PQ, IVFPQ)
[Jégou+, TPAMI 2011]

Hierarchical Navigable
Small World (HNSW)
[Malkov+, TPAMI 2019]

ScaNN (4-bit PQ)
[Guo+, ICML 2020]

Library

- Implementations of algorithms
- Usually, a search function only
- By researchers, developers, etc

faiss

NMSLIB

hnswlib

ScaNN

Service (e.g., vector DB)

- Library + (handling metadata, serving, scaling, IO, CRUD, etc)
- Usually, by companies

Pinecone

Milvus

Vald

Weaviate

Qdrant

jina

Vertex AI
Matching Engine

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faiss

Pinecone

Qdrant

hnswlib

One library may implement
multiple algorithms

☹️ “I benchmarked faiss”

😊 “I benchmarked PQ in faiss”

What is Vector DB?

One algorithm may be implemented in multiple libraries

- Algorithms
- Math
- Often, by researchers, developers, etc
- Usually, a search function only

Product Quantization
Inverted Index (PQ, IV)
[Jégou+, TPAMI 2011]

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faiss

Often, one library = one algorithm

ScaNN

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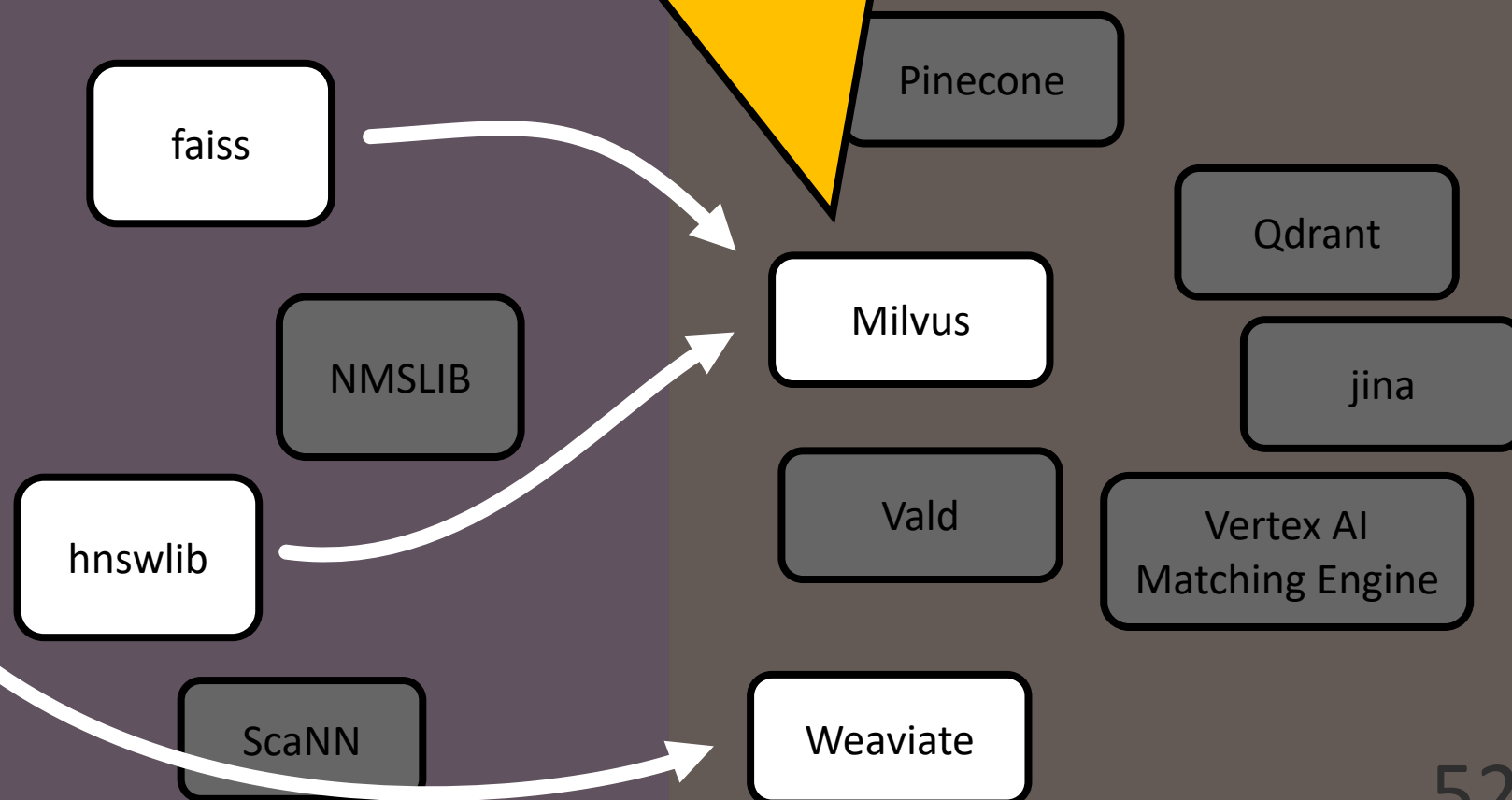
Product Quantization +
Inverted Index (PQ, IVFPQ)
[Jégou+, TPAMI 2011]

Hierarchical Navigable
Small World (HNSW)
[Malkov+, TPAMI 2019]

... or re-implement
algorithms from
scratch (e.g., by Go)

One service may use some libraries

- Usually, a search function
- By researchers, developers, etc.
- Handling metadata, indexing, scaling, IO, CRUD, etc)
- Usually, by companies



What is Vector DB?

Algorithm

- Scientific paper
- Math
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Inverted Index (PQ, IVFPQ)
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Vald

aviate

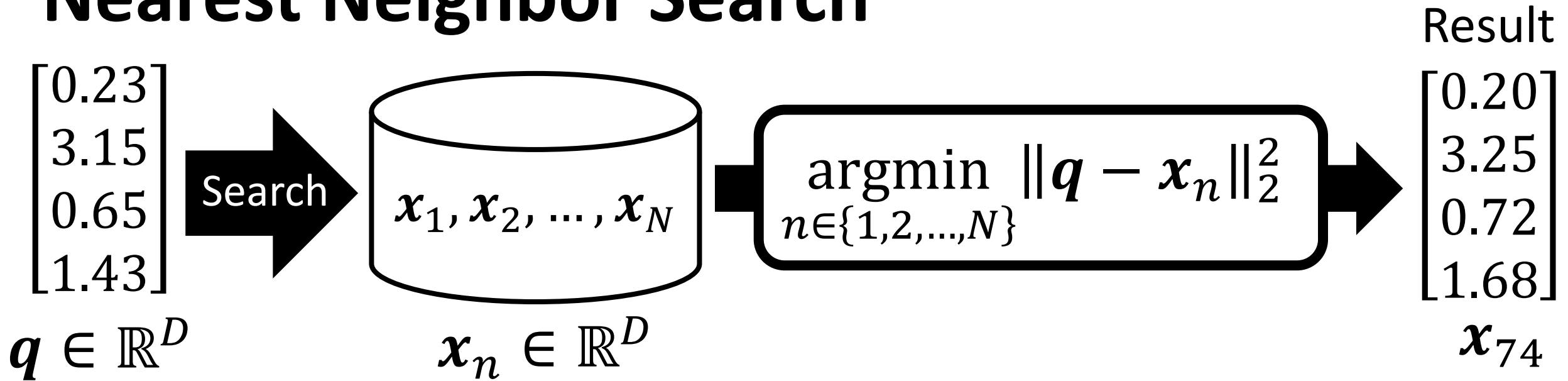
➤ Recently, the DB community has been the most active

➤ The CV community has settled down a bit 🤖

Outline

1. History from an applications perspective
- 2. Importance of implementation:
nearest neighbor search in faiss**
3. Basics of modern baseline: graph-based search

Nearest Neighbor Search



- Before trying ANN, we should try NN
- Introduce a naïve implementation
- Introduce a fast implementation: **Faiss** library
- Experience the drastic difference between the two implementations

M D -dim query vectors $\mathcal{Q} = \{\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_M\}$
 N D -dim database vectors $\mathcal{X} = \{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_N\}$ $M \ll N$
Task : Given $\mathbf{q} \in \mathcal{Q}$ and $\mathbf{x} \in \mathcal{X}$, compute $\|\mathbf{q} - \mathbf{x}\|_2^2$

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Naïve impl.

```
def l2sqr(q, x):  
    diff = 0.0  
    for (d = 0; d < D; ++d):  
        diff += (q[d] - x[d])**2  
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```

```
parfor q in Q:
```

```
    for x in X:
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```
        l2sqr(q, x)
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Parallelize
query-side

Select min by heap,
but omit it now

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faiss impl.

Disclaimer: the code used in the explanation is from several years ago and is not up to date.

if $M < 20$:

 compute $\|\mathbf{q} - \mathbf{x}\|_2^2$ by SIMD

else :

 compute $\|\mathbf{q} - \mathbf{x}\|_2^2 = \|\mathbf{q}\|_2^2 - 2\mathbf{q}^\top \mathbf{x} + \|\mathbf{x}\|_2^2$ by BLAS

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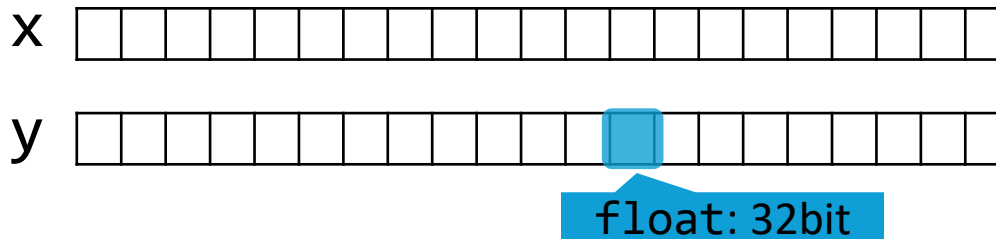
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$\|x - y\|_2^2$ by SIMD

Rename variables for the sake of explanation

D=31



Ref.

```
def l2sqr(x, y):  
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    for (d = 0; d < D; ++d):  
        diff += (x[d] - y[d])**2  
    return diff
```

```
float fvec_L2sqr (const float * x,  
                 const float * y,  
                 size_t d)  
{  
    __m256 msum1 = _mm256_setzero_ps();  
  
    while (d >= 8) {  
        __m256 mx = _mm256_loadu_ps (x); x += 8;  
        __m256 my = _mm256_loadu_ps (y); y += 8;  
        const __m256 a_m_b1 = mx - my;  
        msum1 += a_m_b1 * a_m_b1;  
        d -= 8;  
    }  
  
    __m128 msum2 = _mm256_extractf128_ps(msum1, 1);  
    msum2 += _mm256_extractf128_ps(msum1, 0);  
  
    if (d >= 4) {  
        __m128 mx = _mm_loadu_ps (x); x += 4;  
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        const __m128 a_m_b1 = mx - my;  
        msum2 += a_m_b1 * a_m_b1;  
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    }  
  
    if (d > 0) {  
        __m128 mx = masked_read (d, x);  
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    msum2 = _mm_hadd_ps (msum2, msum2);  
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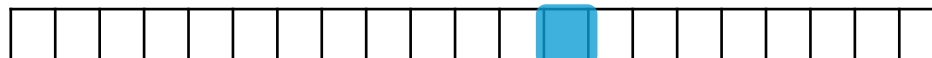
```
}
```

D=31

x



y



float: 32bit

mx



my



```
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    for (d = 0; d < D; ++d):  
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Ref.

- 256bit SIMD Register
- Process eight floats at once

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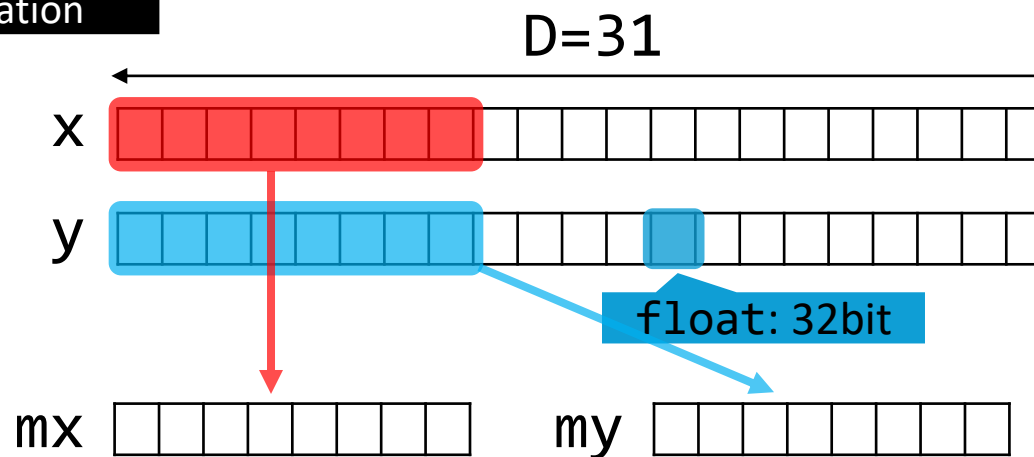
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    while (d >= 8) {  
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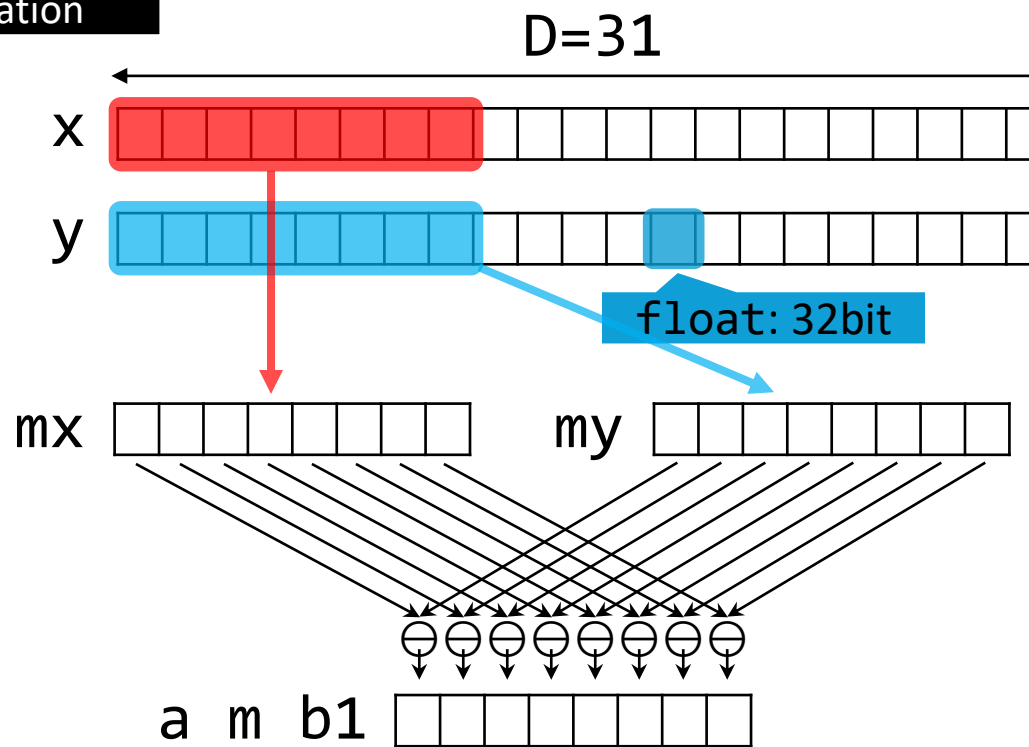
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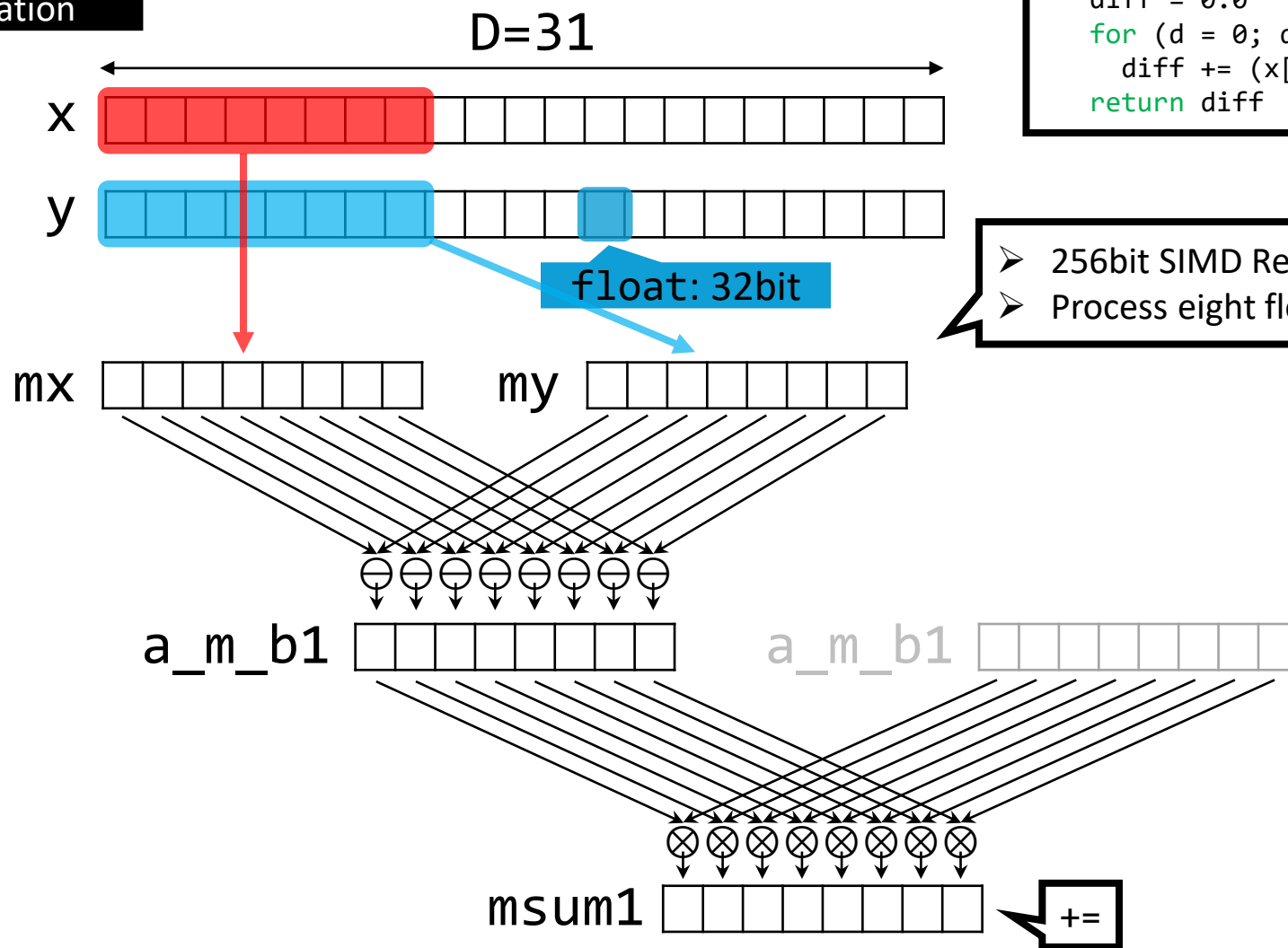

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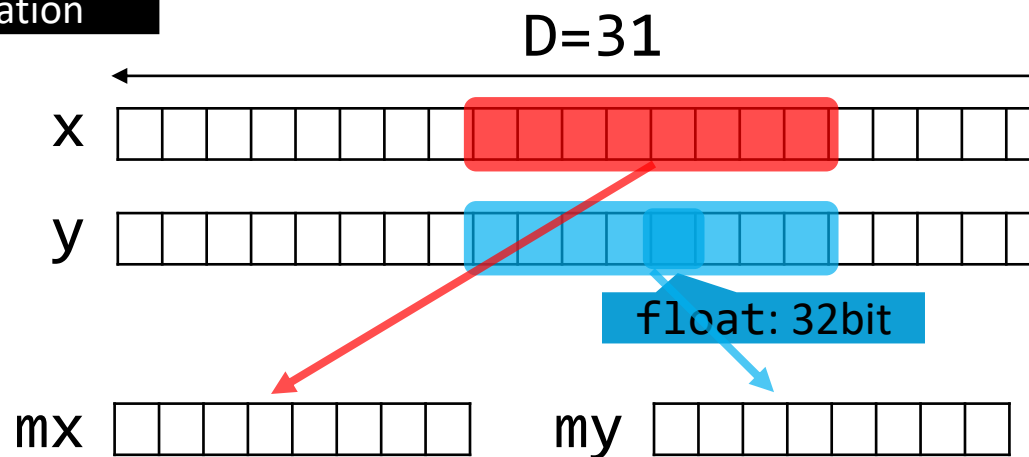
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msum1

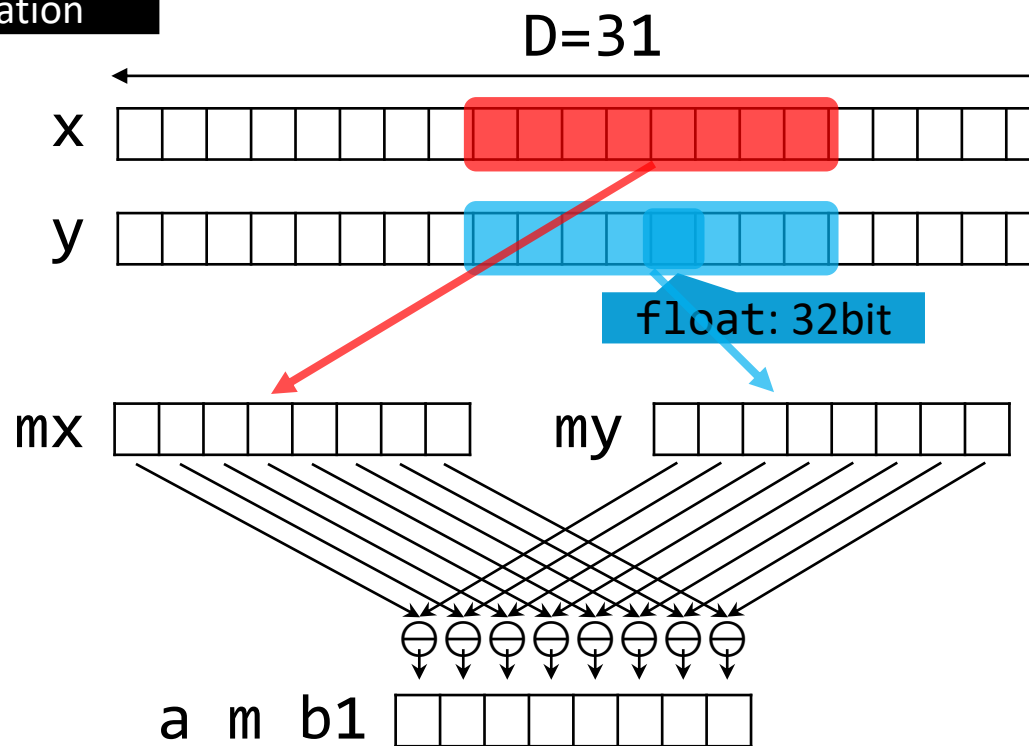
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$msum1$ [] [] [] [] [] [] [] [] $\boxed{+=}$

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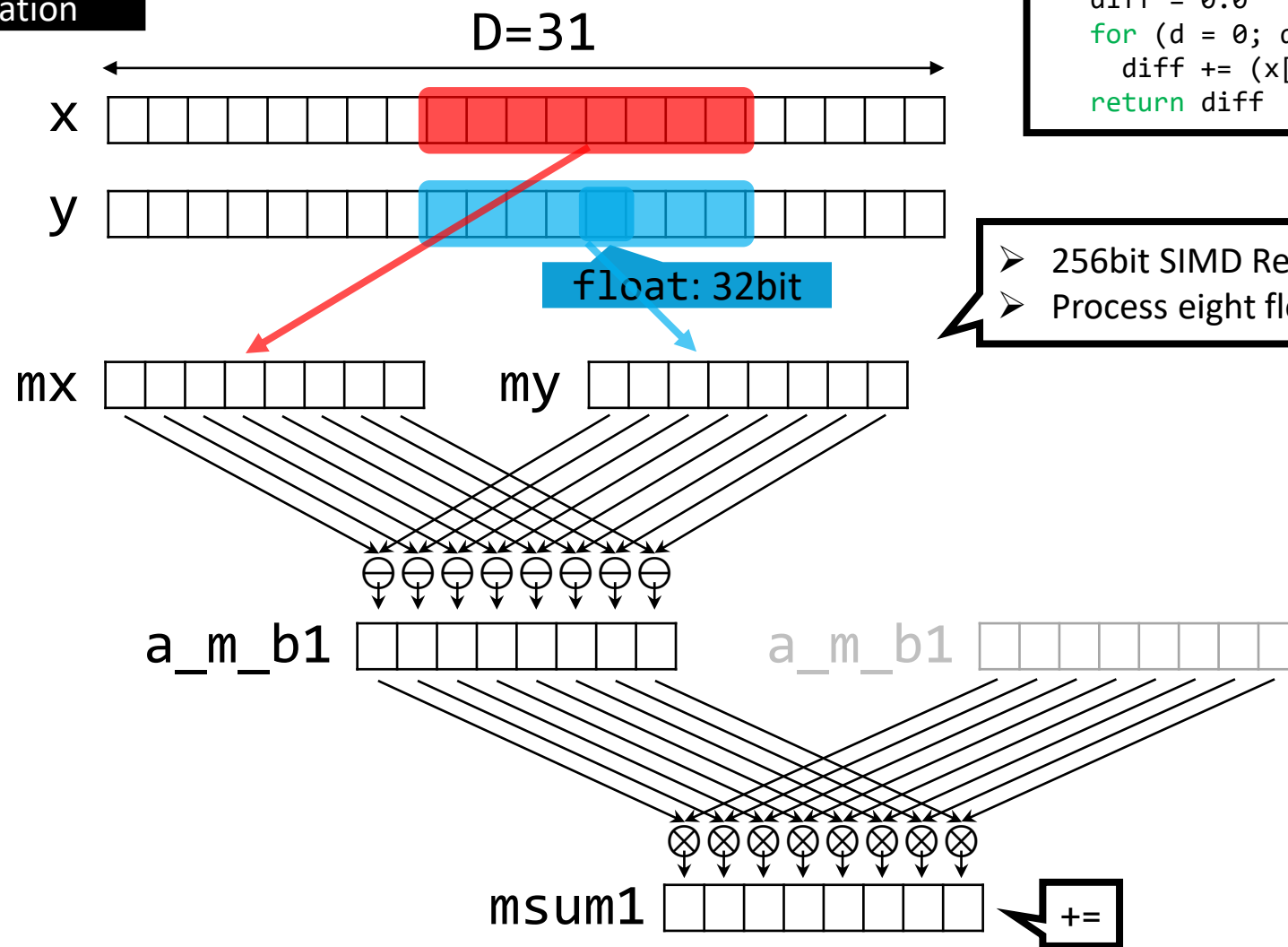
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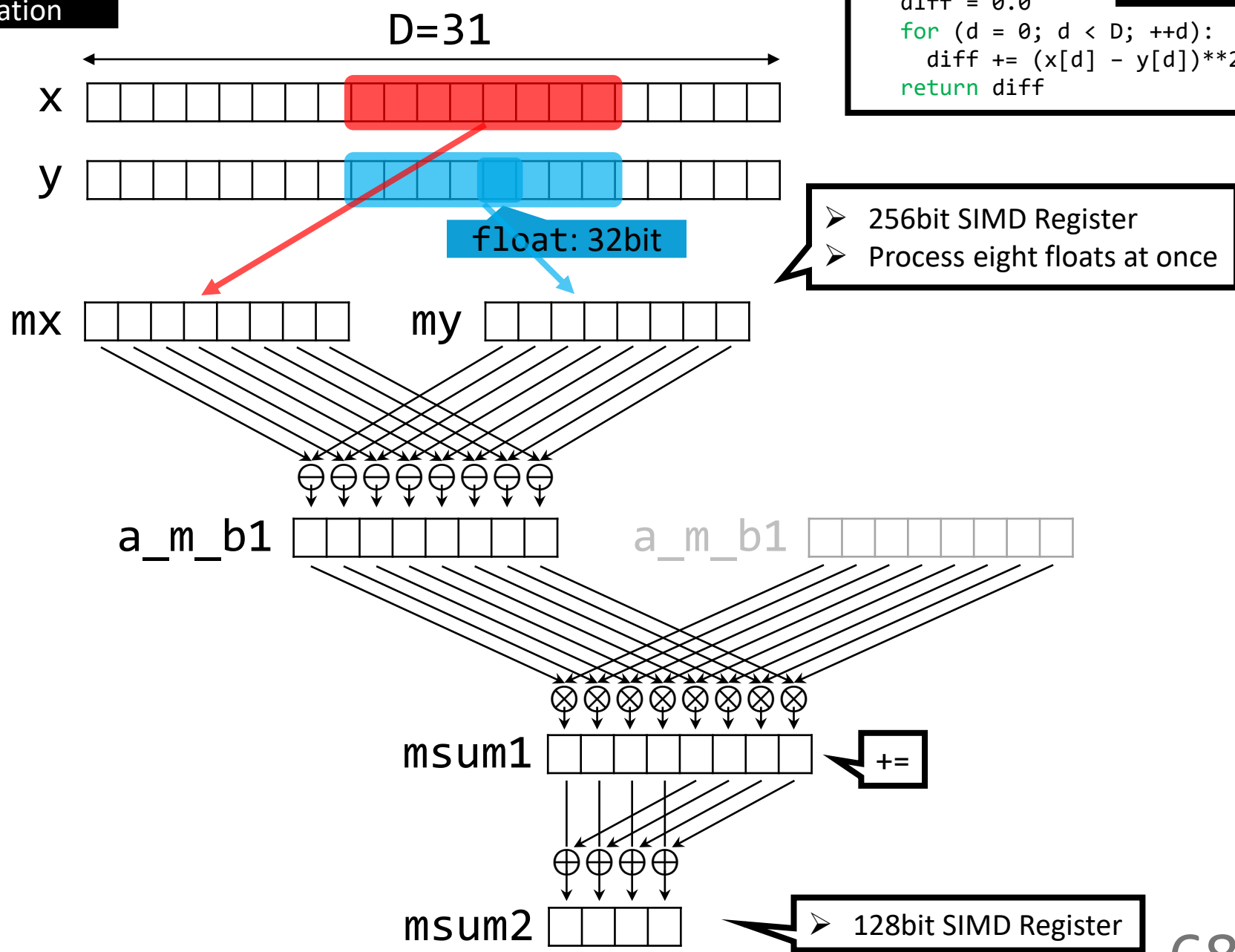
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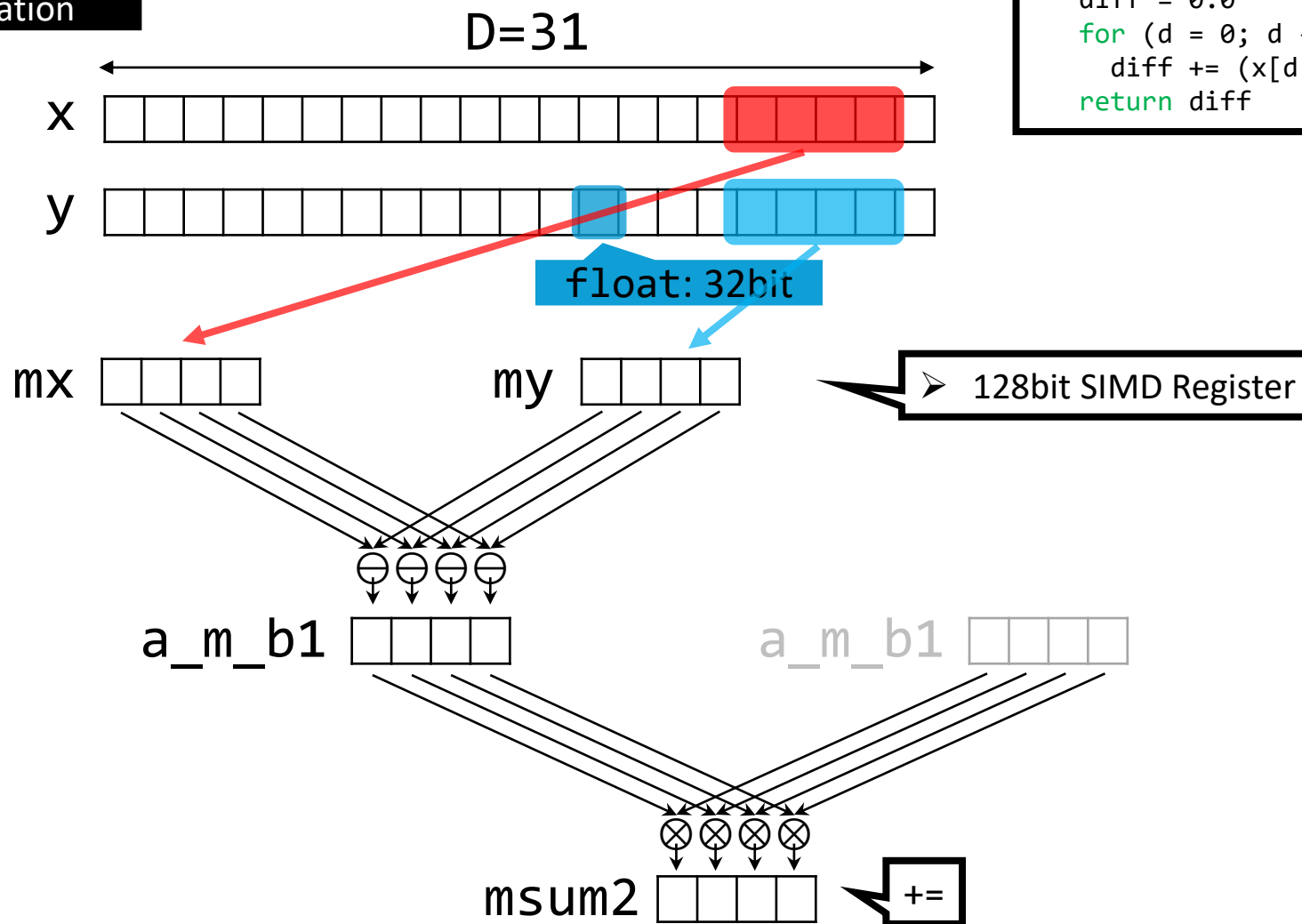
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        __m256 mx = _mm256_loadu_ps (x); x += 8;
        __m256 my = _mm256_loadu_ps (y); y += 8;
        const __m256 a_m_b1 = mx - my;
        msum1 += a_m_b1 * a_m_b1;
        d -= 8;
    }

    __m128 msum2 = _mm256_extractf128_ps(msum1, 1);
    msum2 += _mm256_extractf128_ps(msum1, 0);

    if (d >= 4) {
        __m128 mx = _mm_loadu_ps (x); x += 4;
        __m128 my = _mm_loadu_ps (y); y += 4;
        const __m128 a_m_b1 = mx - my;
        msum2 += a_m_b1 * a_m_b1;
        d -= 4;
    }

    if (d > 0) {
        __m128 mx = masked_read (d, x);
        __m128 my = masked_read (d, y);
        __m128 a_m_b1 = mx - my;
        msum2 += a_m_b1 * a_m_b1;
    }

    msum2 = _mm_hadd_ps (msum2, msum2);
    msum2 = _mm_hadd_ps (msum2, msum2);
    return _mm_cvtss_f32 (msum2);
}
```



$\|x - y\|_2^2$ by SIMD

Rename variables for the sake of explanation

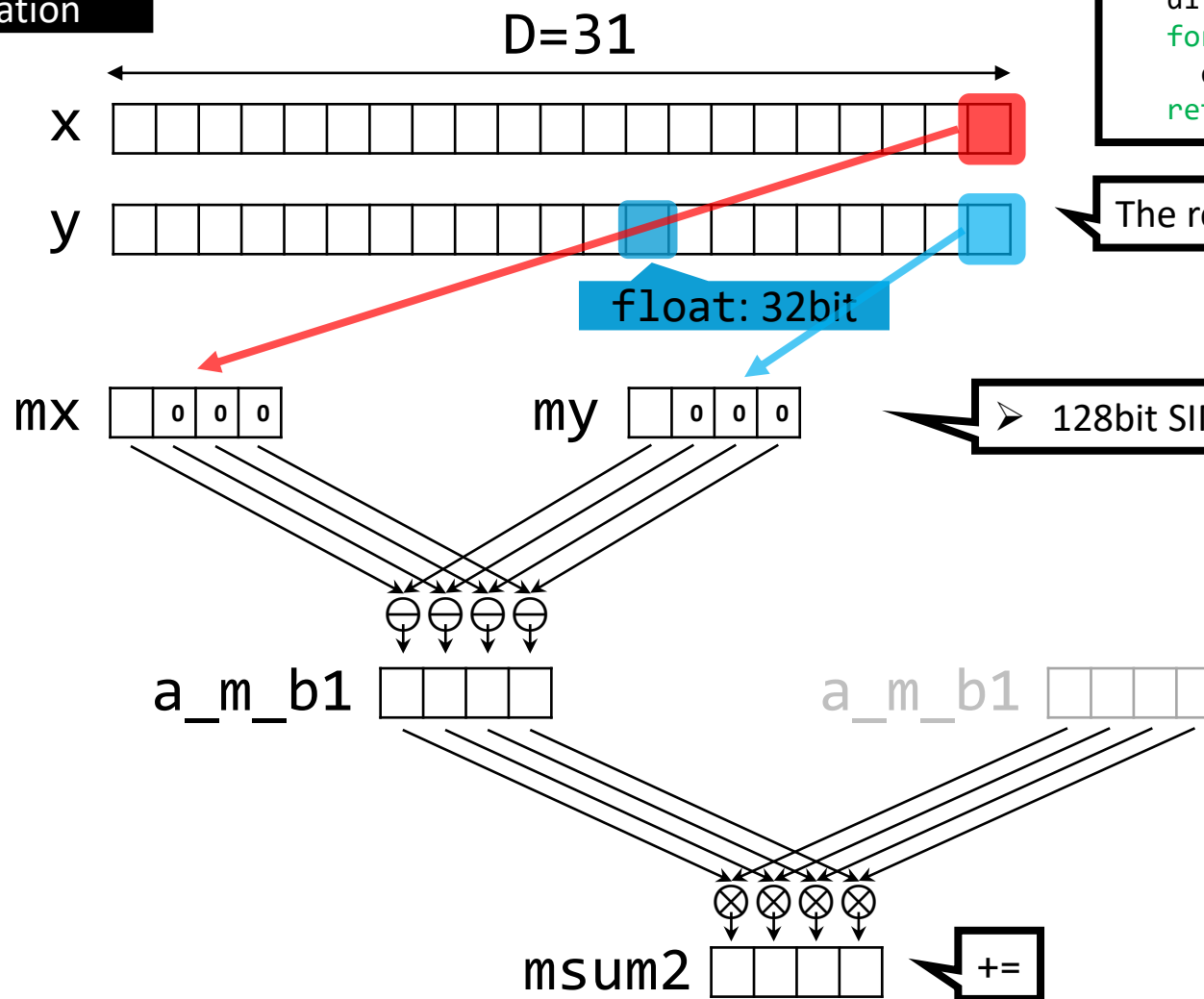
Ref.

```
def l2sqr(x, y):
    diff = 0.0
    for (d = 0; d < D; ++d):
        diff += (x[d] - y[d])**2
    return diff
```

The rest

➤ 128bit SIMD Register

+=



```
float fvec_L2sqr (const float * x,
                  const float * y,
                  size_t d)
{
    __m256 msum1 = _mm256_setzero_ps();

    while (d >= 8) {
        __m256 mx = _mm256_loadu_ps (x); x += 8;
        __m256 my = _mm256_loadu_ps (y); y += 8;
        const __m256 a_m_b1 = mx - my;
        msum1 += a_m_b1 * a_m_b1;
        d -= 8;
    }

    __m128 msum2 = _mm256_extractf128_ps(msum1, 1);
    msum2 += _mm256_extractf128_ps(msum1, 0);

    if (d >= 4) {
        __m128 mx = _mm_loadu_ps (x); x += 4;
        __m128 my = _mm_loadu_ps (y); y += 4;
        const __m128 a_m_b1 = mx - my;
        msum2 += a_m_b1 * a_m_b1;
        d -= 4;
    }

    if (d > 0) {
        __m128 mx = masked_read (d, x);
        __m128 my = masked_read (d, y);
        __m128 a_m_b1 = mx - my;
        msum2 += a_m_b1 * a_m_b1;
    }

    msum2 = _mm_hadd_ps (msum2, msum2);
    msum2 = _mm_hadd_ps (msum2, msum2);
    return _mm_cvtss_f32 (msum2);
}
```

$\|x - y\|_2^2$ by SIMD

Rename variables for the sake of explanation

```
float fvec_L2sqr (const float * x,
                  const float * y,
                  size_t d)
{
    __m256 msum1 = _mm256_setzero_ps();

    while (d >= 8) {
        __m256 mx = _mm256_loadu_ps (x); x += 8;
        __m256 my = _mm256_loadu_ps (y); y += 8;
        const __m256 a_m_b1 = mx - my;
        msum1 += a_m_b1 * a_m_b1;
        d -= 8;
    }

    __m128 msum2 = _mm256_extractf128_ps(msum1, 1);
    msum2 += _mm256_extractf128_ps(msum1, 0);

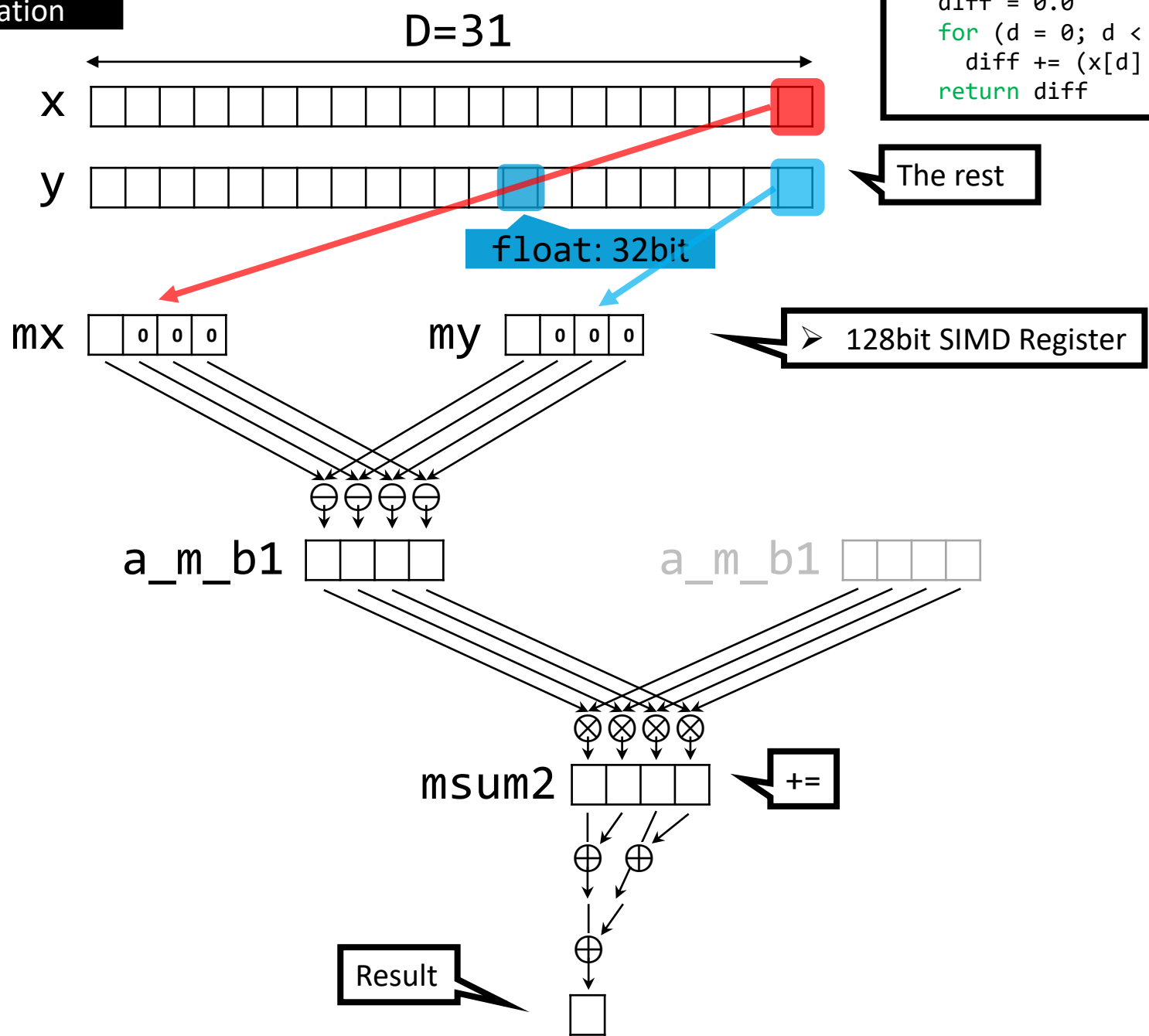
    if (d >= 4) {
        __m128 mx = _mm_loadu_ps (x); x += 4;
        __m128 my = _mm_loadu_ps (y); y += 4;
        const __m128 a_m_b1 = mx - my;
        msum2 += a_m_b1 * a_m_b1;
        d -= 4;
    }

    if (d > 0) {
        __m128 mx = masked_read (d, x);
        __m128 my = masked_read (d, y);
        __m128 a_m_b1 = mx - my;
        msum2 += a_m_b1 * a_m_b1;
    }

    msum2 = _mm_hadd_ps (msum2, msum2);
    msum2 = _mm_hadd_ps (msum2, msum2);
    return _mm_cvtss_f32 (msum2);
}
```

Ref.

```
def l2sqr(x, y):
    diff = 0.0
    for (d = 0; d < D; ++d):
        diff += (x[d] - y[d])**2
    return diff
```



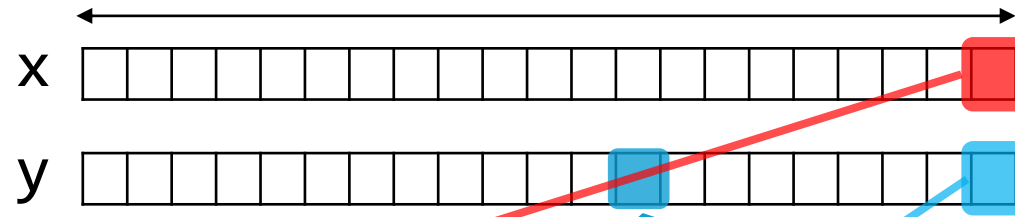
$\|x - y\|_2^2$ by SIMD

Rename variables for the sake of explanation

Ref.

```
def l2sqr(x, y):  
    diff = 0.0  
    for (d = 0; d < D; ++d):  
        diff += (x[d] - y[d])**2  
    return diff
```

D=31



- SIMD codes of faiss are simple and easy to read
- Being able to read SIMD codes comes in handy sometimes; *why this impl is super fast*
- Another example of SIMD L2sqr from HNSW:

https://github.com/nmslib/hnswlib/blob/master/hnswlib/space_l2.h

msum2

+=

Result

```
__m128 mx = masked_read(d, x);  
__m128 my = masked_read(d, y);  
__m128 a_m_b1 = mx - my;  
msum2 += a_m_b1 * a_m_b1;  
}
```

```
msum2 = _mm_hadd_ps(msum2, msum2);  
msum2 = _mm_hadd_ps(msum2, msum2);  
return _mm_cvtss_f32(msum2);  
}
```

M D -dim query vectors $\mathcal{Q} = \{\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_M\}$

N D -dim database vectors $\mathcal{X} = \{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_N\}$

Task : Given $\mathbf{q} \in \mathcal{Q}$ and $\mathbf{x} \in \mathcal{X}$, compute $\|\mathbf{q} - \mathbf{x}\|_2^2$

Naïve impl.

```
def l2sqr(q, x):  
    diff = 0.0  
    for (d = 0; d < D; ++d):  
        diff += (q[d] - x[d])**2  
    return diff
```

```
parfor q in Q:  
    for x in X:  
        l2sqr(q, x)
```

Parallelize
query-side

Select min by heap,
but omit it now

faiss impl.

if $M < 20$:

compute $\|\mathbf{q} - \mathbf{x}\|_2^2$ by SIMD

else :

compute $\|\mathbf{q} - \mathbf{x}\|_2^2 = \|\mathbf{q}\|_2^2 - 2\mathbf{q}^\top \mathbf{x} + \|\mathbf{x}\|_2^2$ by BLAS

Compute $\|\mathbf{q} - \mathbf{x}\|_2^2 = \|\mathbf{q}\|_2^2 - 2\mathbf{q}^\top \mathbf{x} + \|\mathbf{x}\|_2^2$ with BLAS

Stack M D -dim query vectors to a $D \times M$ matrix: $Q = [\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_M] \in \mathbb{R}^{D \times M}$

Stack N D -dim database vectors to a $D \times N$ matrix: $X = [\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_N] \in \mathbb{R}^{D \times N}$

Compute tables

SIMD-accelerated function

q_norms = norms(Q) # $\|\mathbf{q}_1\|_2^2, \|\mathbf{q}_2\|_2^2, \dots, \|\mathbf{q}_M\|_2^2$

x_norms = norms(X) # $\|\mathbf{x}_1\|_2^2, \|\mathbf{x}_2\|_2^2, \dots, \|\mathbf{x}_N\|_2^2$

ip = sgemm_(Q, X, ...) # $Q^\top X$

Scan and sum

parfor (m = 0; m < M; ++m):

for (n = 0; n < N; ++n):

dist = q_norms[m] + x_norms[n] - 2 * ip[m][n]

- Matrix multiplication by BLAS
- Dominant if Q and X are large
- The difference of the background matters:
 - ✓ Intel MKL is 30% faster than OpenBLAS

$$\|\mathbf{q}_m - \mathbf{x}_n\|_2^2$$

$$\|\mathbf{q}_m\|_2^2$$

$$\|\mathbf{x}_n\|_2^2$$

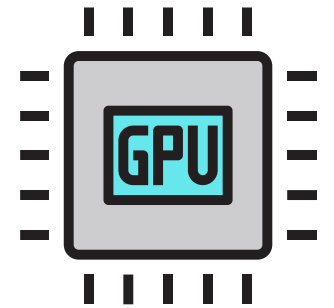
$$(Q^\top X)_{mn}$$

NN in GPU (faiss-gpu) is 10x faster than NN in CPU (faiss-cpu)

Benchmark: <https://github.com/facebookresearch/faiss/wiki/Low-level-benchmarks>

- NN-GPU always compute $\|\mathbf{q}\|_2^2 - 2\mathbf{q}^\top \mathbf{x} + \|\mathbf{x}\|_2^2$
- k-means for 1M vectors (D=256, K=20000)
 - ✓ 11 min on CPU
 - ✓ 55 sec on 1 Pascal-class P100 GPU (float32 math)
 - ✓ 34 sec on 1 Pascal-class P100 GPU (float16 math)
 - ✓ 21 sec on 4 Pascal-class P100 GPUs (float32 math)
 - ✓ 16 sec on 4 Pascal-class P100 GPUs (float16 math)
- If GPU is available and its memory is enough, try GPU-NN
- The behavior is little bit different (e.g., a restriction for top-k)

➤ x10 faster

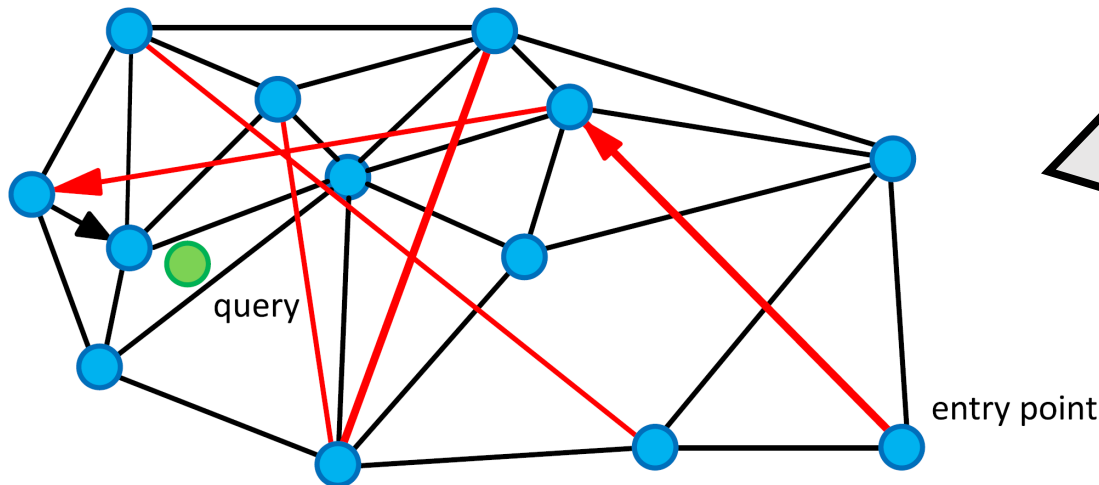


Outline

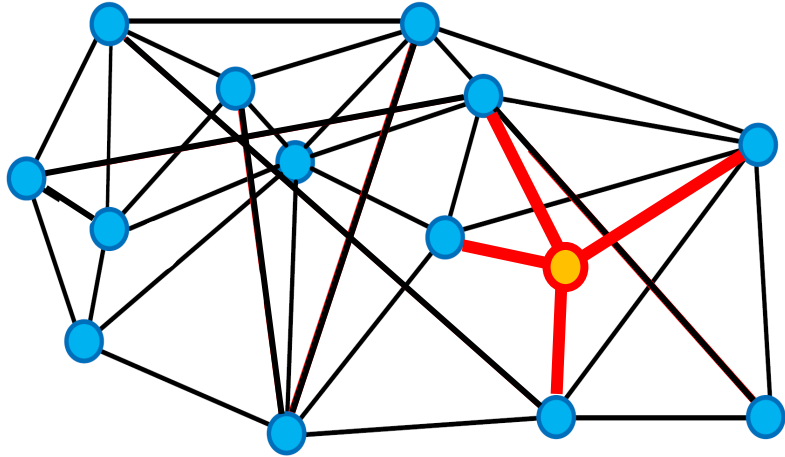
1. History from an applications perspective
2. Importance of implementation:
nearest neighbor search in faiss
3. Basics of modern baseline: graph-based search

Graph search

- De facto standard if all data can be loaded on memory
- Fast and accurate for real-world data
- Important for billion-scale situation as well
 - ✓ Graph-search is a building block for billion-scale systems

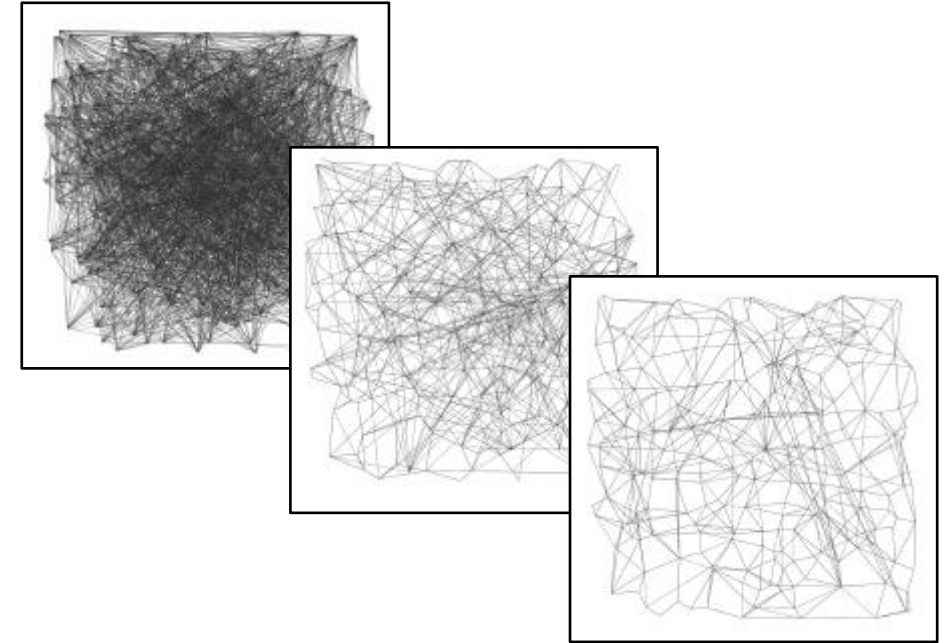


- Traverse graph towards the query
- Seems intuitive, but not so much easy to understand
- Review the algorithm carefully



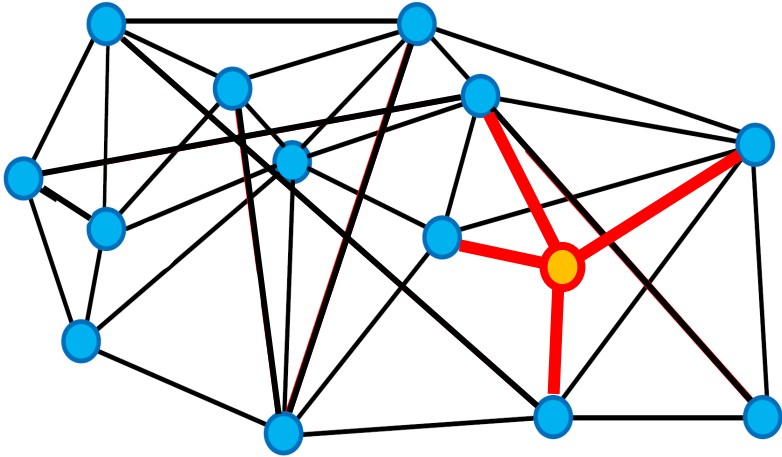
Increment approach

- Add a new item to the current graph incrementally



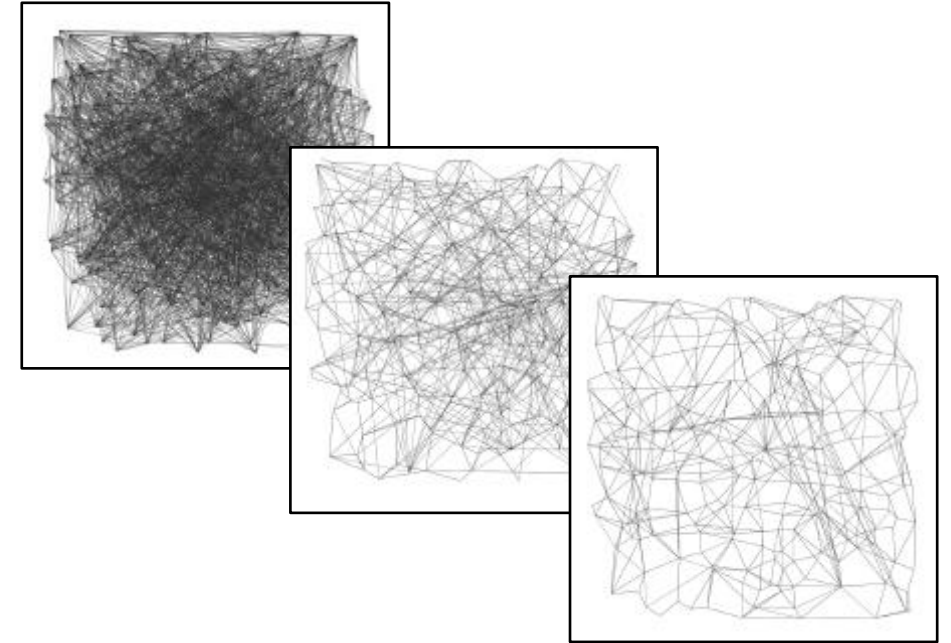
Refinement approach

- Iteratively refine an initial graph



Increment approach

- Add a new item to the current graph incrementally

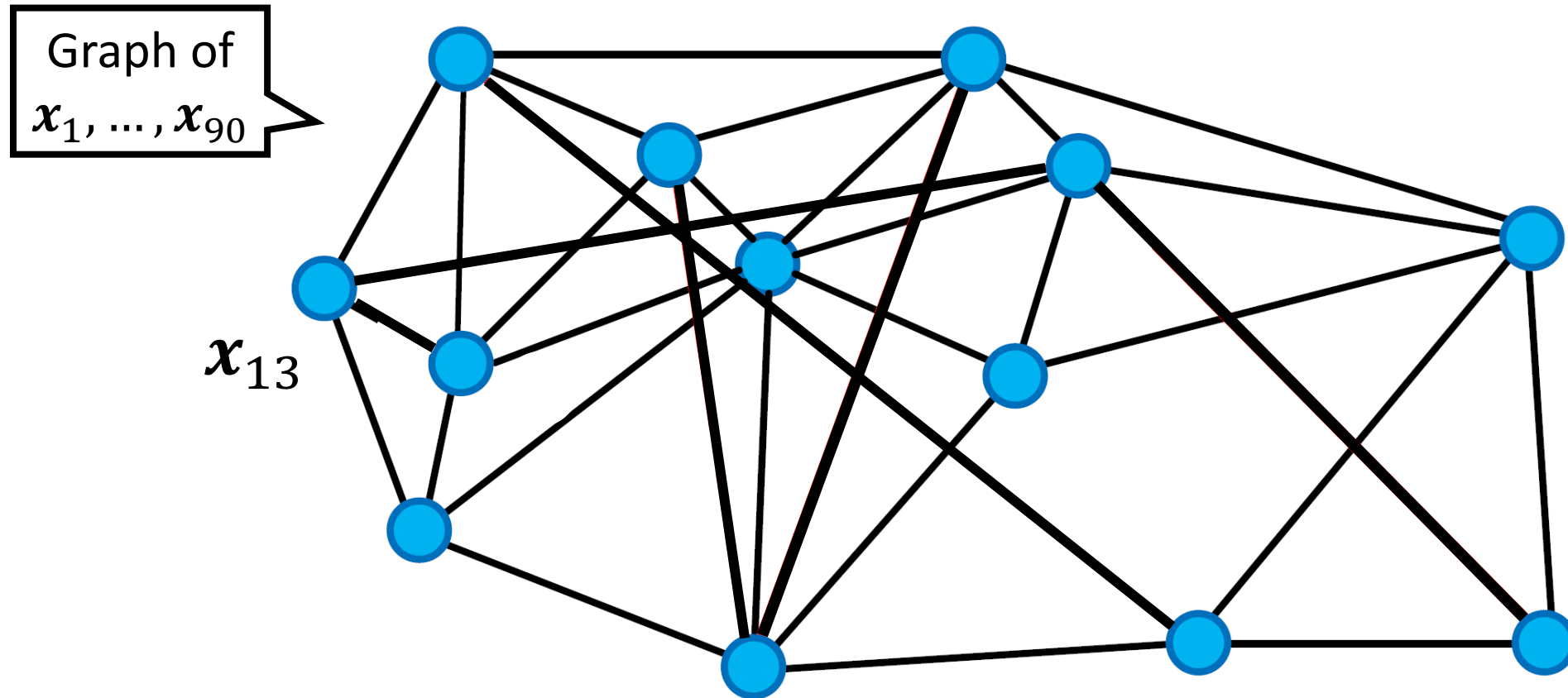


Refinement approach

- Iteratively refine an initial graph

Construction: incremental approach

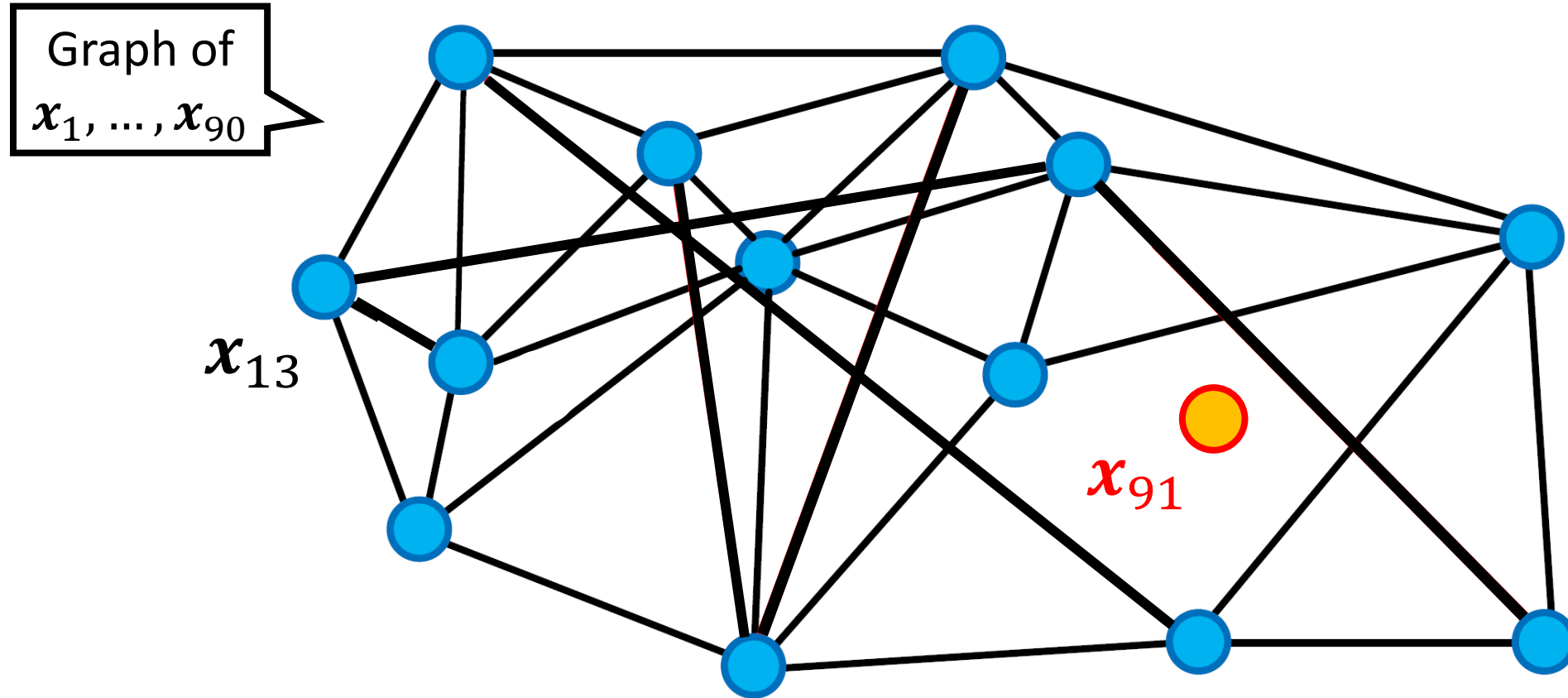
Images are from [Malkov+, Information Systems, 2013]



➤ Each node is a database vector

Construction: incremental approach

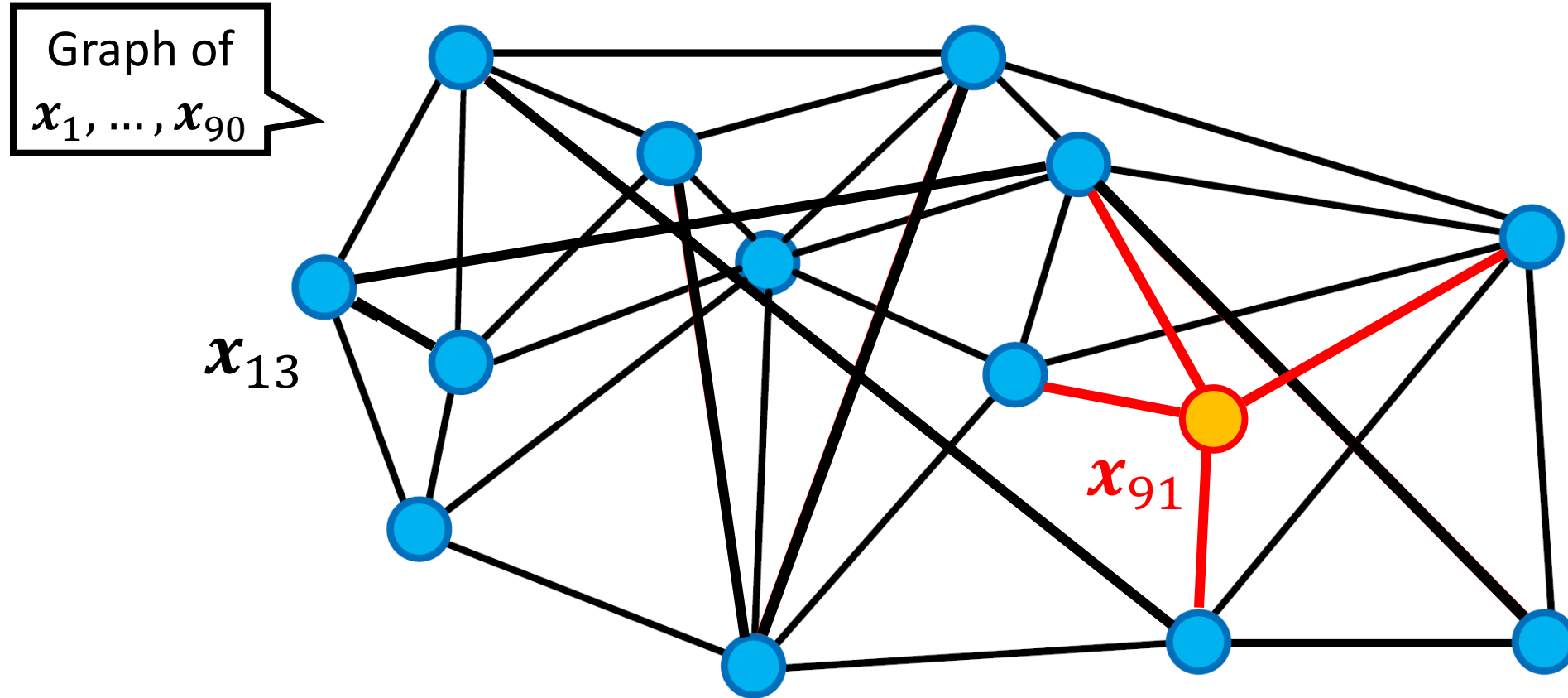
Images are from [Malkov+, Information Systems, 2013]



- Each node is a database vector
- Given a new database vector,

Construction: incremental approach

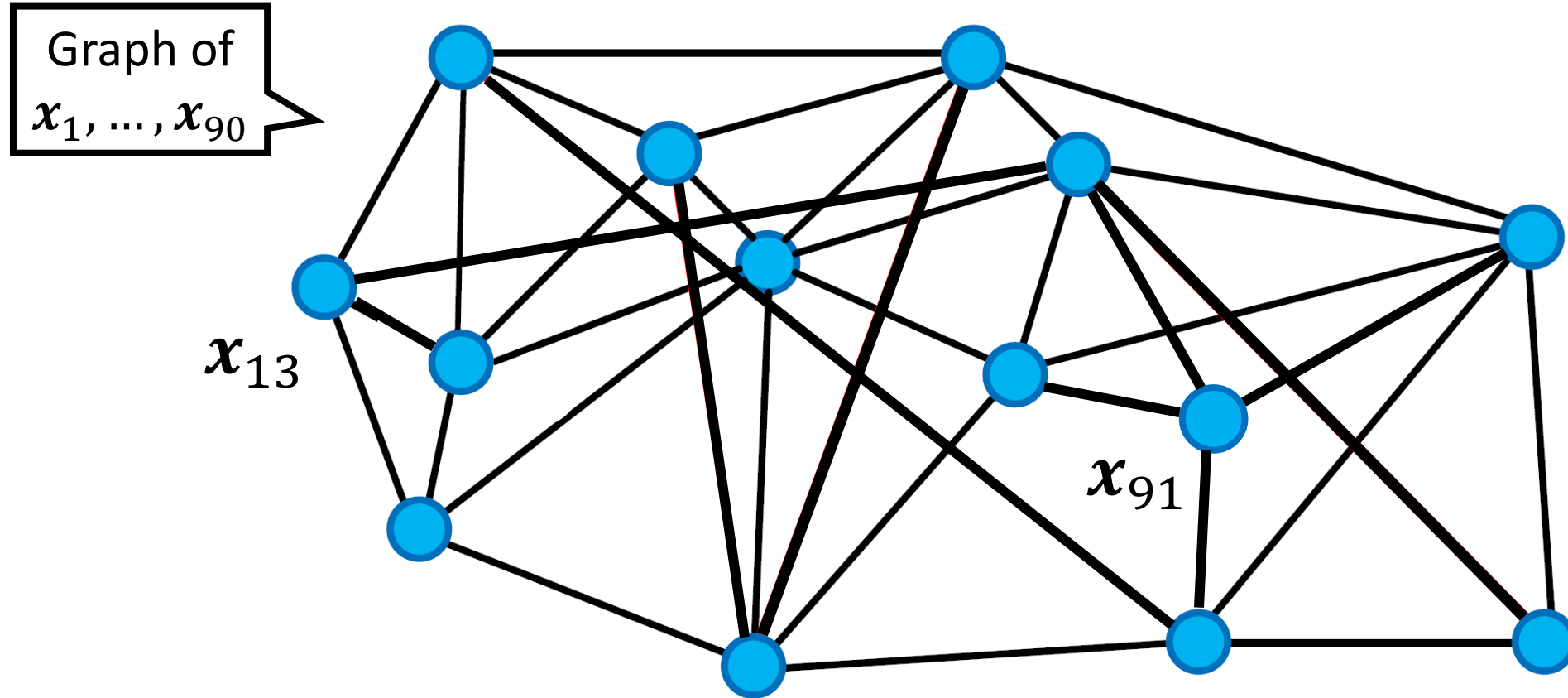
Images are from [Malkov+, Information Systems, 2013]



- Each node is a database vector
- Given a new database vector, create new edges to neighbors

Construction: incremental approach

Images are from [Malkov+, Information Systems, 2013]



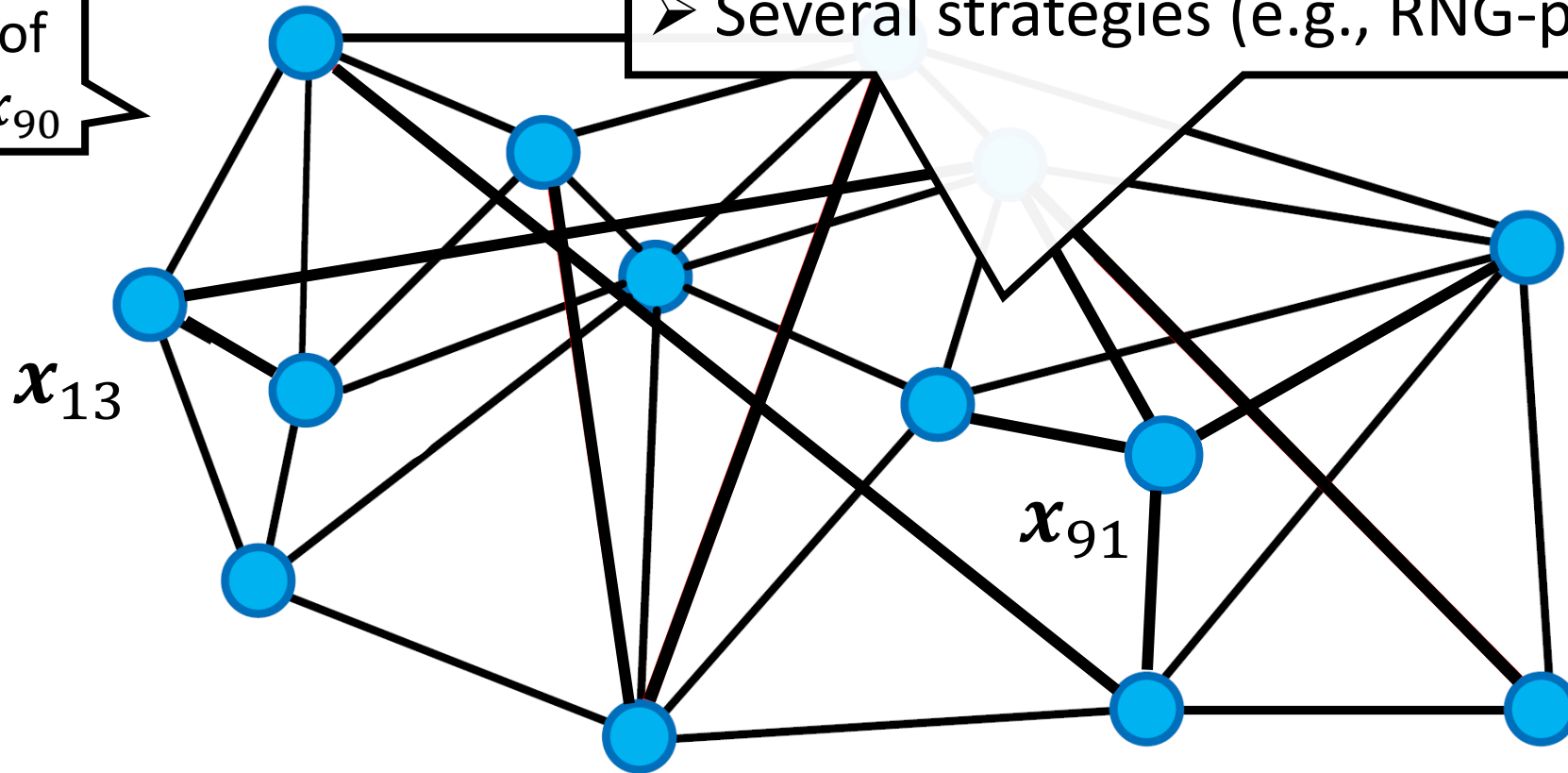
- Each node is a database vector
- Given a new database vector, create new edges to neighbors

Construction: incremental approach

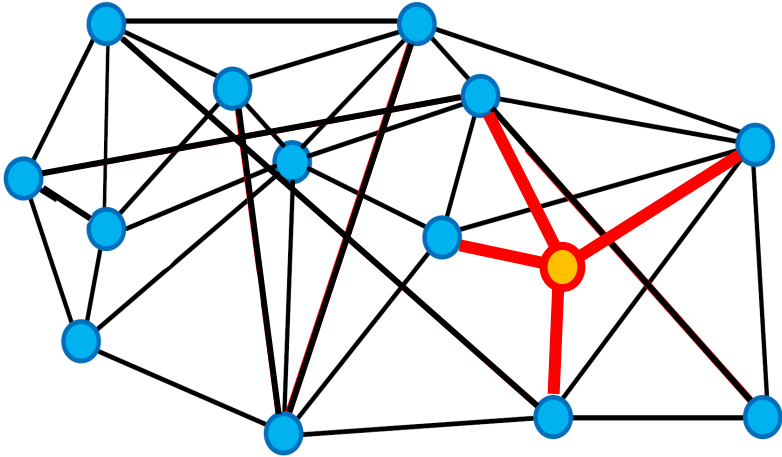
Images are from [Mallory, Information Systems, 2012]

- Prune edges if some node have too many edges
- Several strategies (e.g., RNG-pruning)

Graph of
 $x_1, \dots, x_{90}$

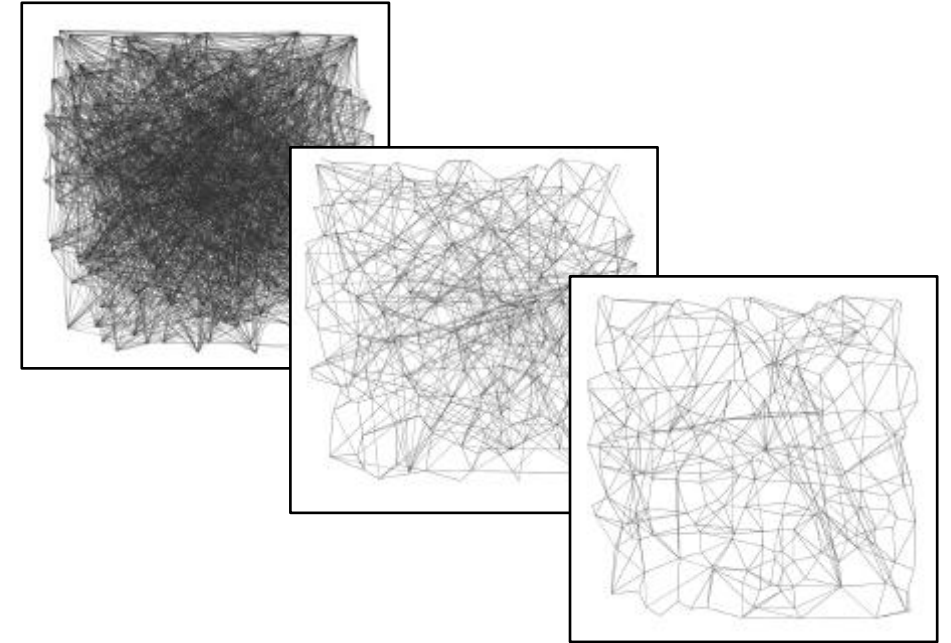


- Each node is a database vector
- Given a new database vector, create new edges to neighbors



Increment approach

- Add a new item to the current graph incrementally

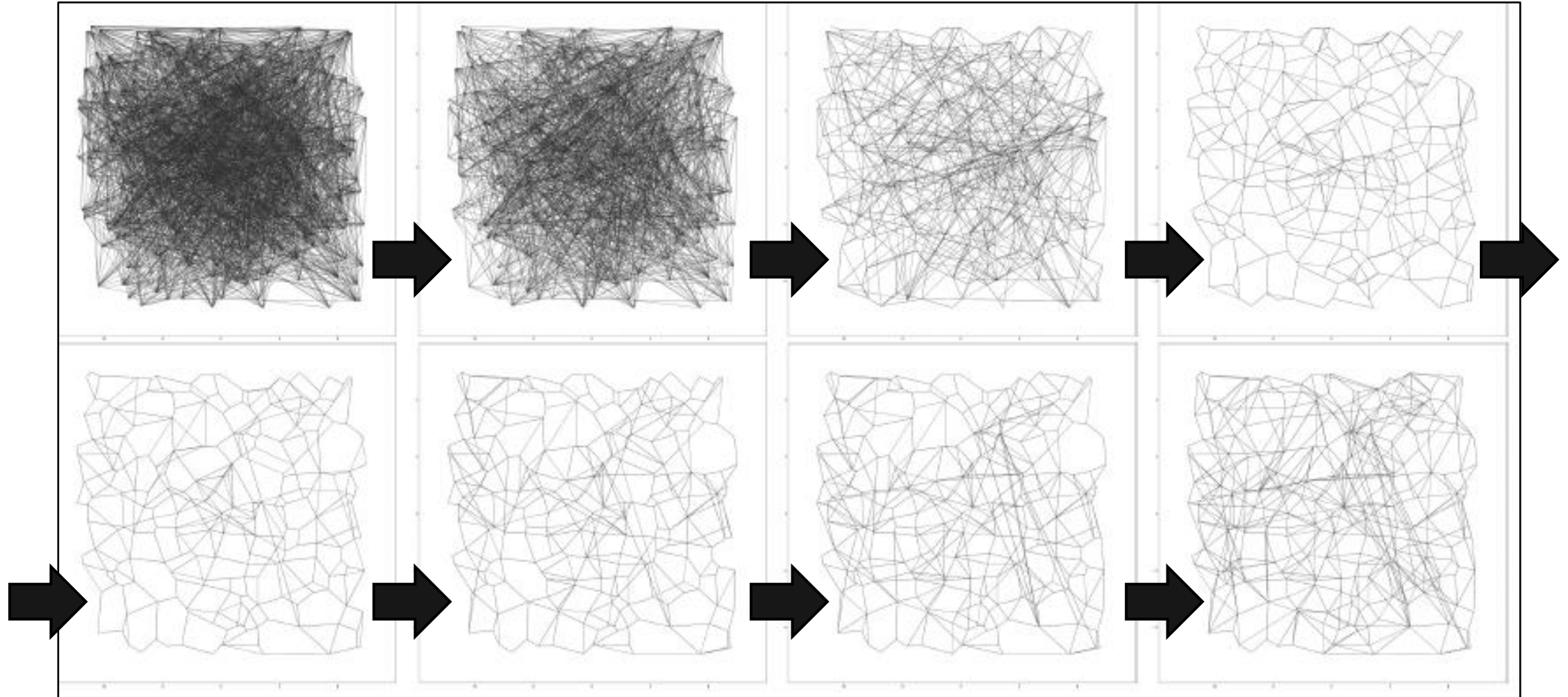


Refinement approach

- Iteratively refine an initial graph

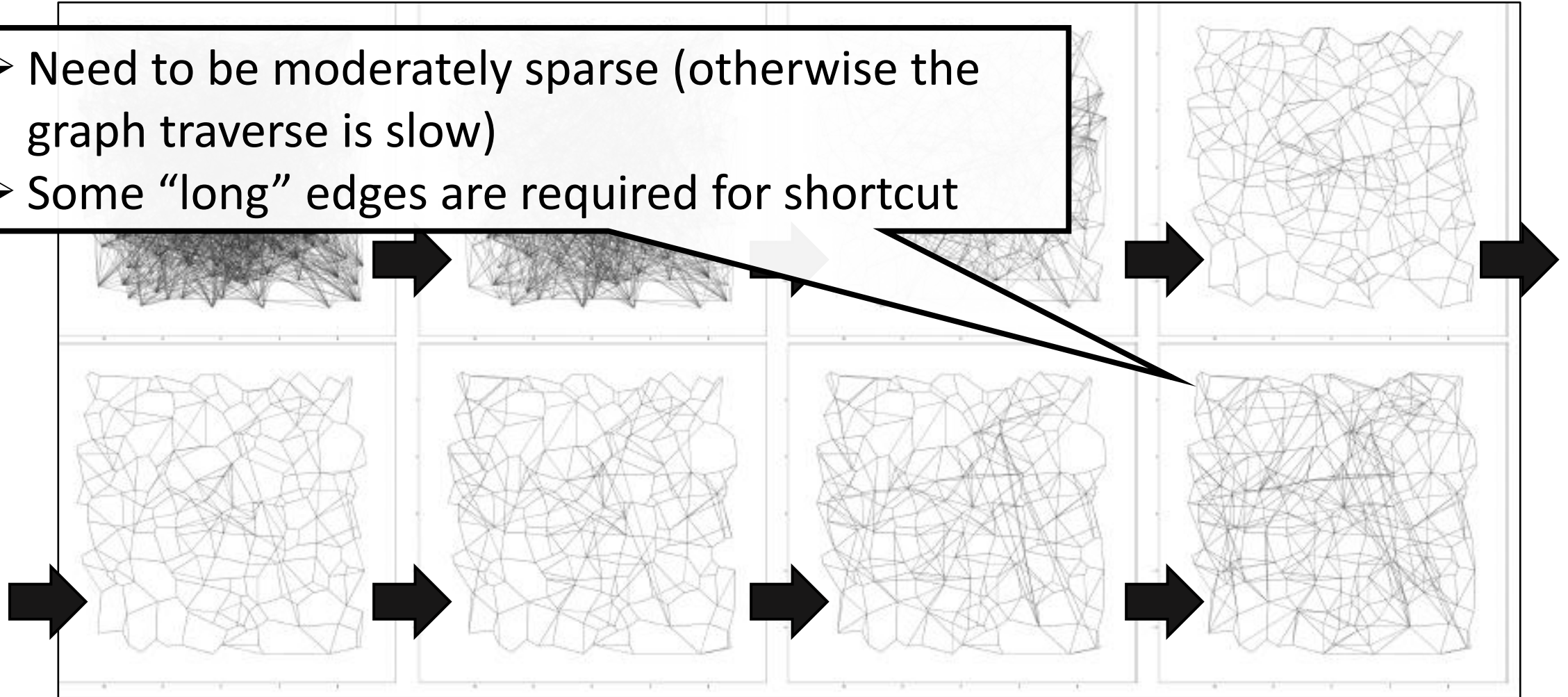
Construction: refinement approach

Images are from [Subramanya+, NerulPS 2019]



- Create an initial graph (e.g., random graph or approx. kNN graph)
- Refine it iteratively (pruning/adding edges)

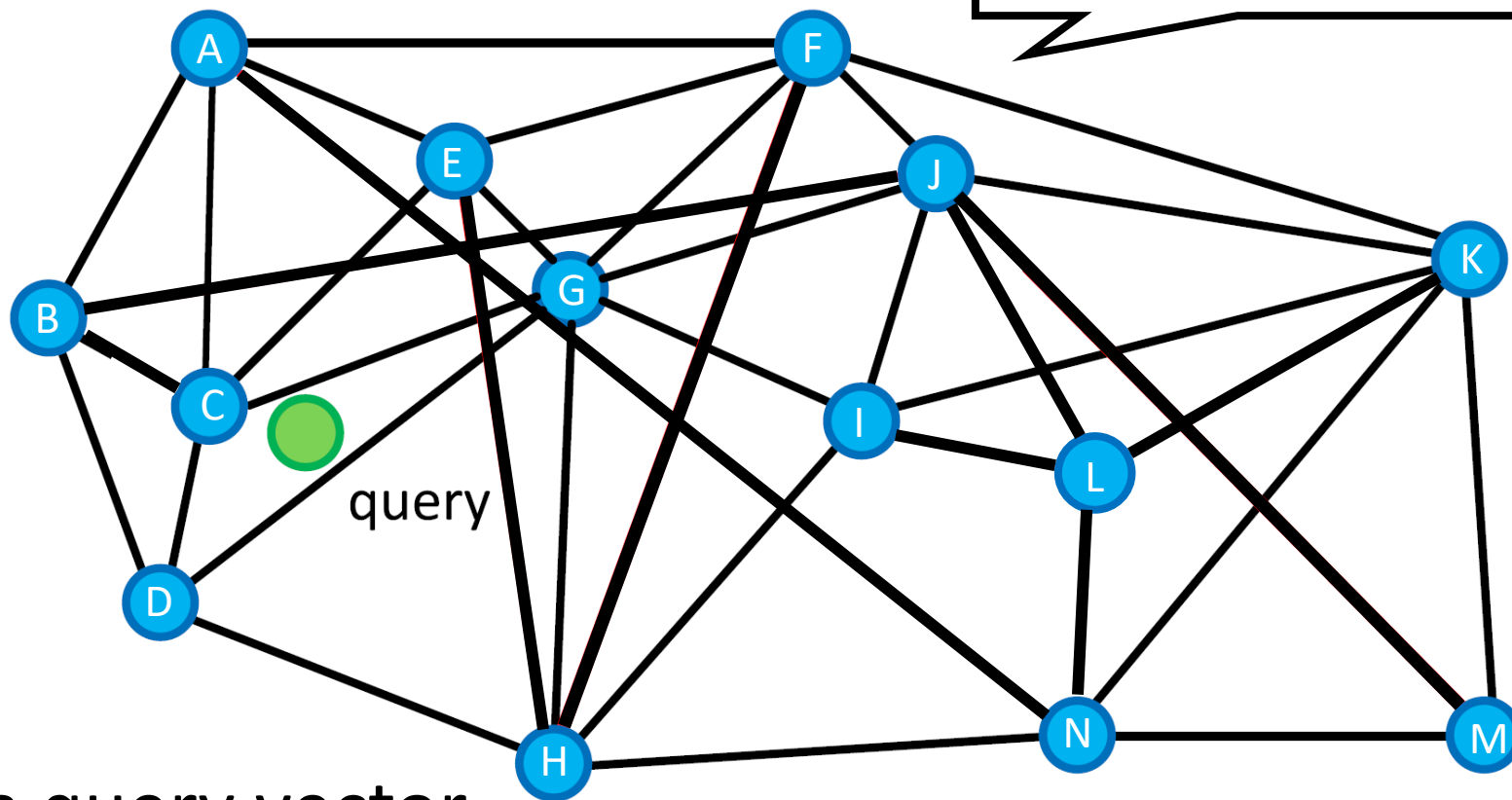
- Need to be moderately sparse (otherwise the graph traverse is slow)
- Some “long” edges are required for shortcut



- Create an initial graph (e.g., random graph or approx. kNN graph)
- Refine it iteratively (pruning/adding edges)

Search

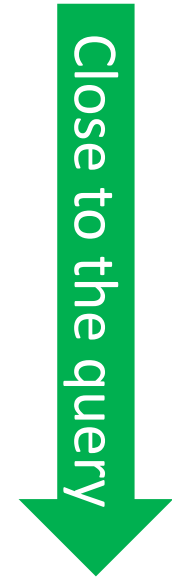
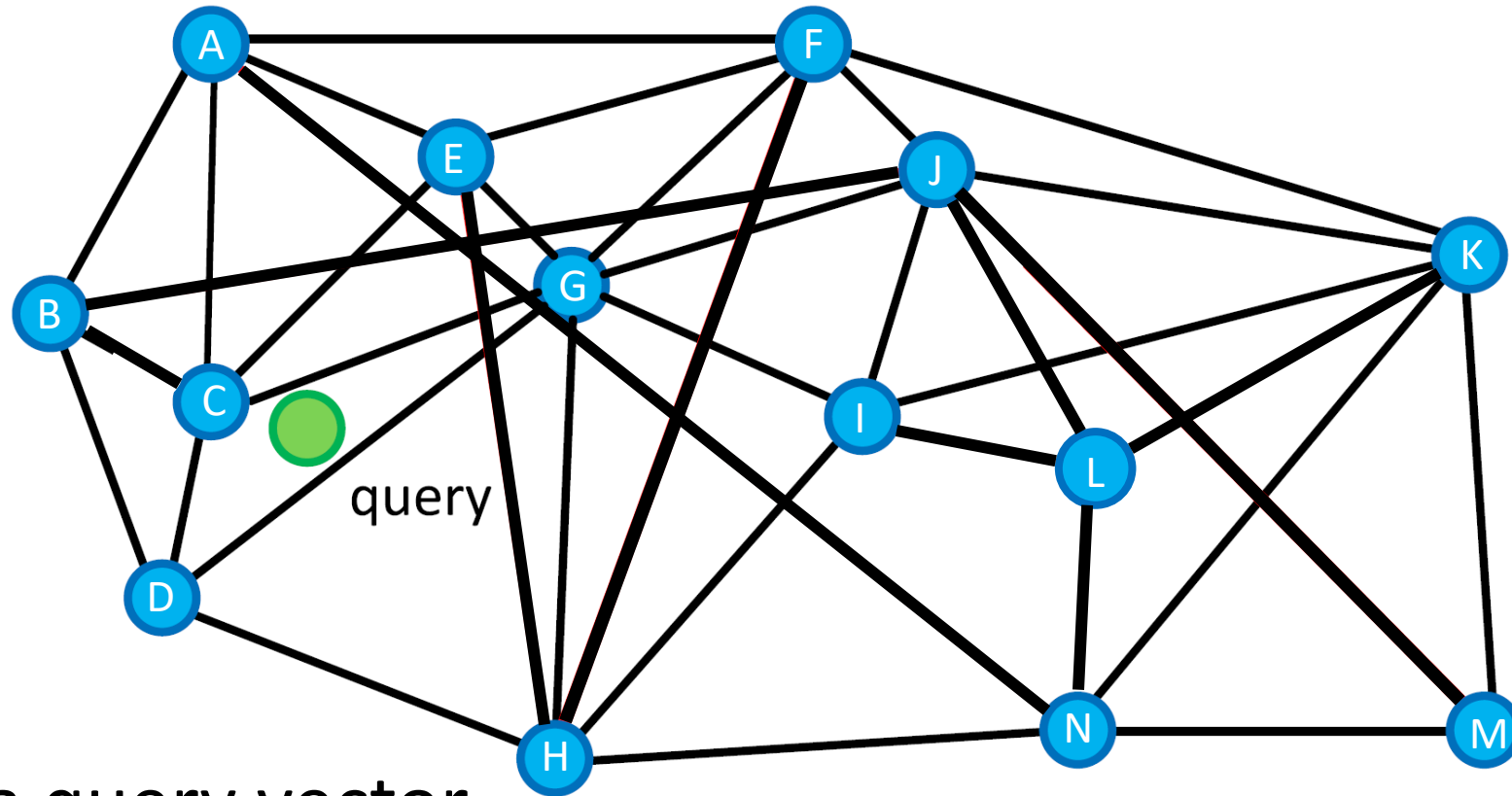
Images are from [Malkov+, Information Systems, 2013]



Close to the query

Candidates
(size = 3)

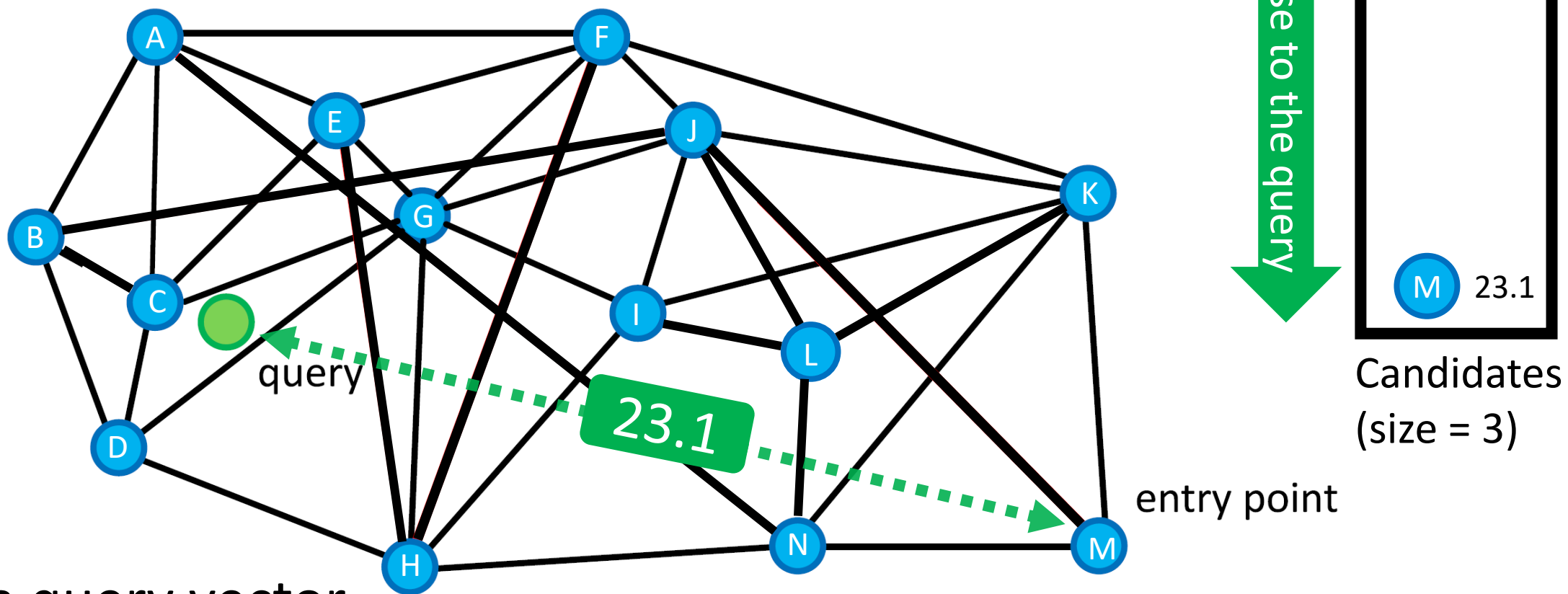
➤ Given a query vector



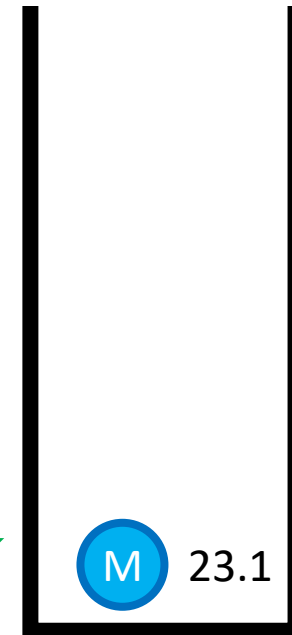
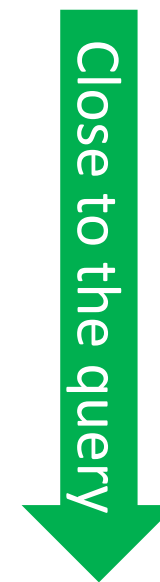
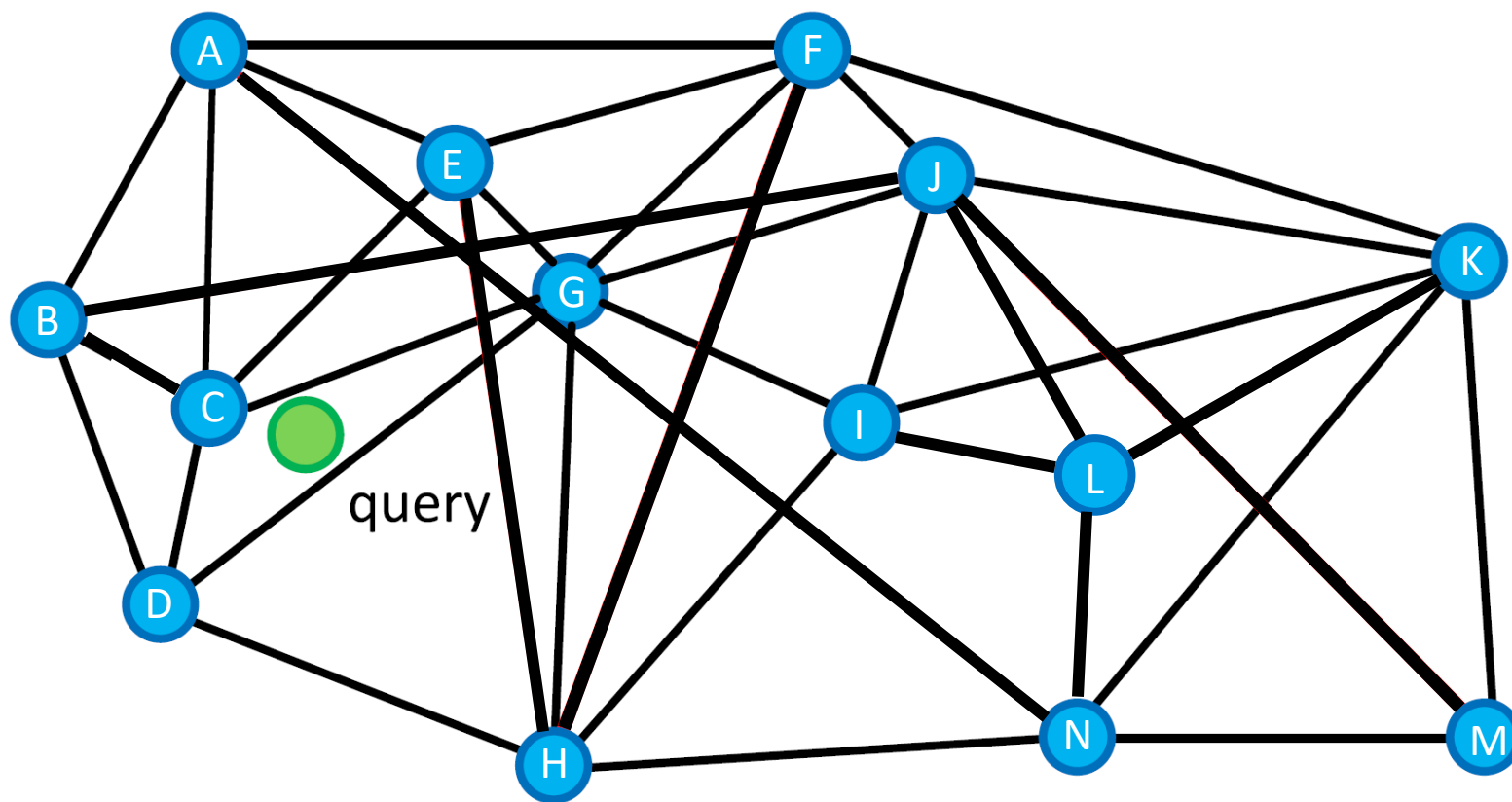
Candidates
(size = 3)

entry point

- Given a query vector
- Start from an entry point (e.g., M)

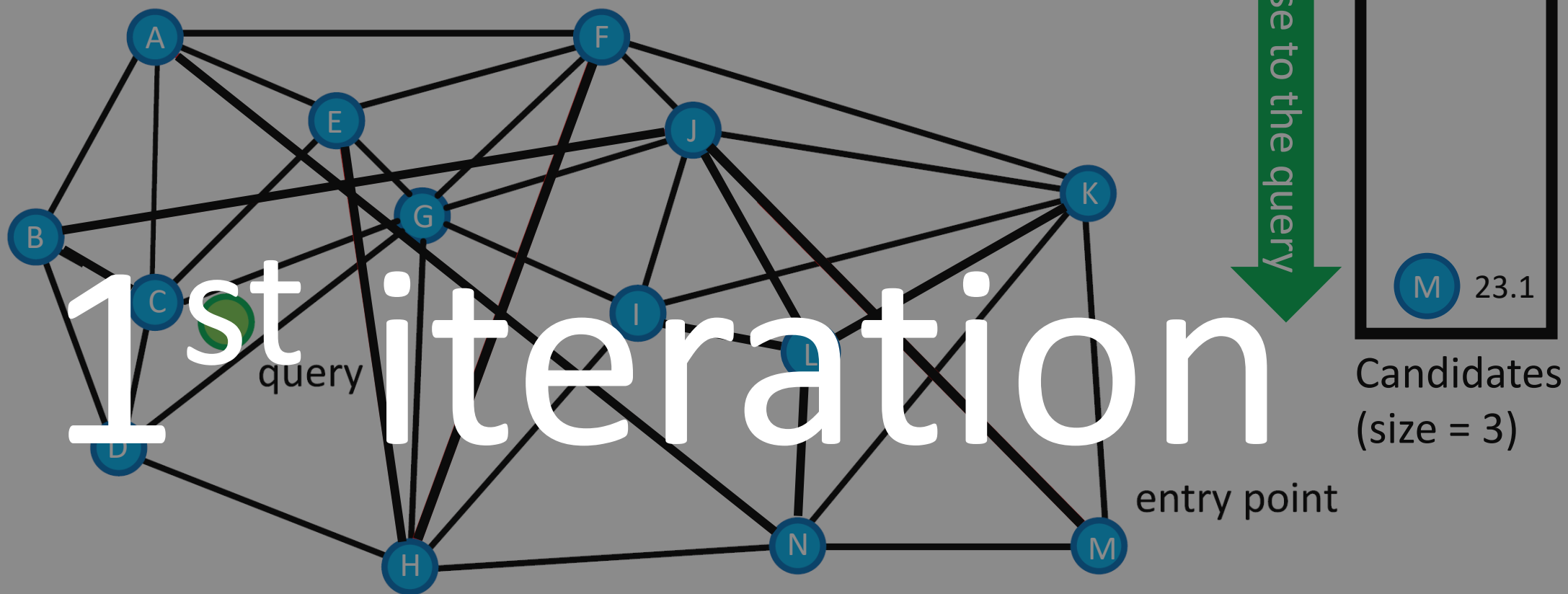


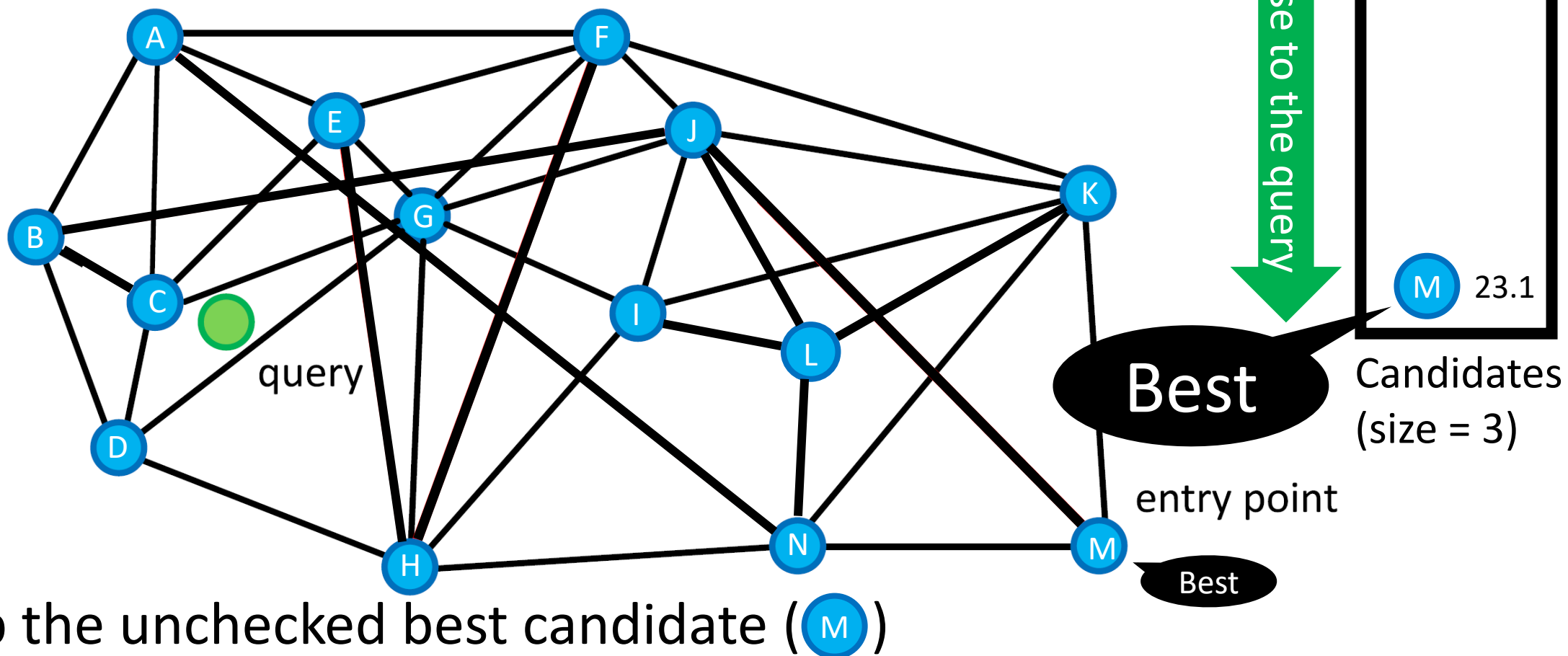
- Given a query vector
- Start from an entry point (e.g., **M**). Record the distance to **q**.

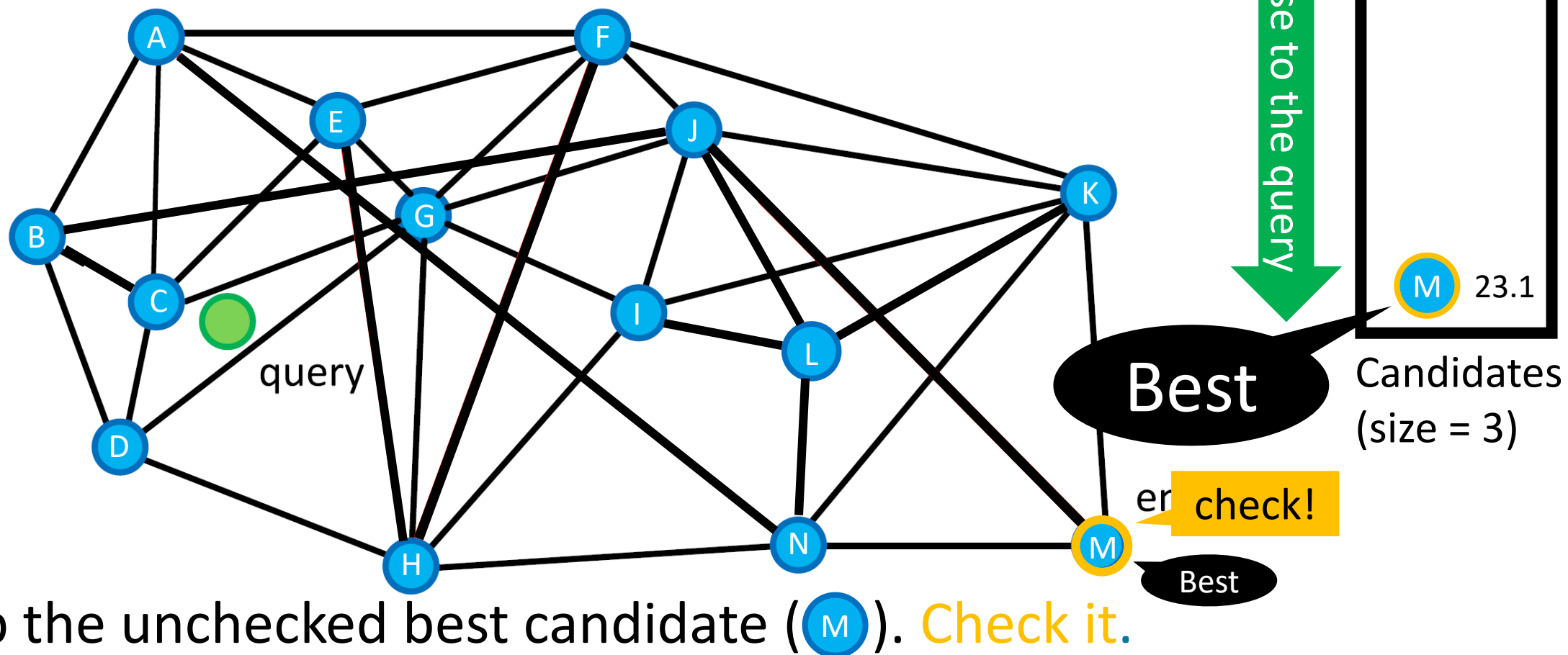


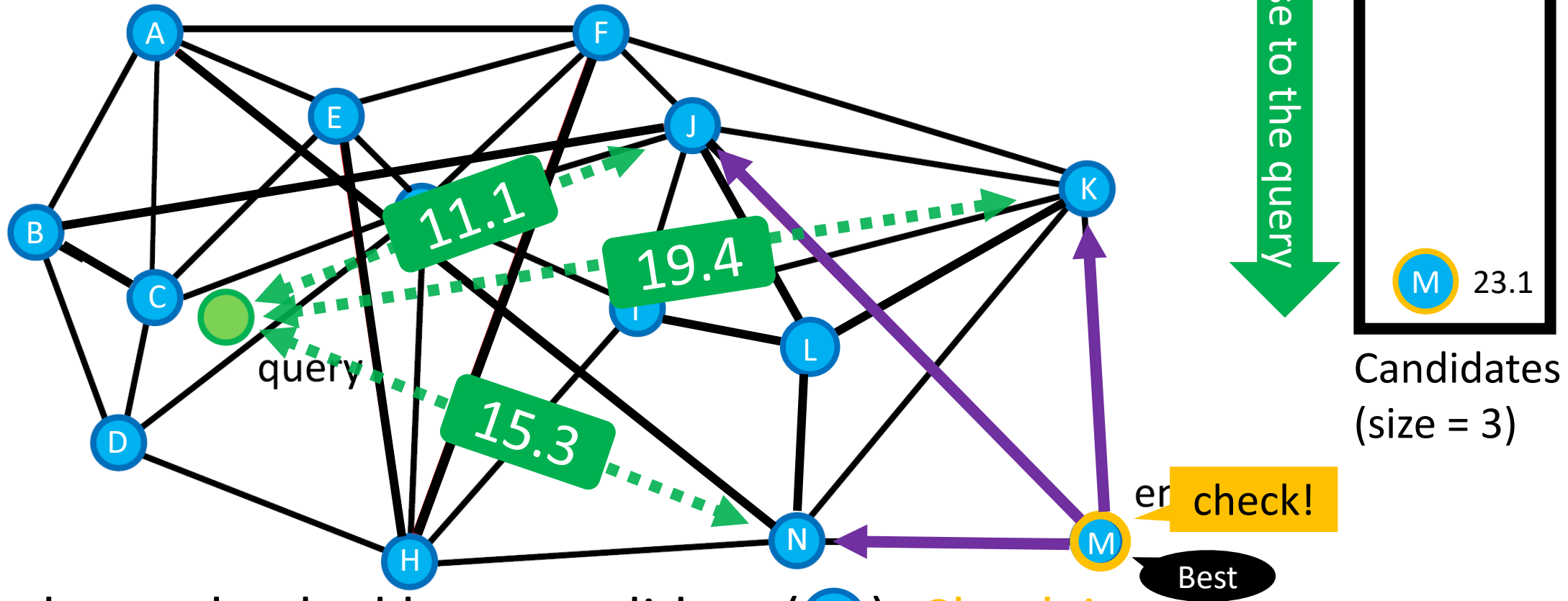
Candidates
(size = 3)

entry point

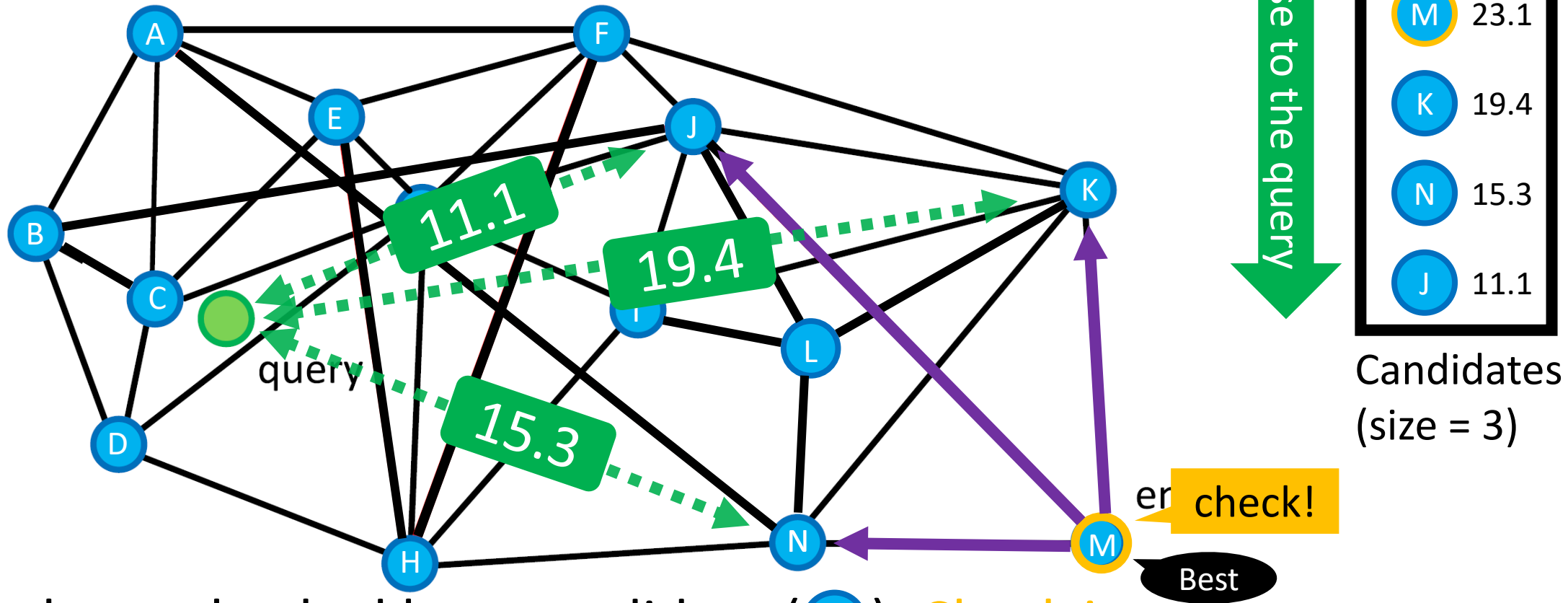




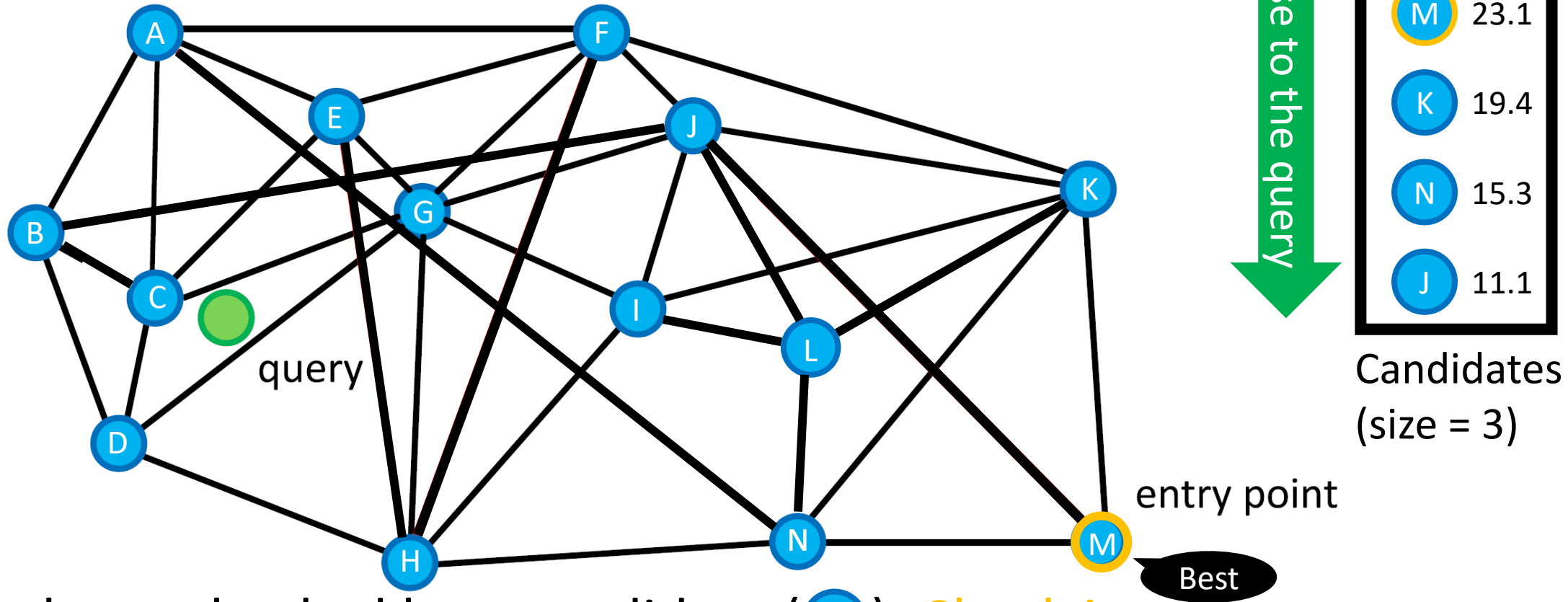




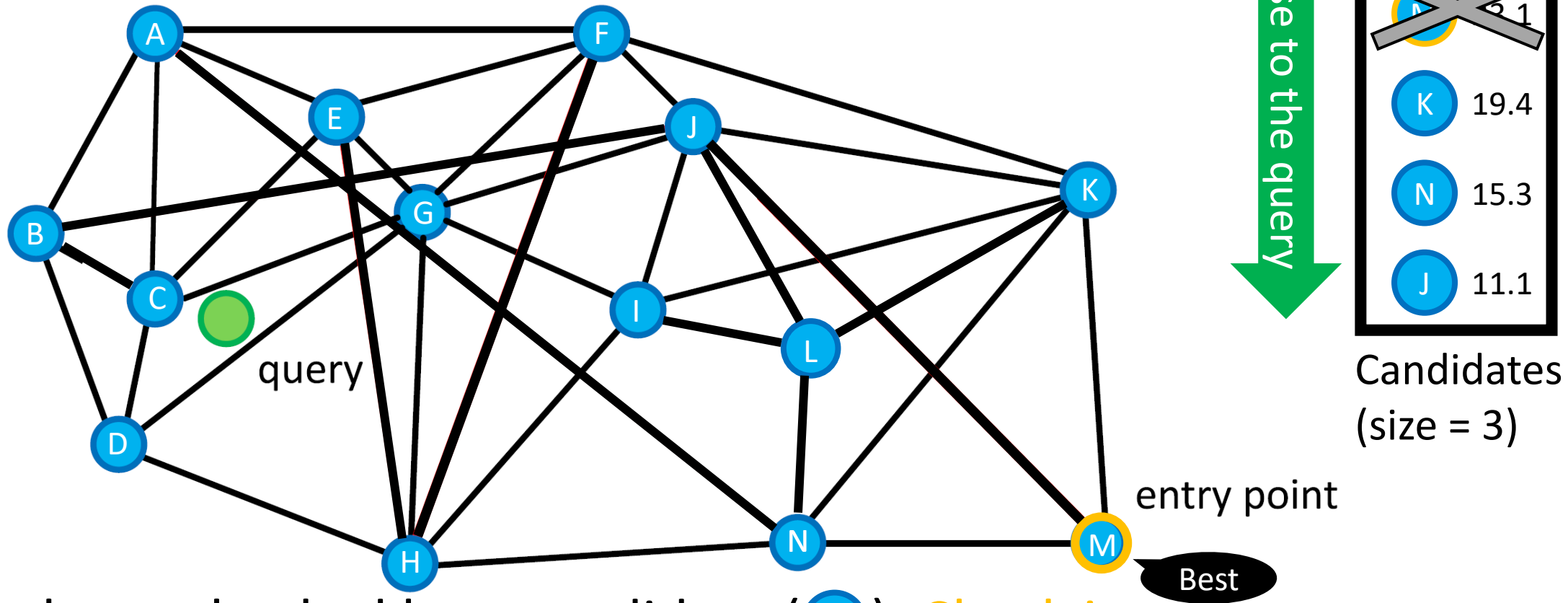
- Pick up the unchecked best candidate (M). Check it.
- Find the connected points.
- Record the distances to q.



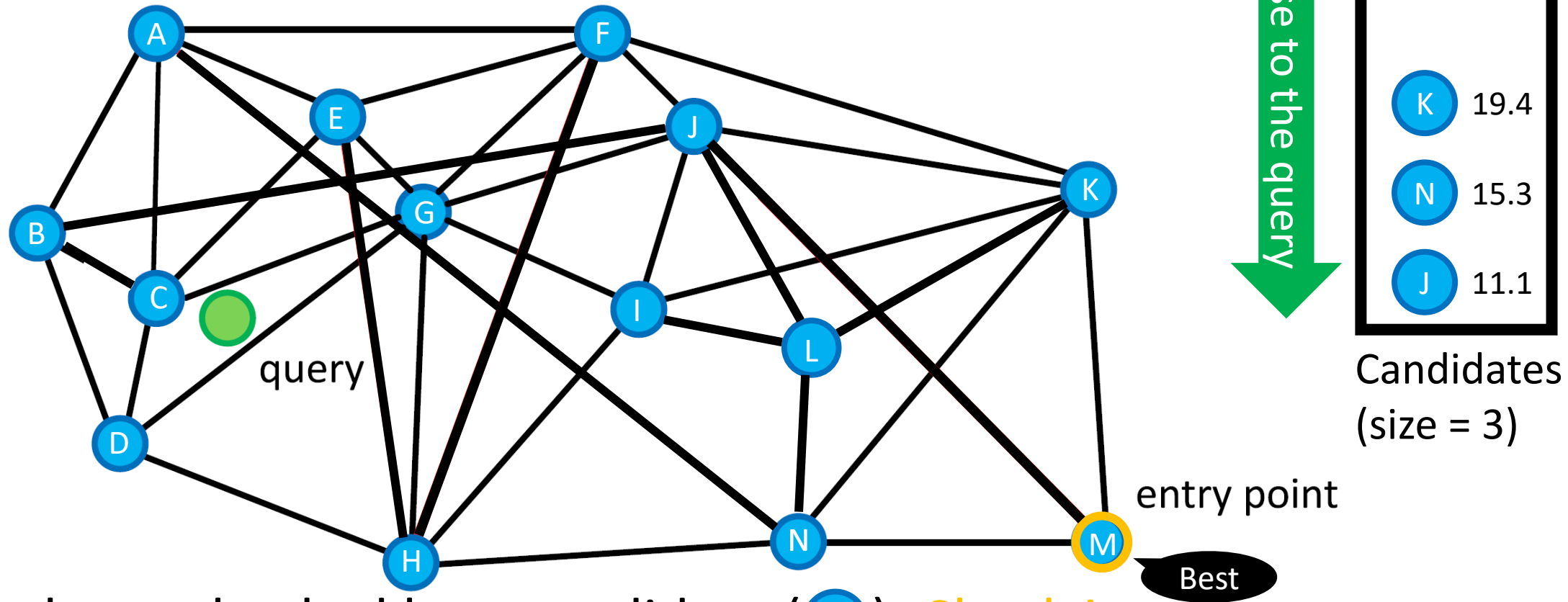
- Pick up the unchecked best candidate (M). Check it.
- Find the connected points.
- Record the distances to q.



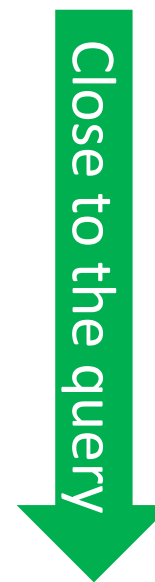
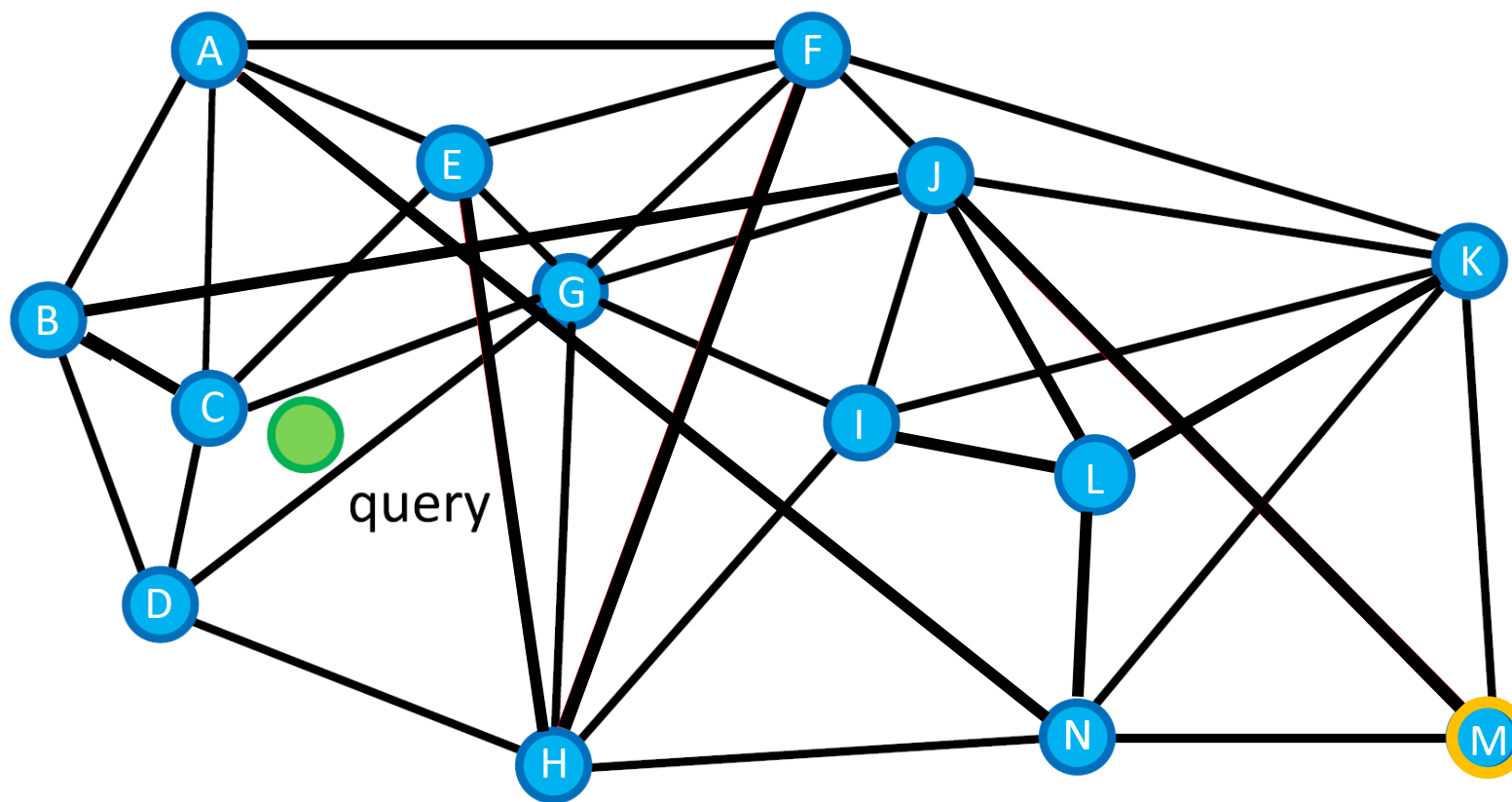
- Pick up the unchecked best candidate (M). Check it.
- Find the connected points.
- Record the distances to q.



- Pick up the unchecked best candidate (M). Check it. .
- Find the connected points.
- Record the distances to q.
- Maintain the candidates (size=3)



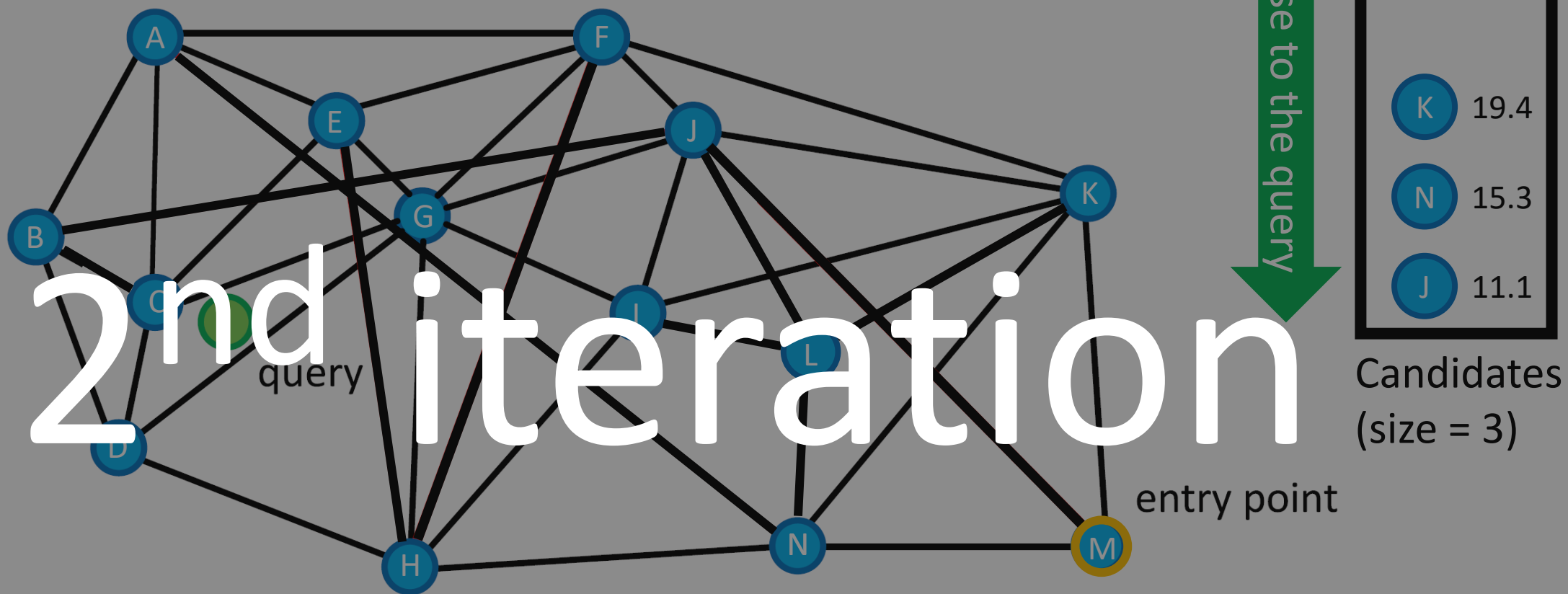
- Pick up the unchecked best candidate (M). Check it.
- Find the connected points.
- Record the distances to q.
- Maintain the candidates (size=3)

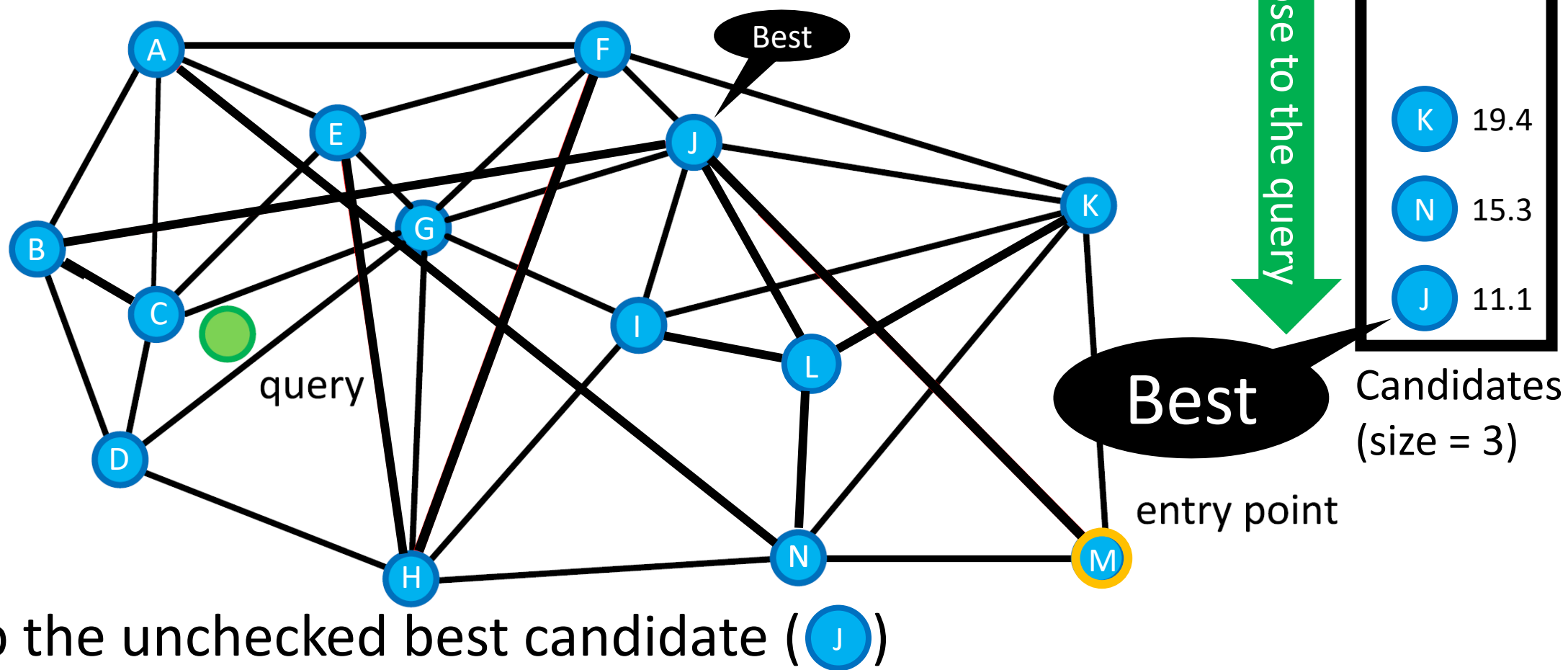


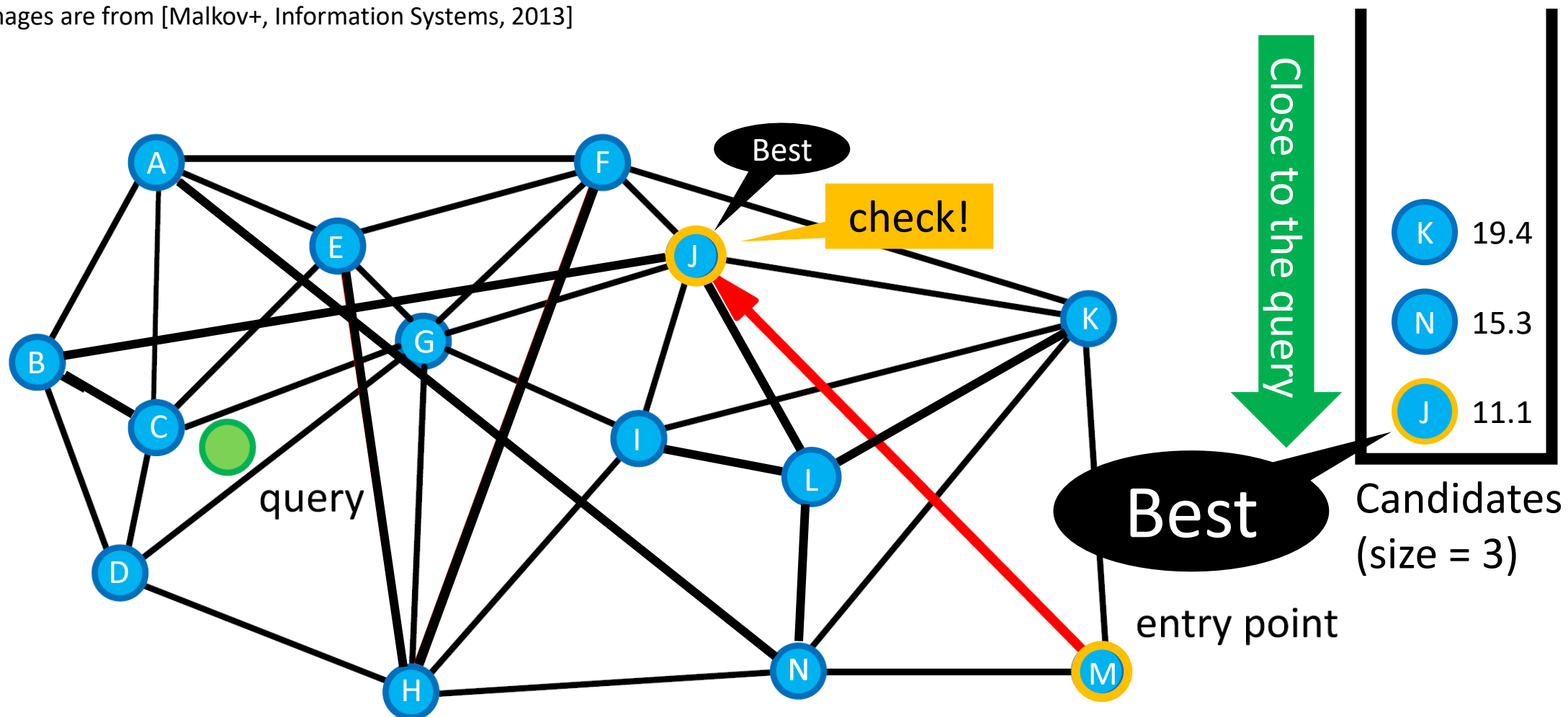
K	19.4
N	15.3
J	11.1

Candidates
(size = 3)

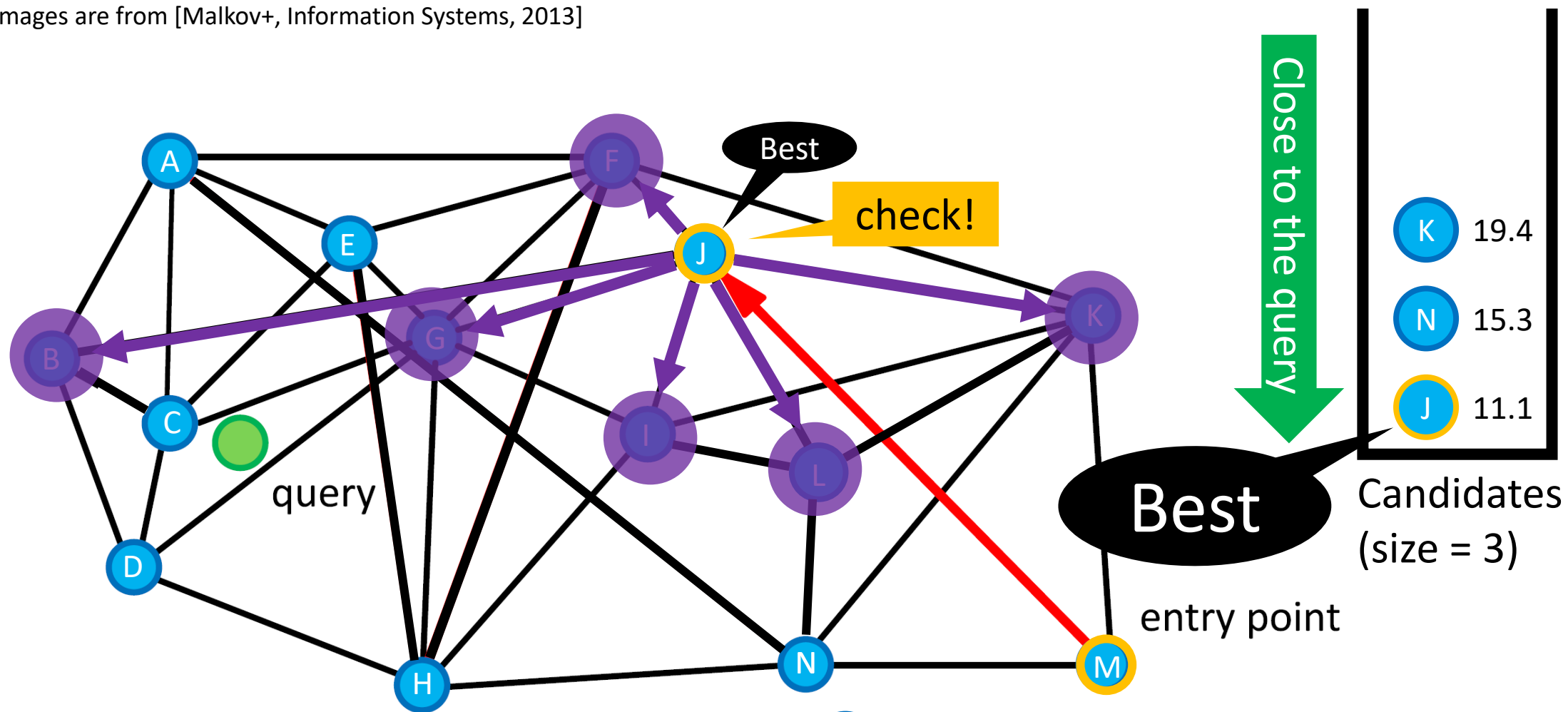
entry point



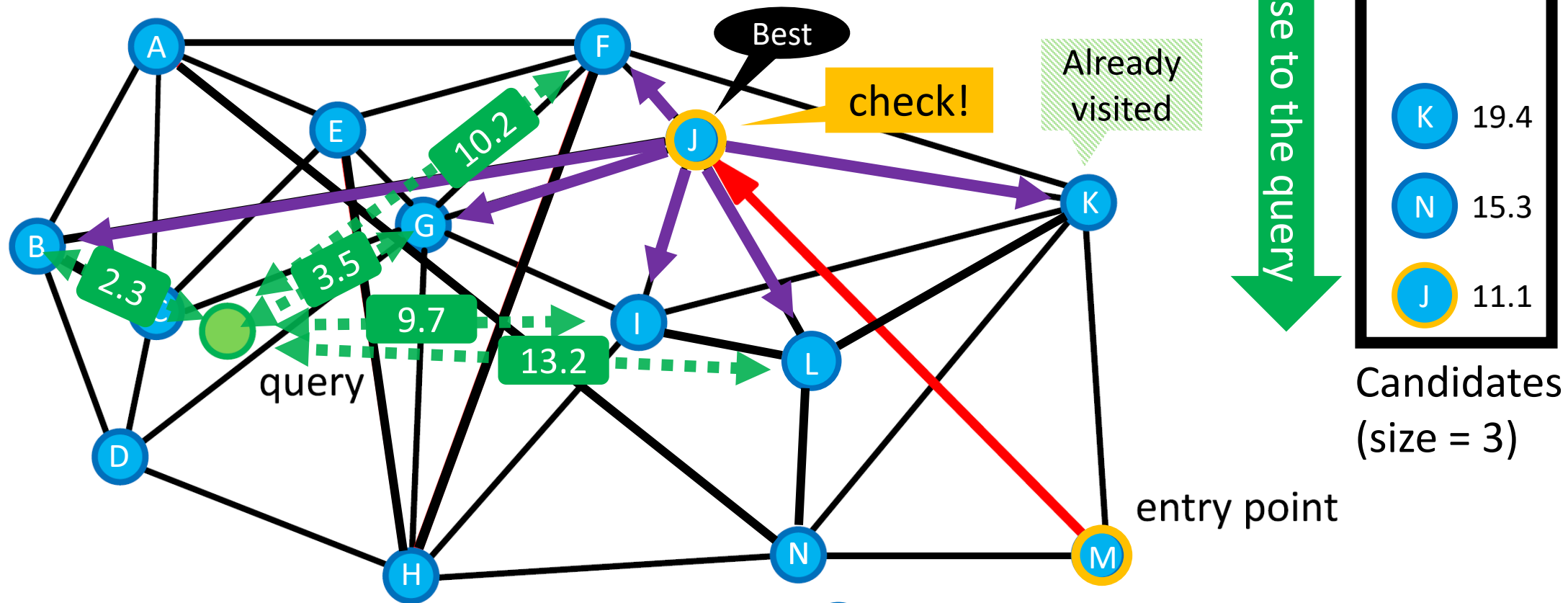




➤ Pick up the unchecked best candidate (J). Check it.



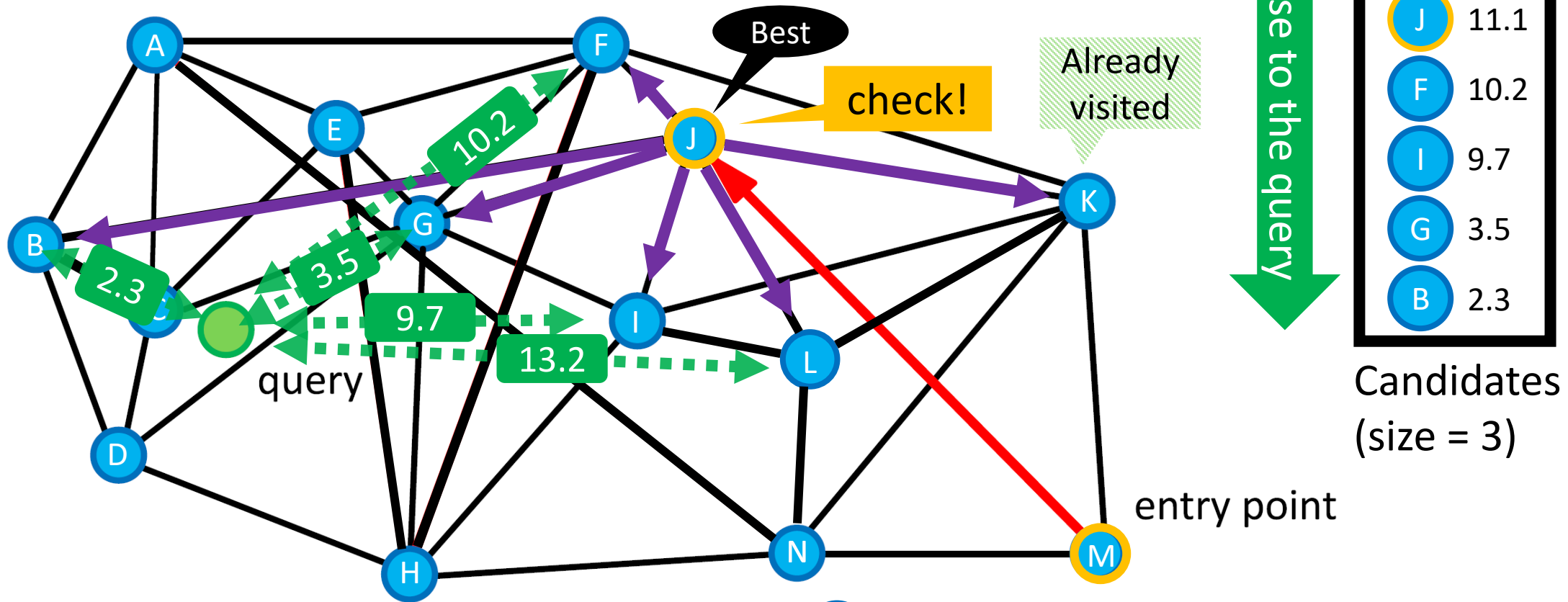
- Pick up the unchecked best candidate (J). Check it.
- Find the connected points.



- Pick up the unchecked best candidate (J). Check it.
- Find the connected points.
- Record the distances to q.

Search

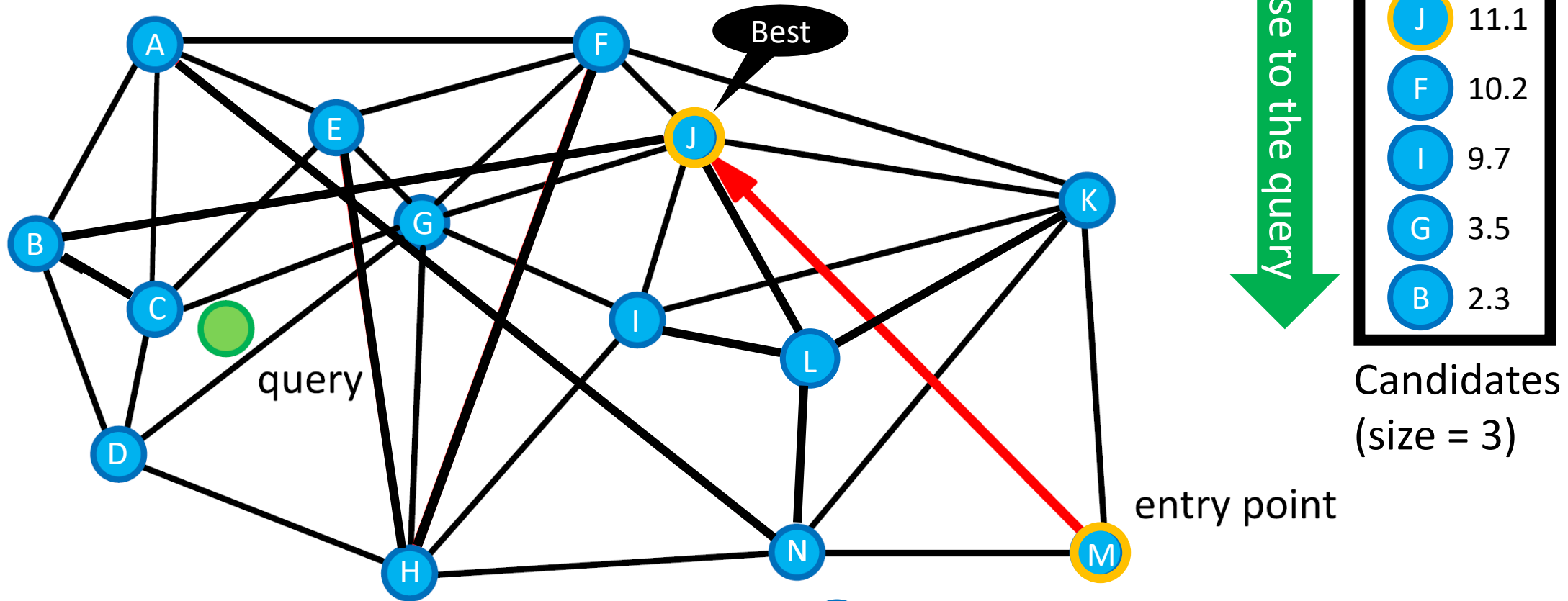
Images are from [Malkov+, Information Systems, 2013]



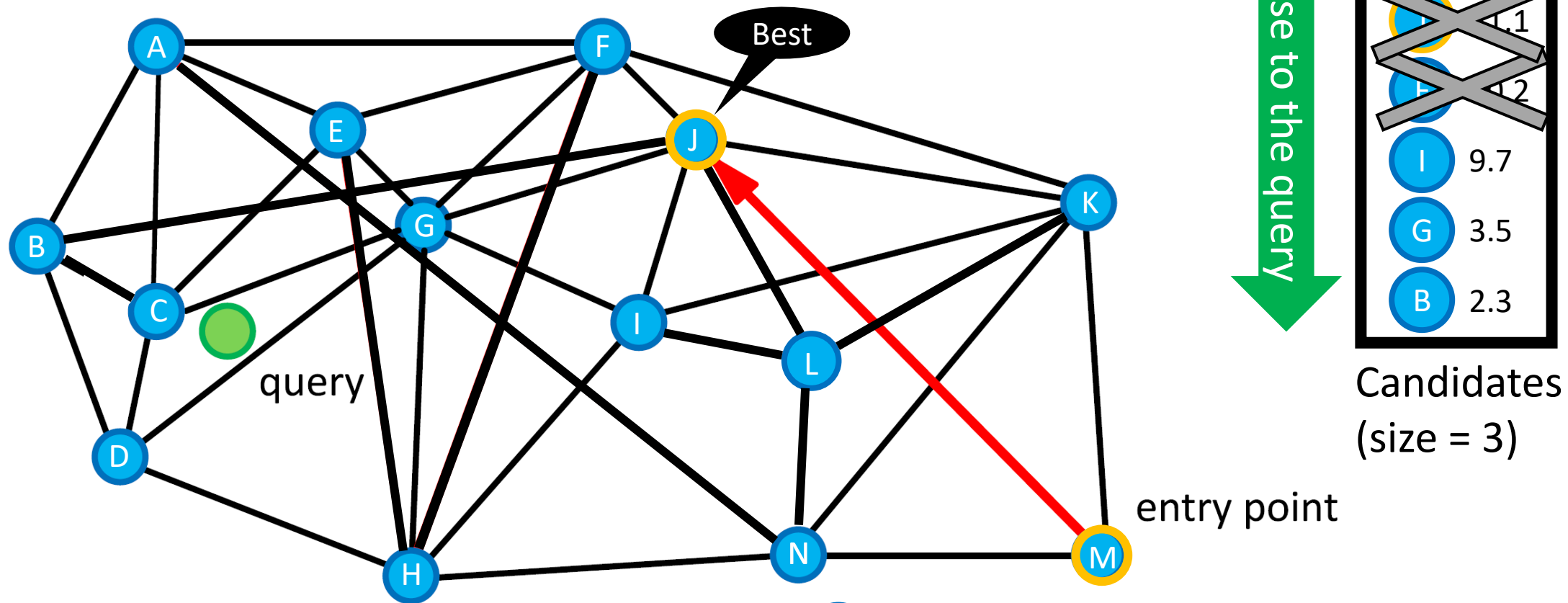
- Pick up the unchecked best candidate (J). Check it.
- Find the connected points.
- Record the distances to q.

Search

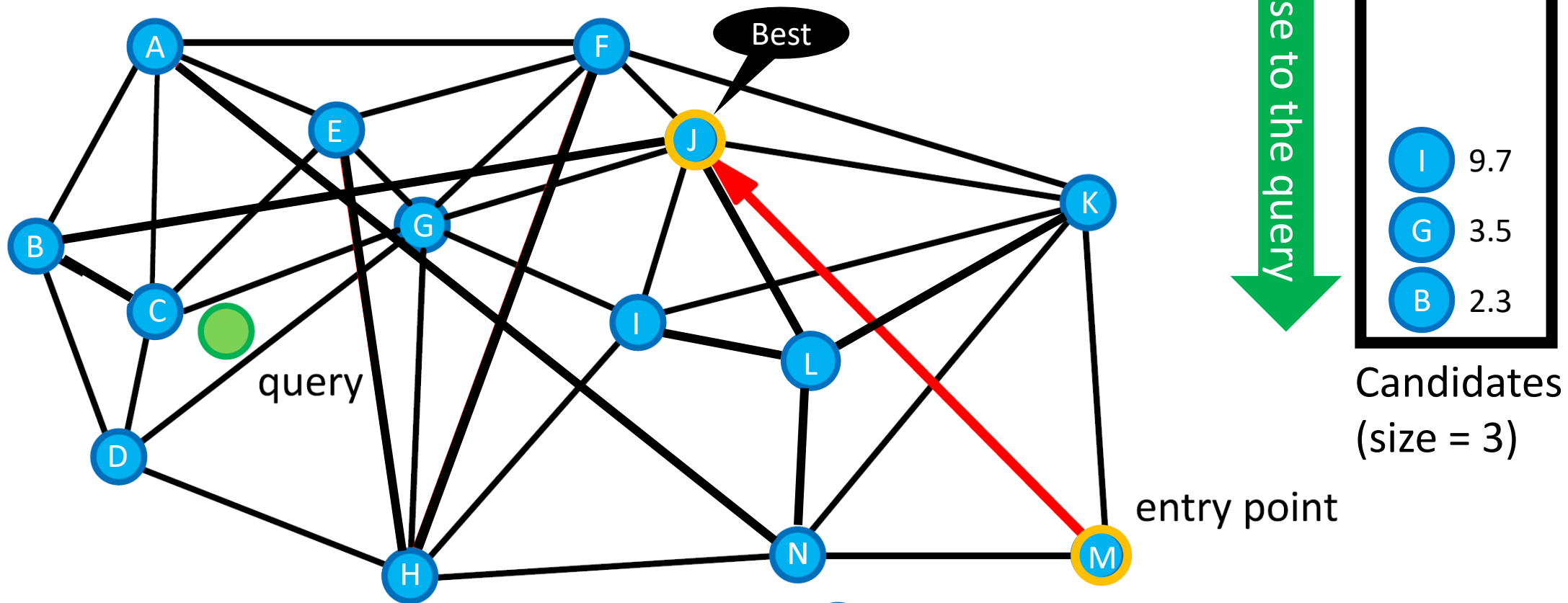
Images are from [Malkov+, Information Systems, 2013]



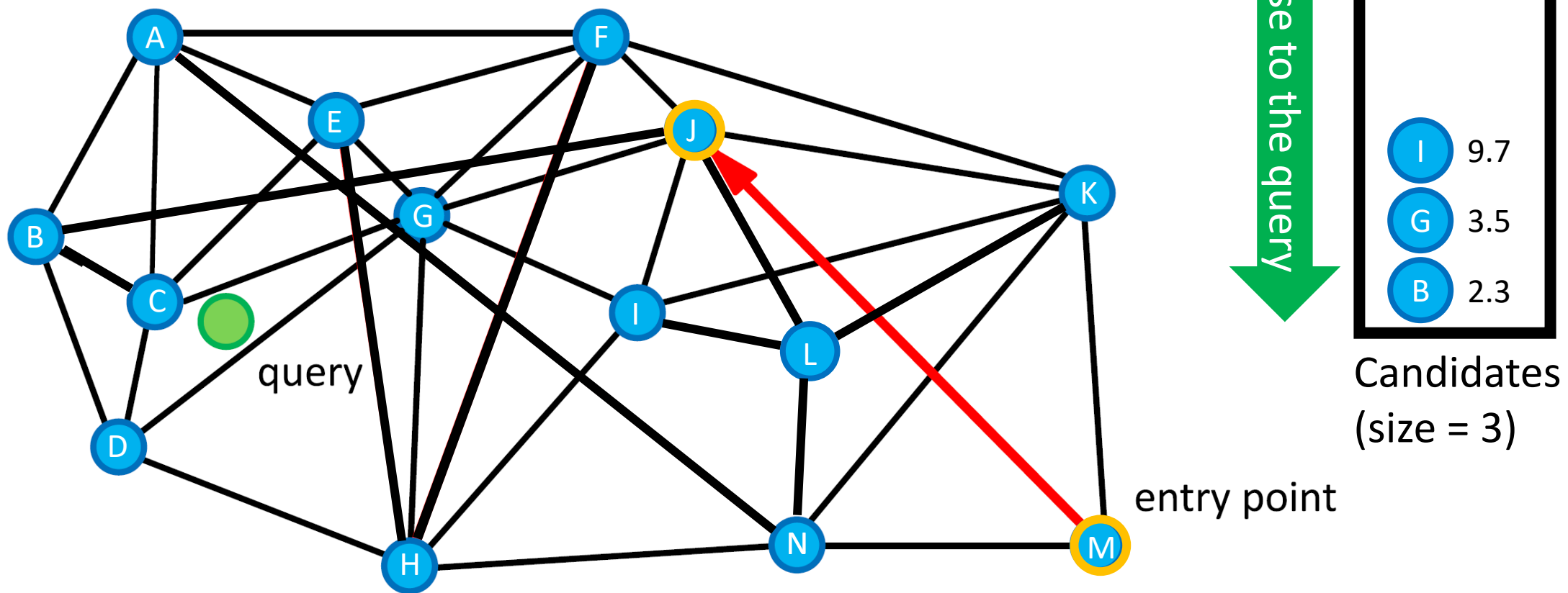
- Pick up the unchecked best candidate (J). Check it.
- Find the connected points.
- Record the distances to q.

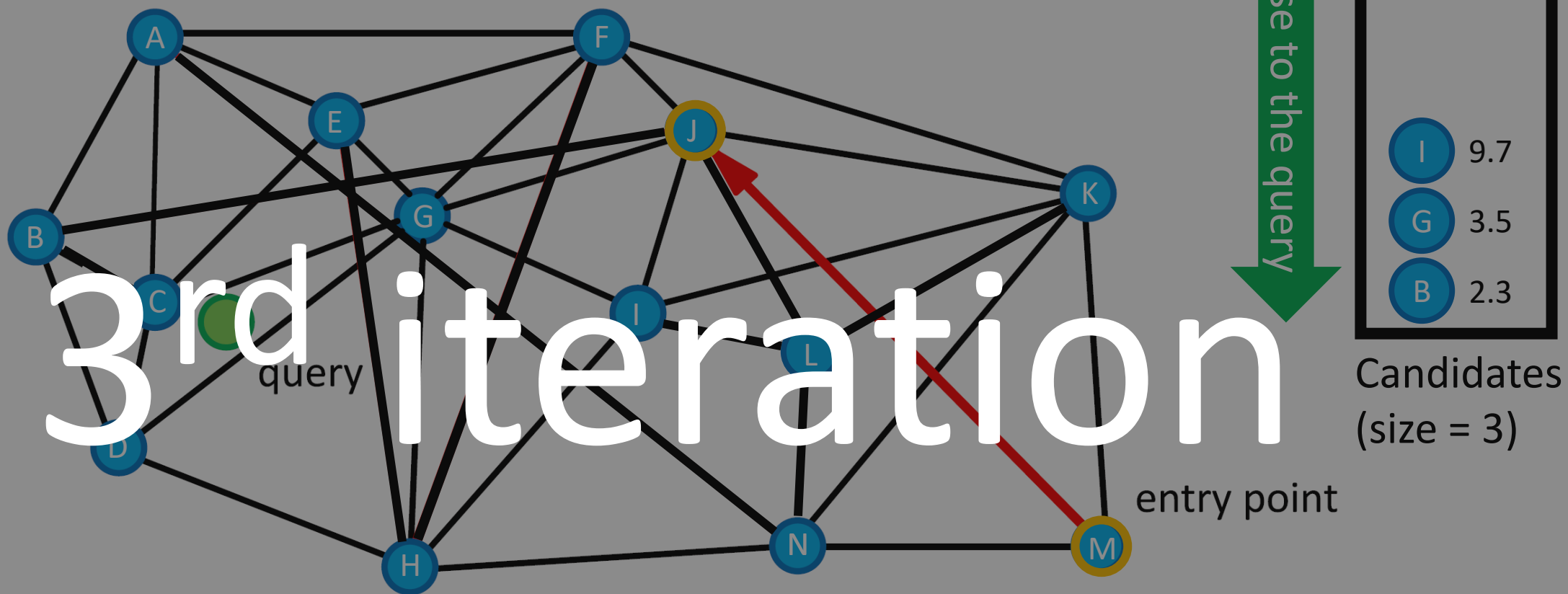


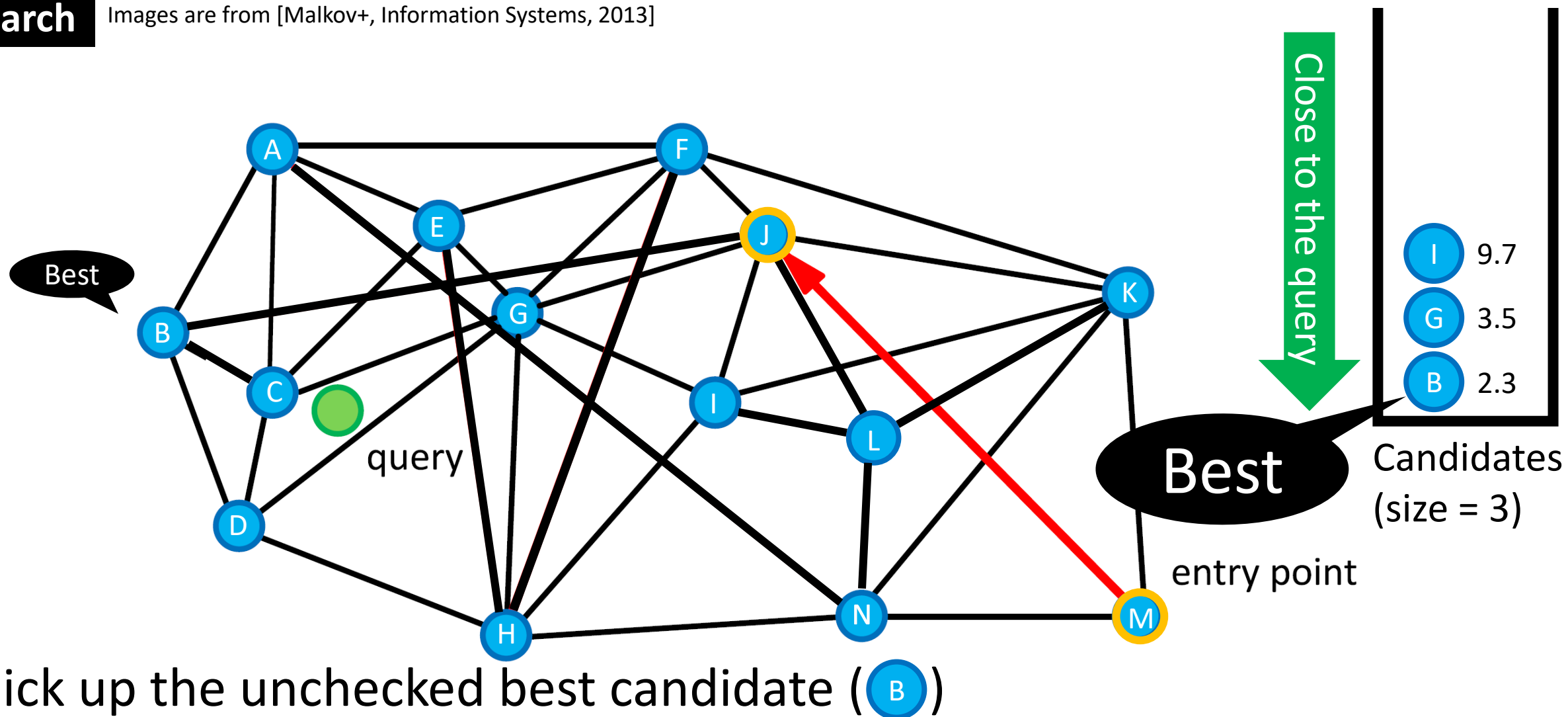
- Pick up the unchecked best candidate (J). Check it.
- Find the connected points.
- Record the distances to q.
- Maintain the candidates (size=3)

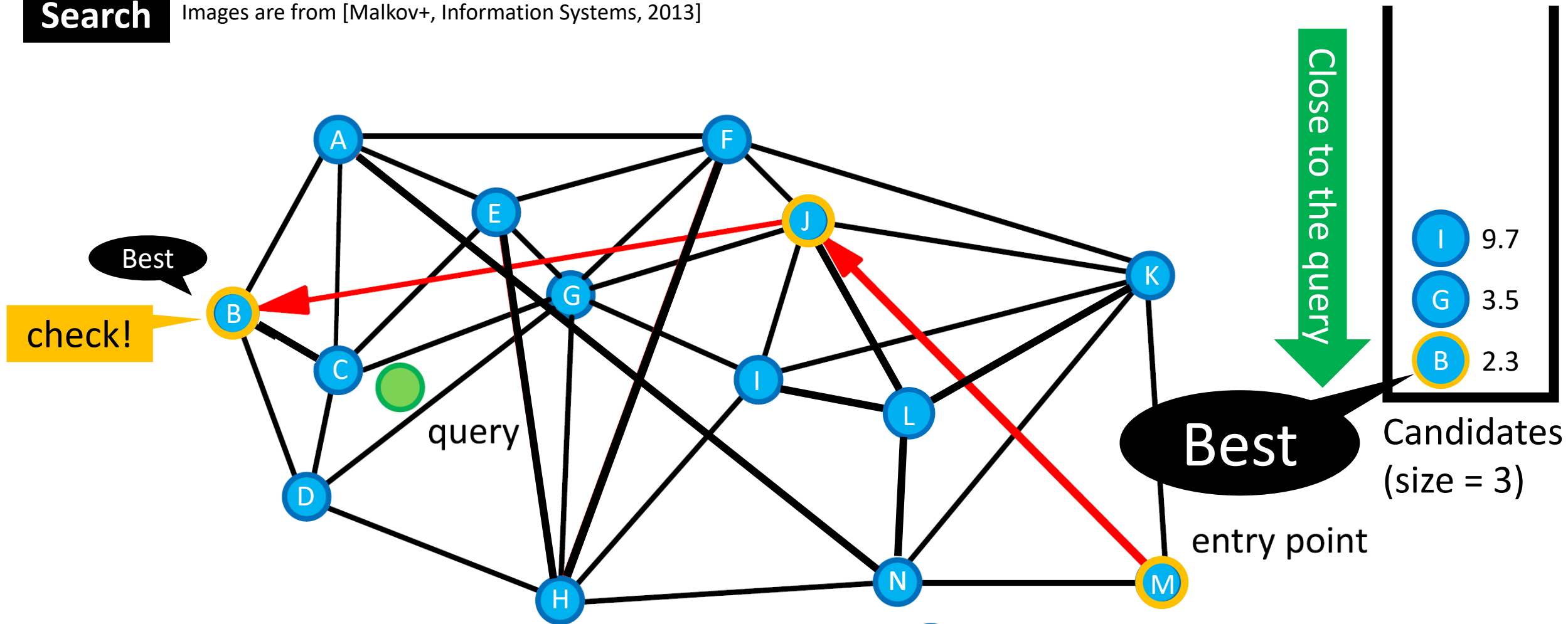


- Pick up the unchecked best candidate (J). Check it.
- Find the connected points.
- Record the distances to q.
- Maintain the candidates (size=3)

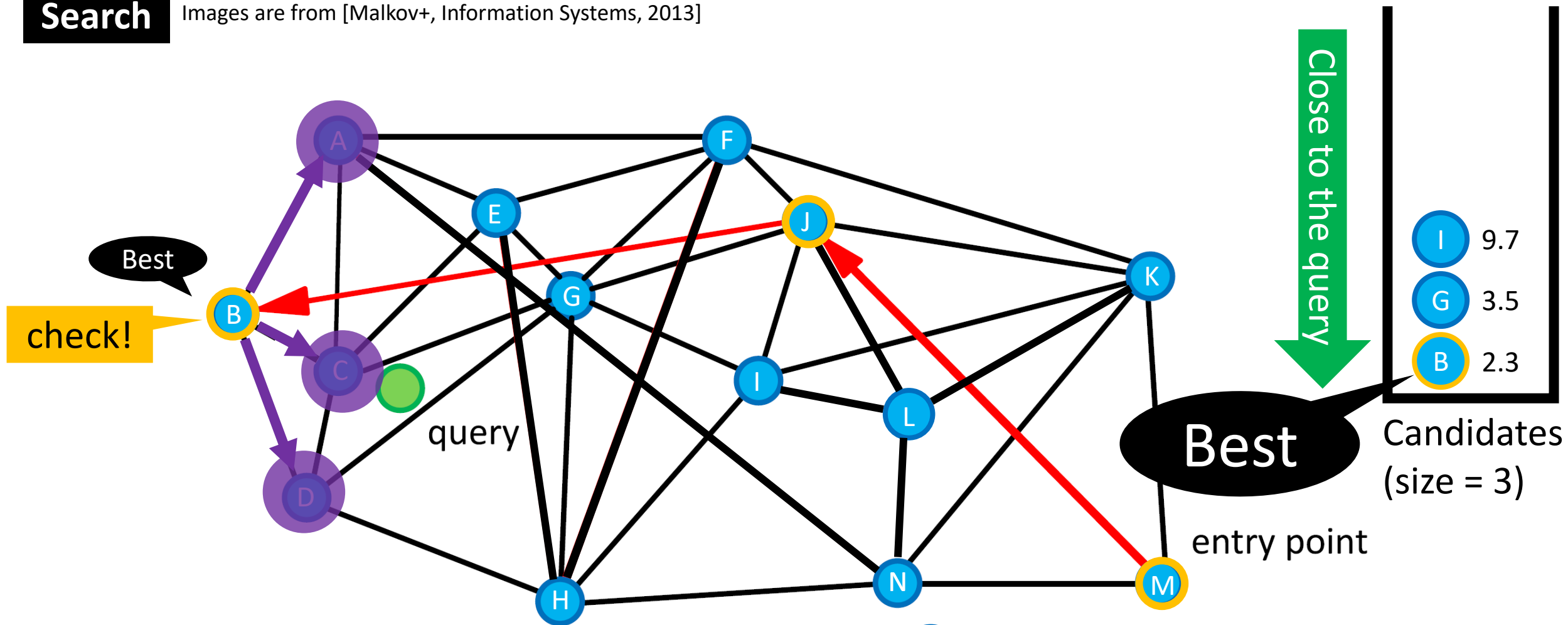




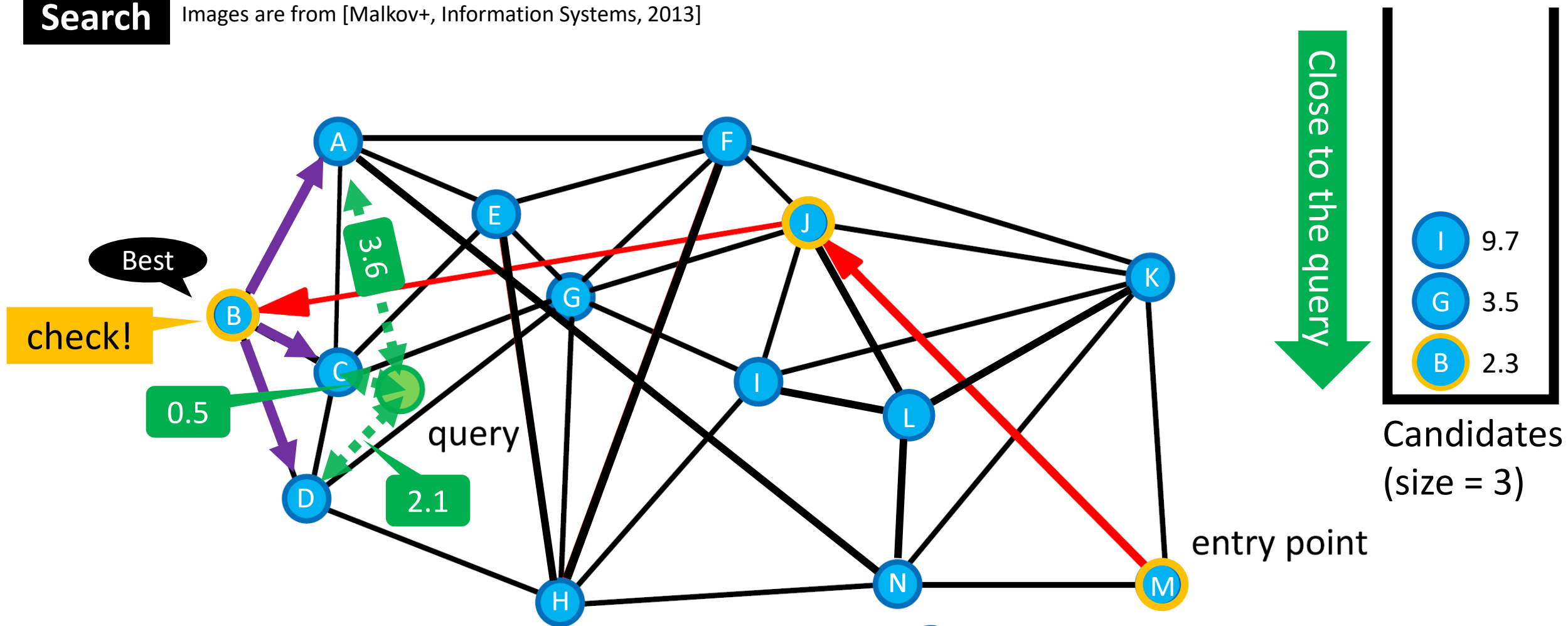




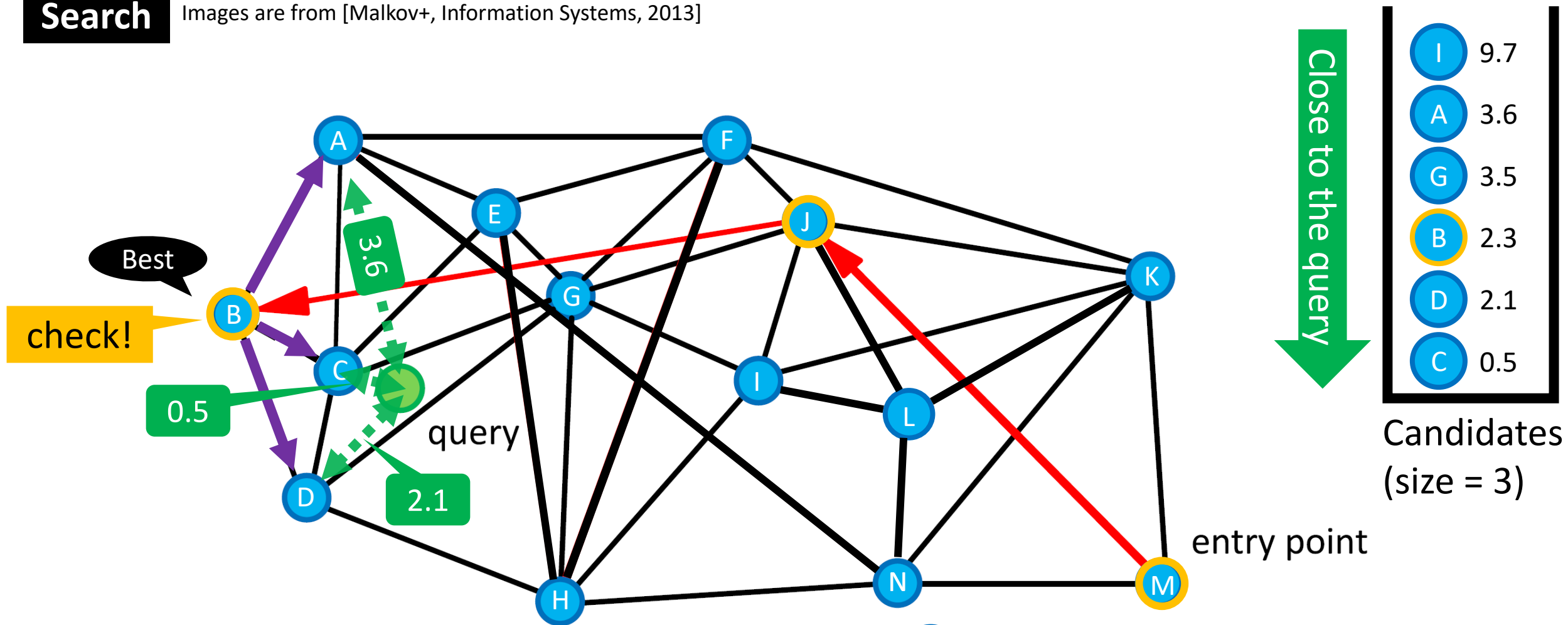
➤ Pick up the unchecked best candidate (B). Check it.



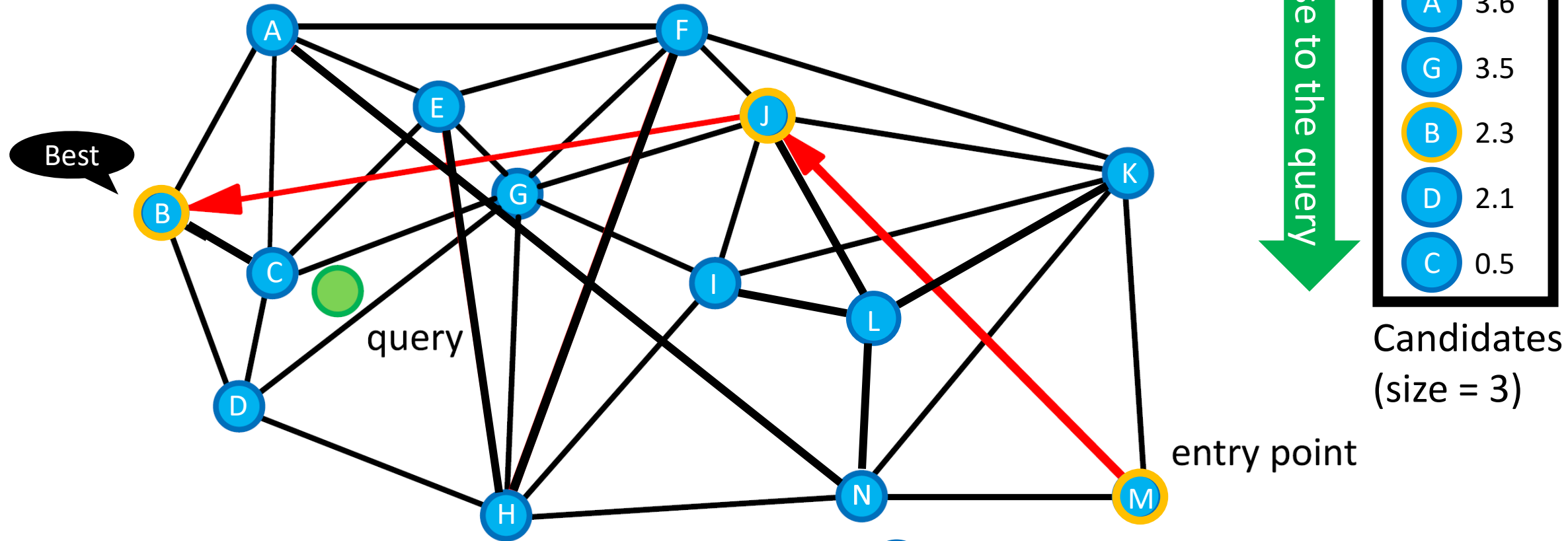
- Pick up the unchecked best candidate (B). Check it.
- Find the connected points.



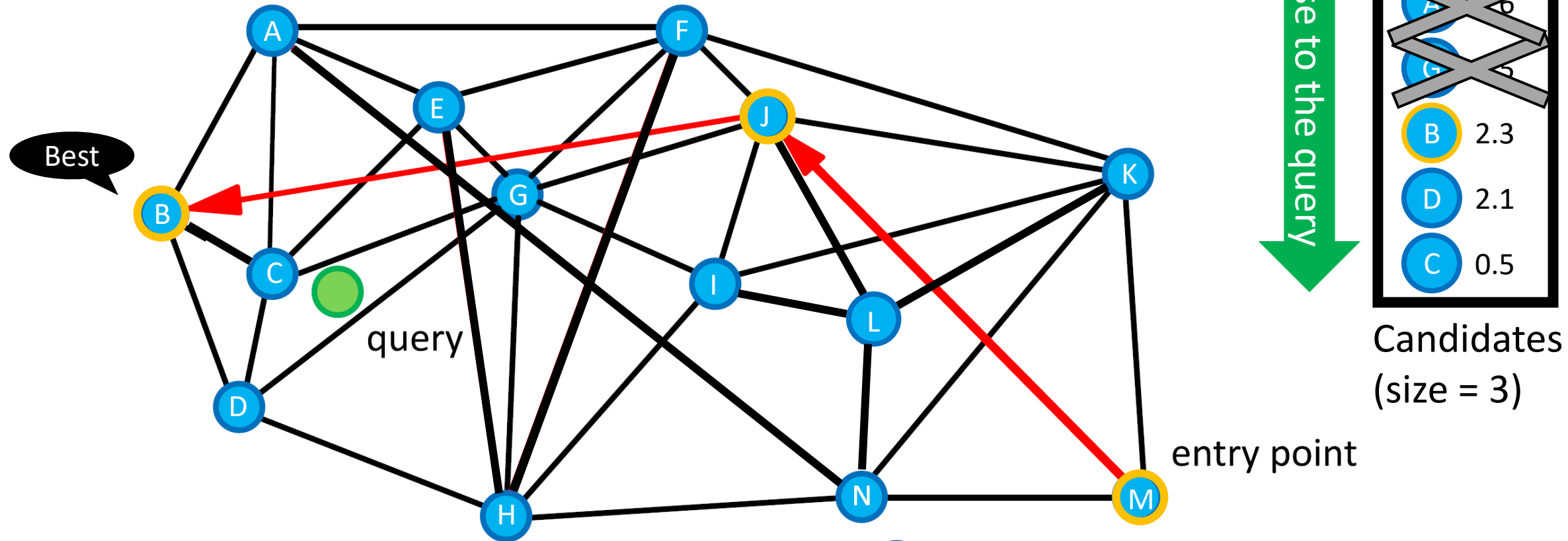
- Pick up the unchecked best candidate (B). Check it.
- Find the connected points.
- Record the distances to q.



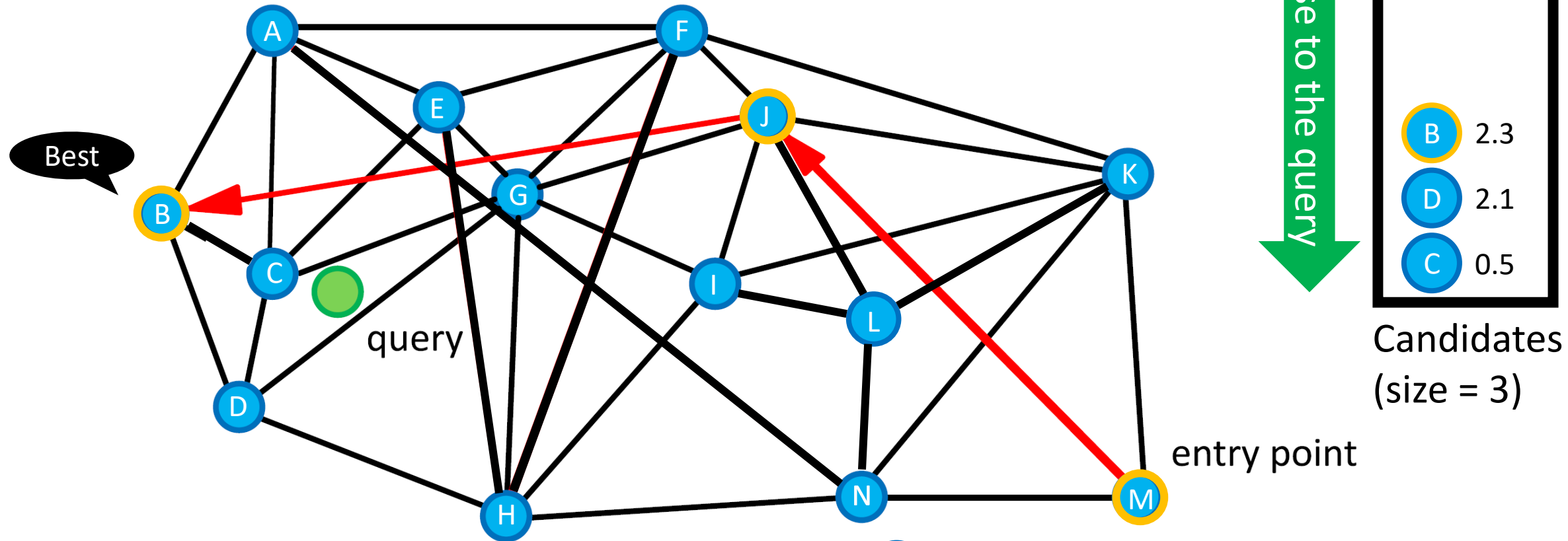
- Pick up the unchecked best candidate (B). Check it.
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- Record the distances to q.



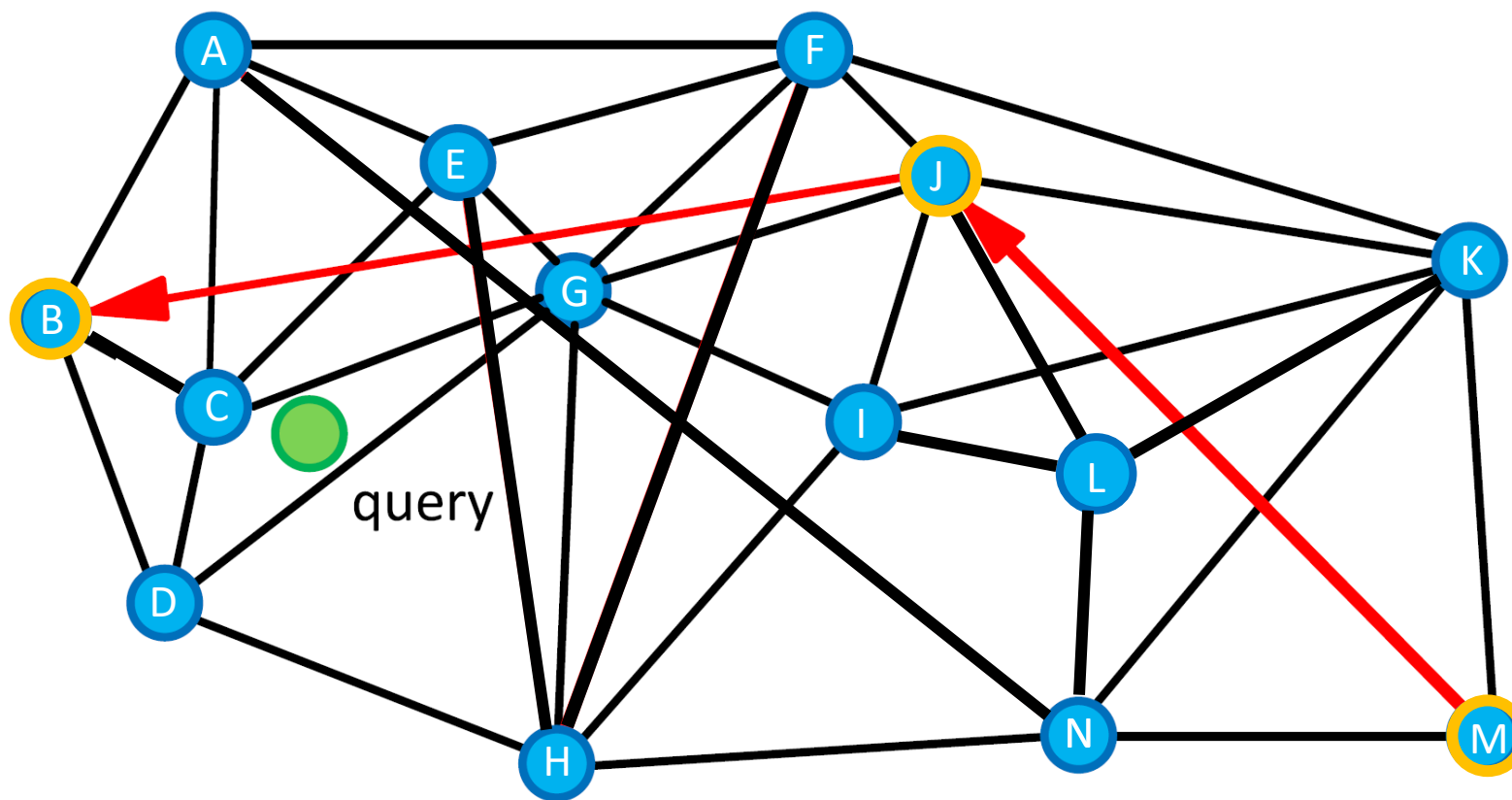
- Pick up the unchecked best candidate (B). Check it.
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- Record the distances to q.



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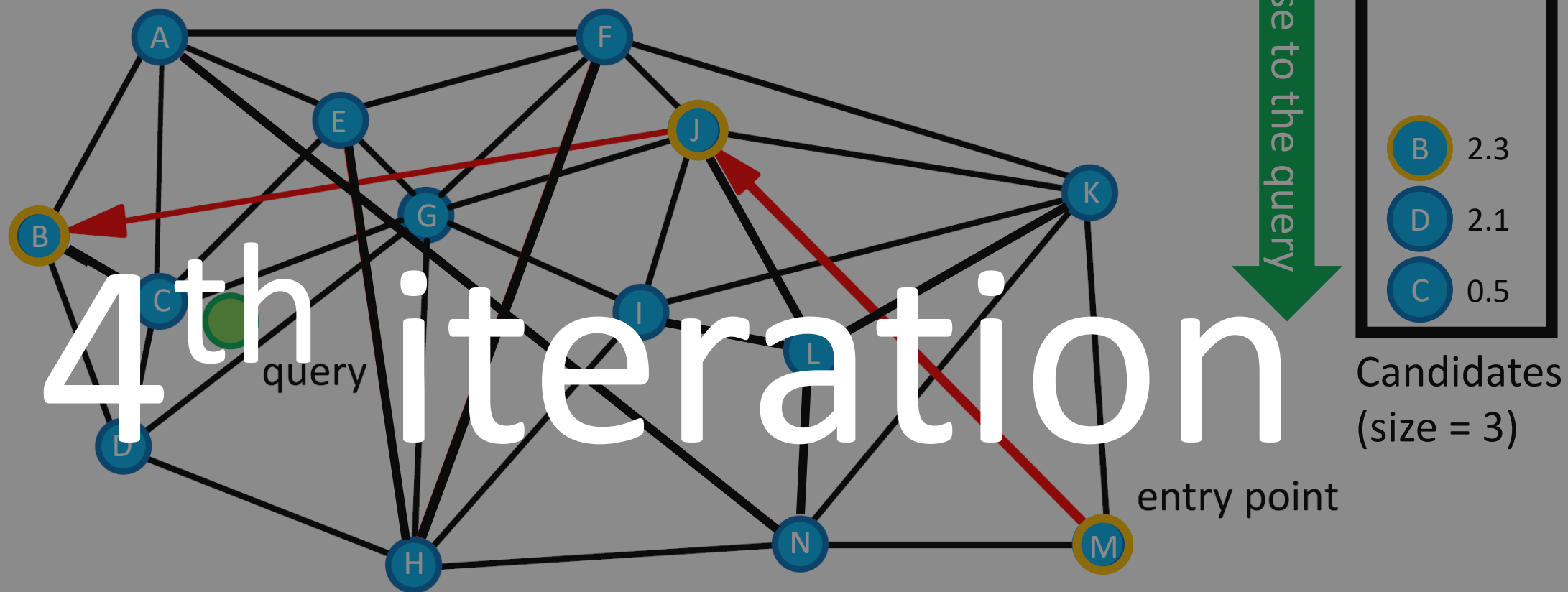


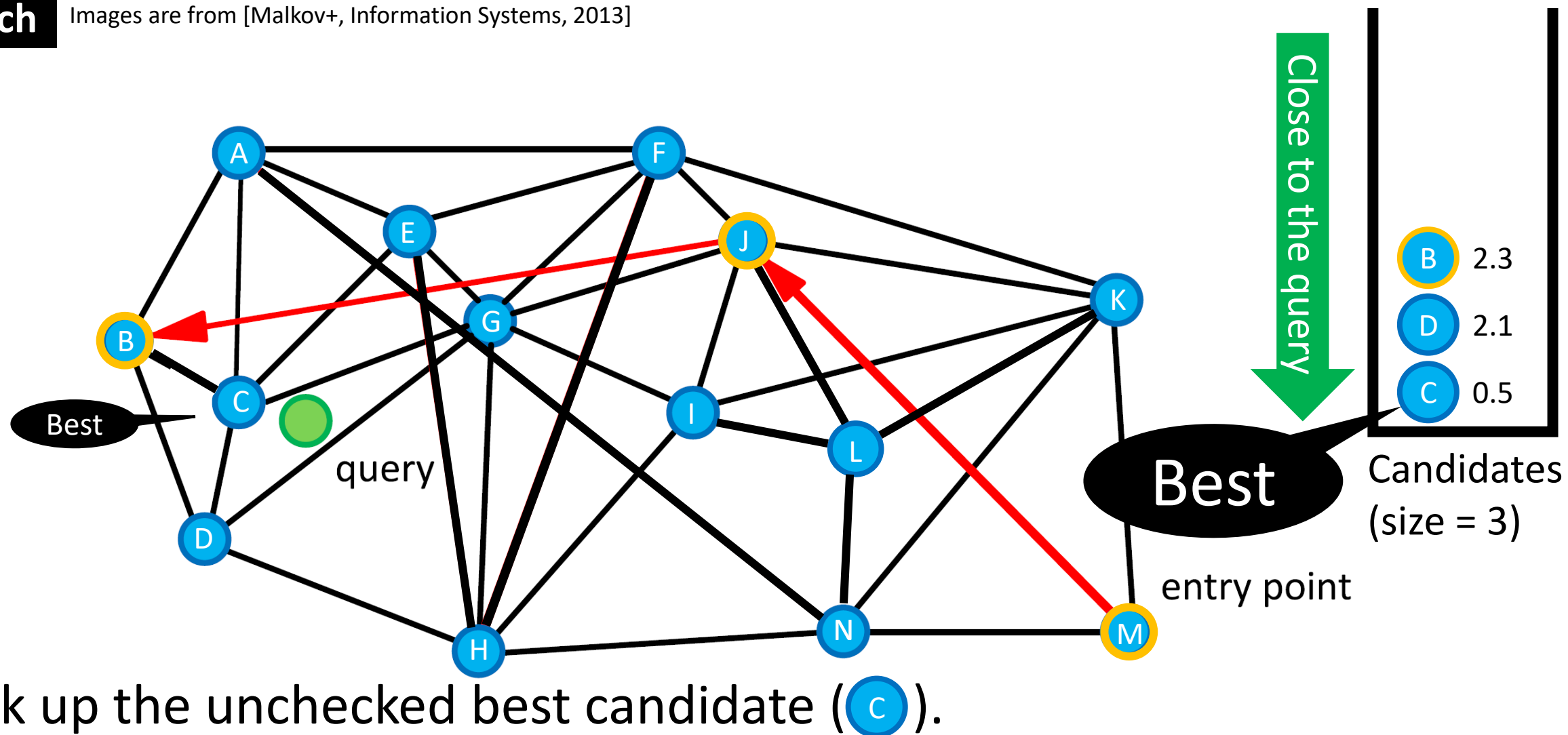
Close to the query

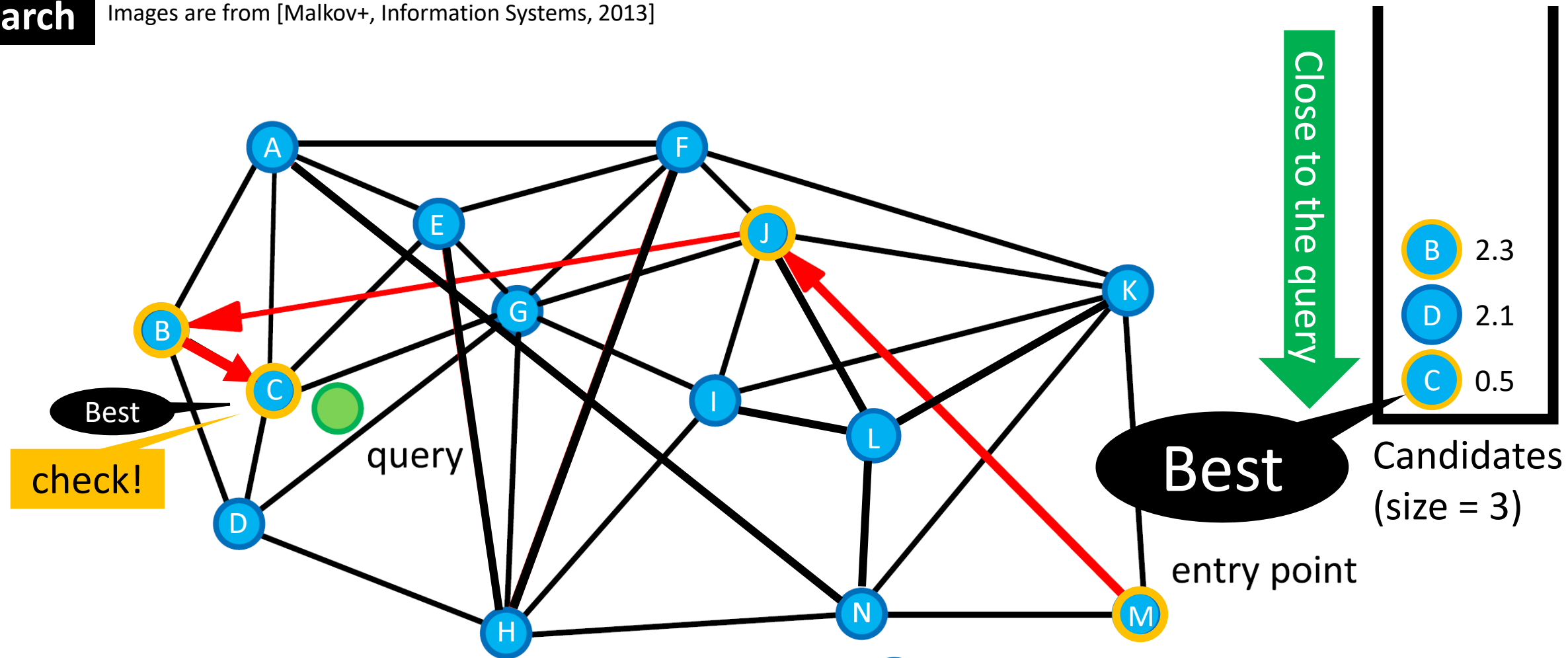
B	2.3
D	2.1
C	0.5

Candidates
(size = 3)

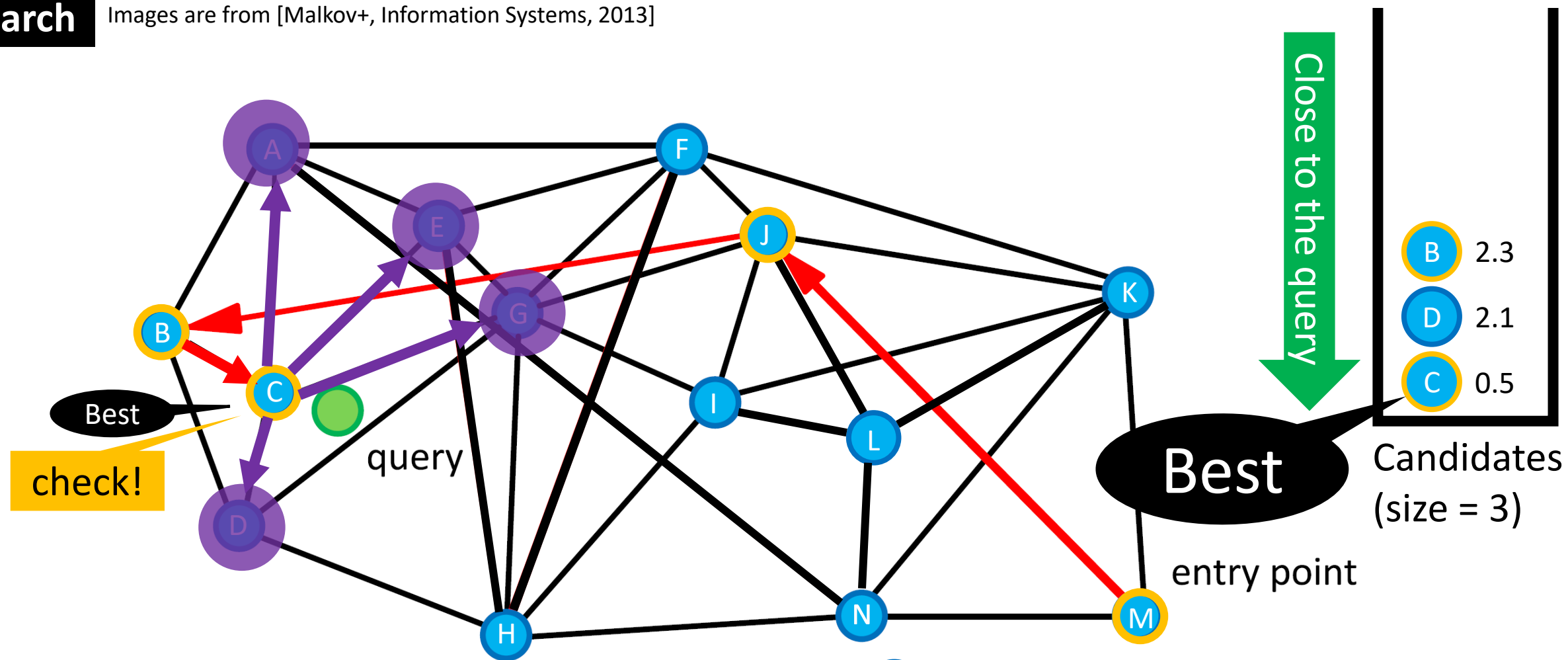
entry point







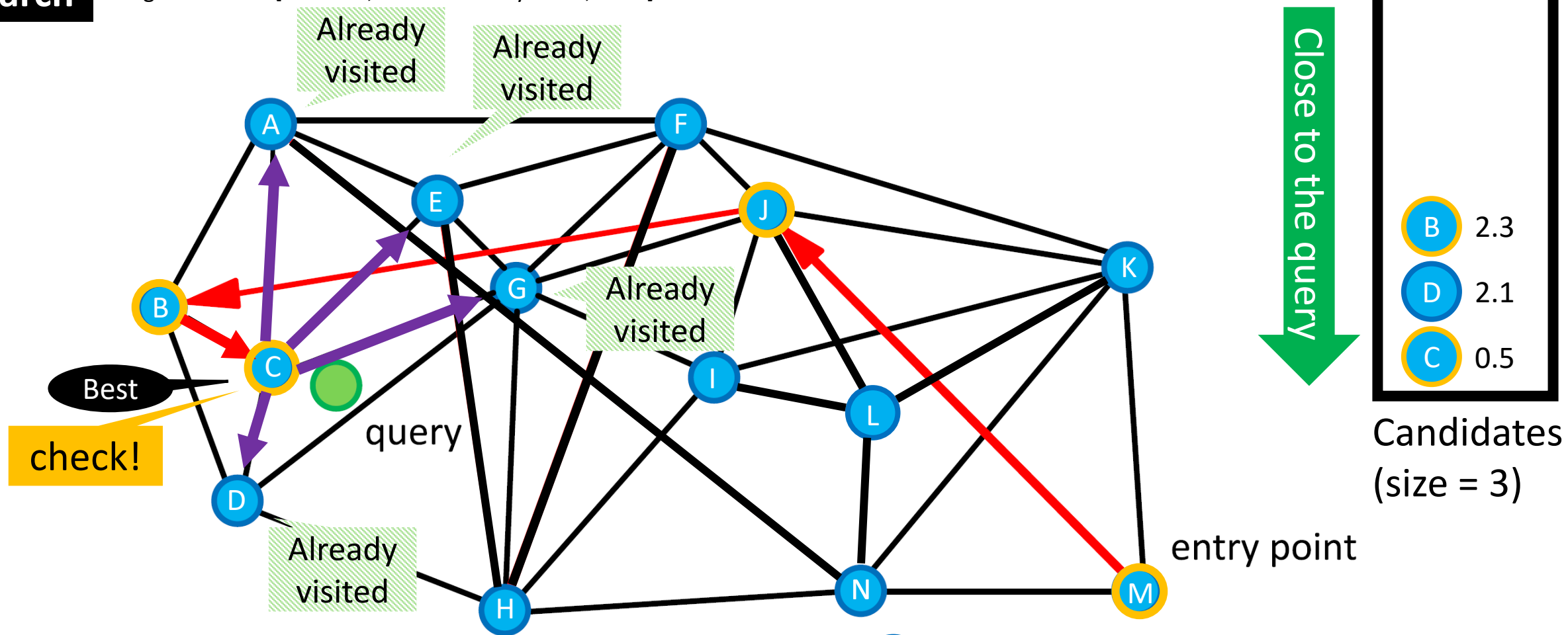
➤ Pick up the unchecked best candidate (C). Check it.



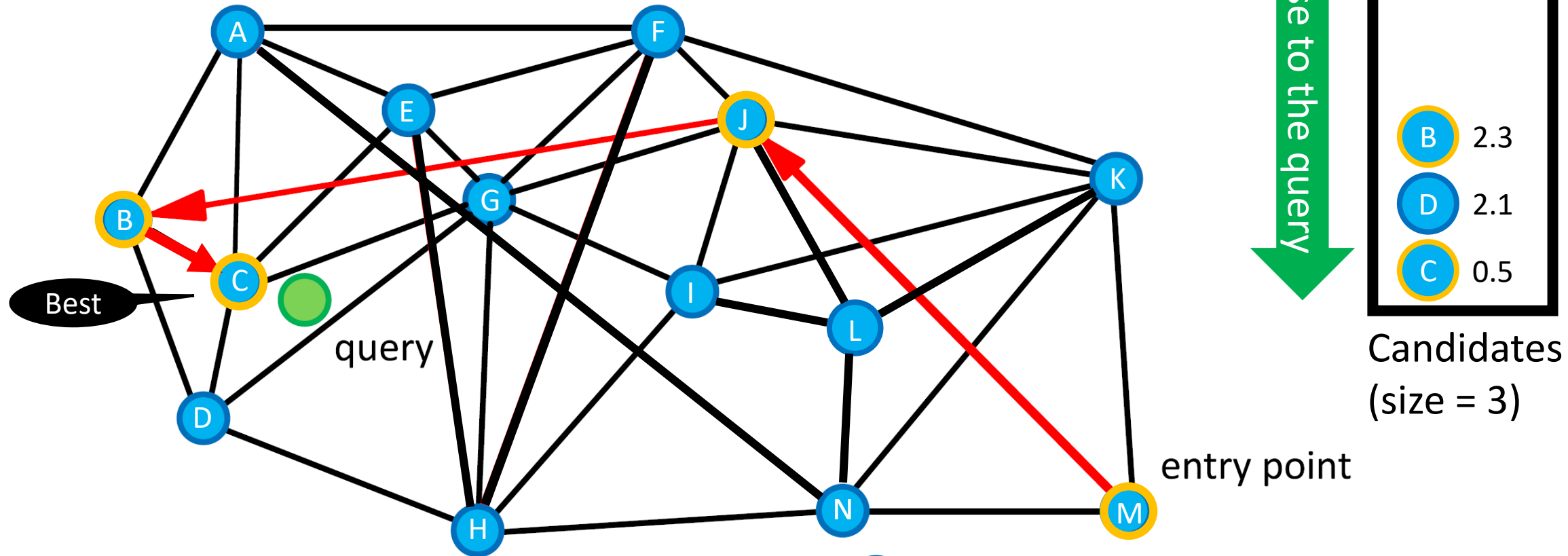
- Pick up the unchecked best candidate (C). Check it.
- Find the connected points.

Search

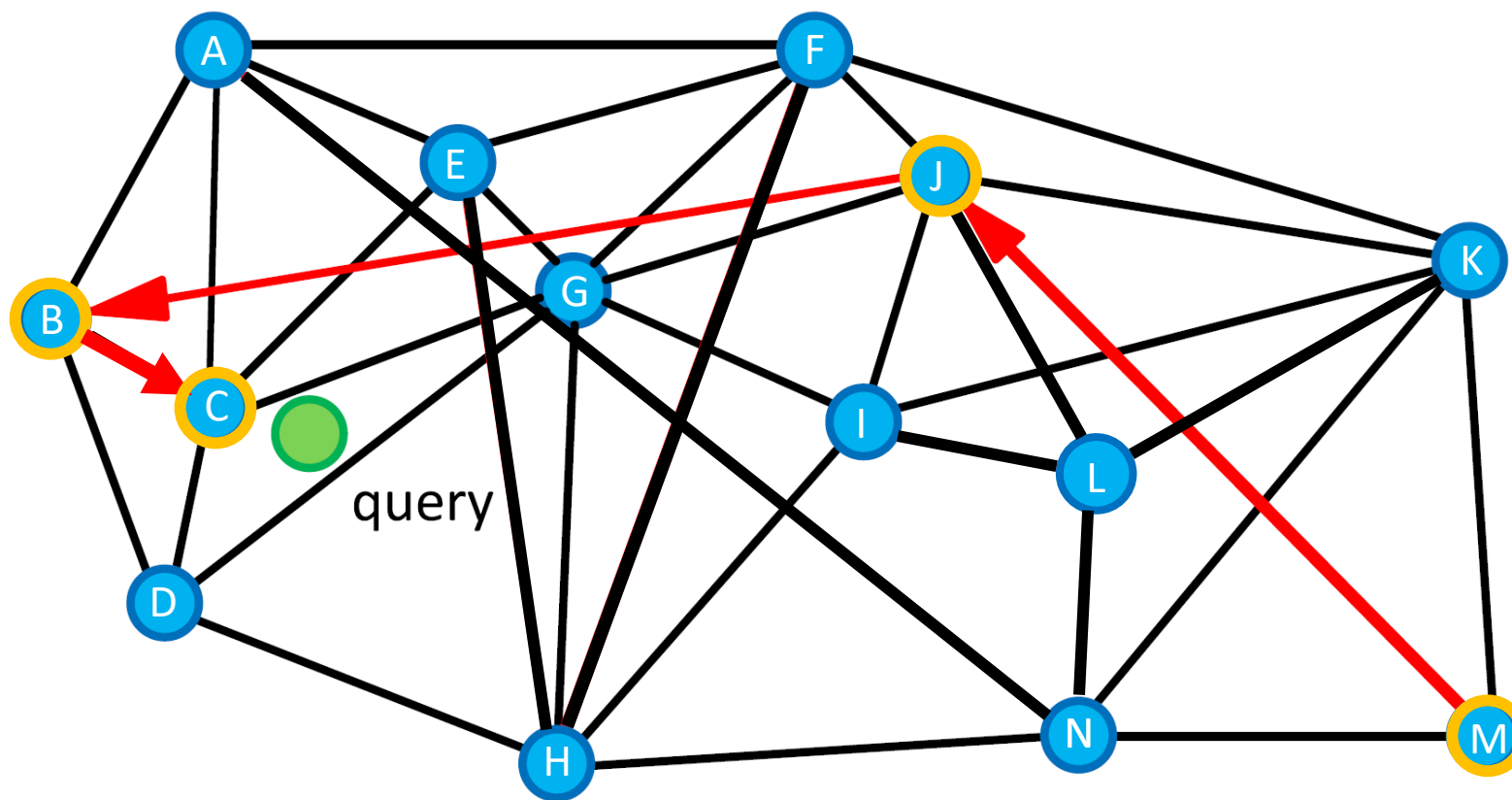
Images are from [Malkov+, Information Systems, 2013]



- Pick up the unchecked best candidate (C). Check it.
- Find the connected points.
- Record the distances to q.



- Pick up the unchecked best candidate (C). Check it.
- Find the connected points.
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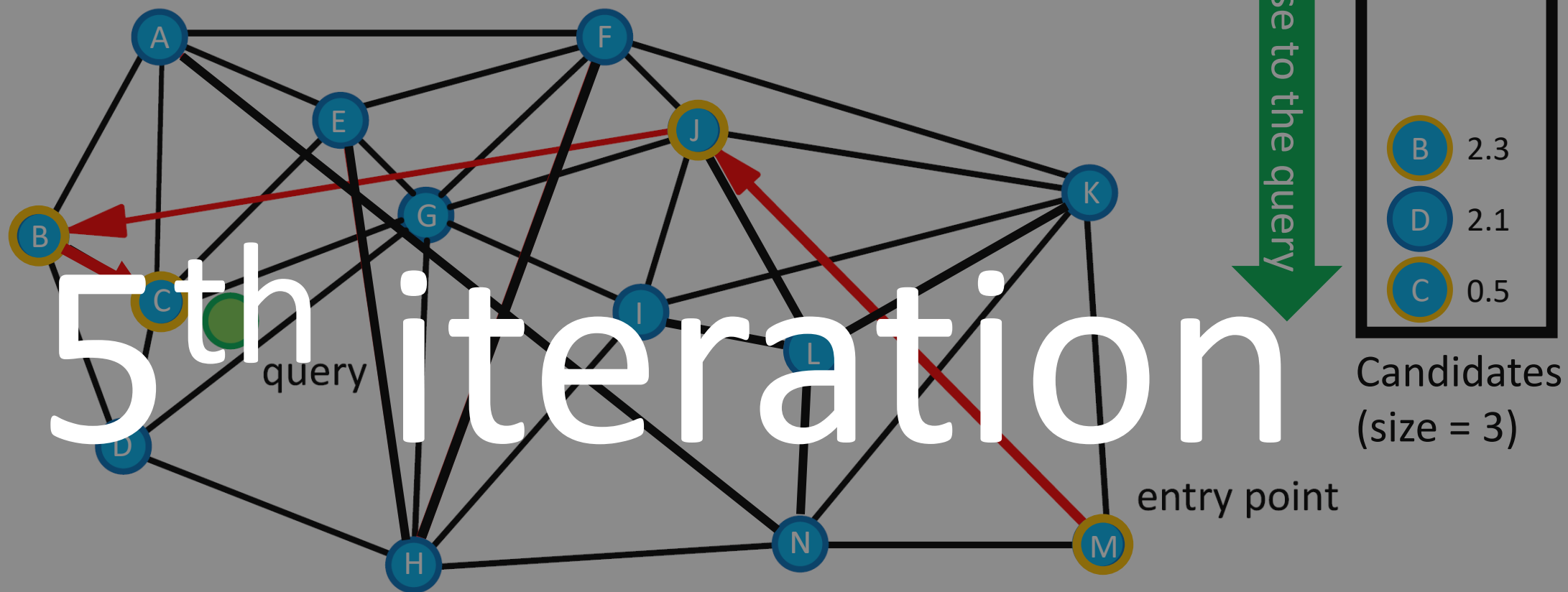


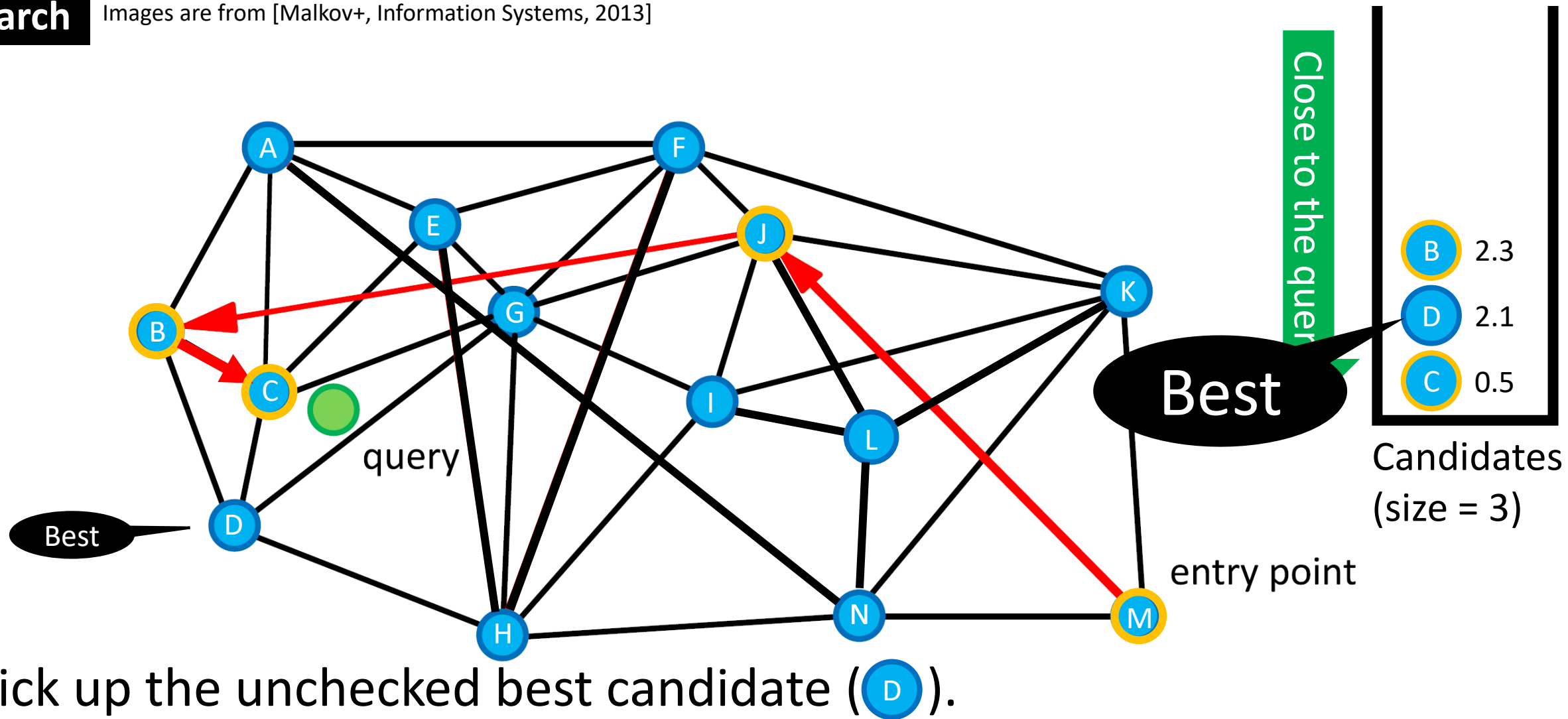
Close to the query

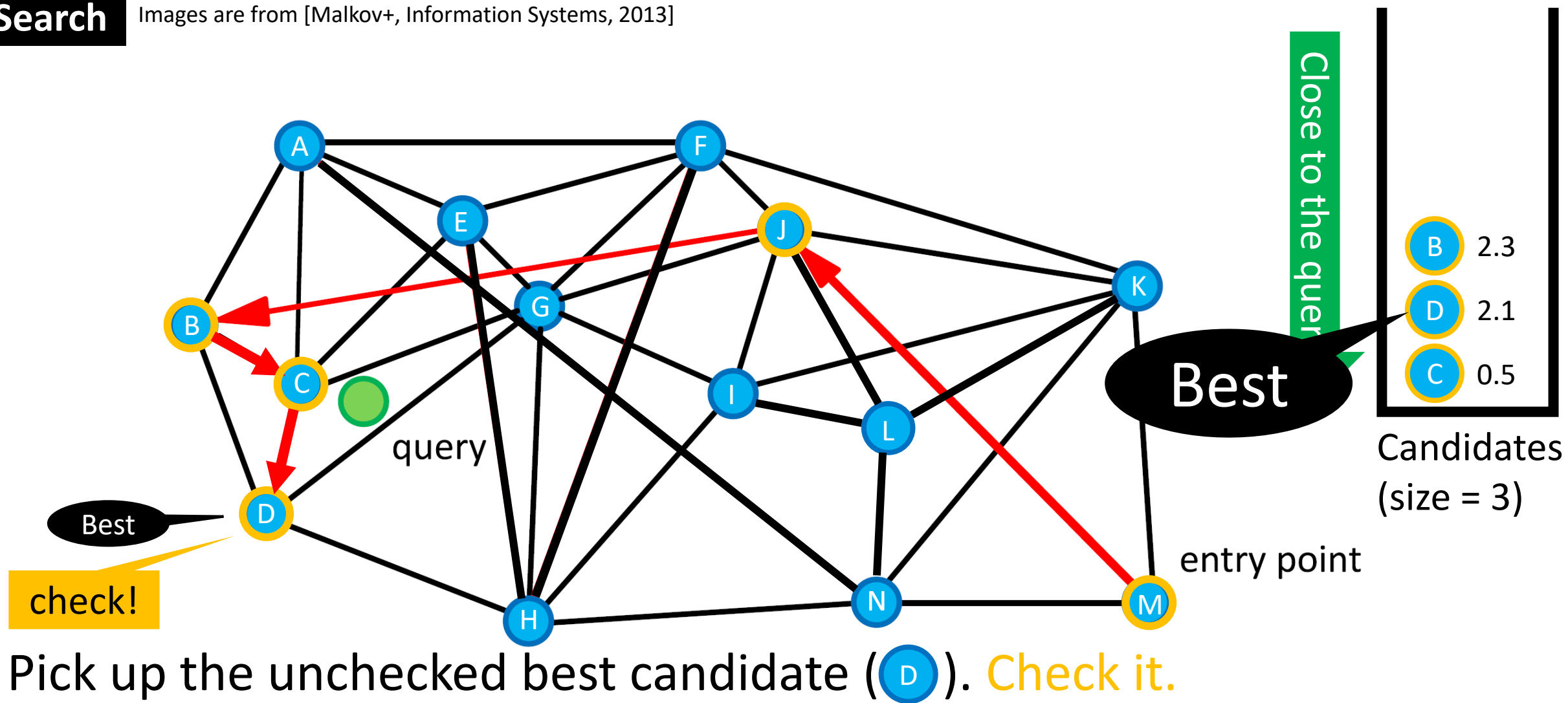
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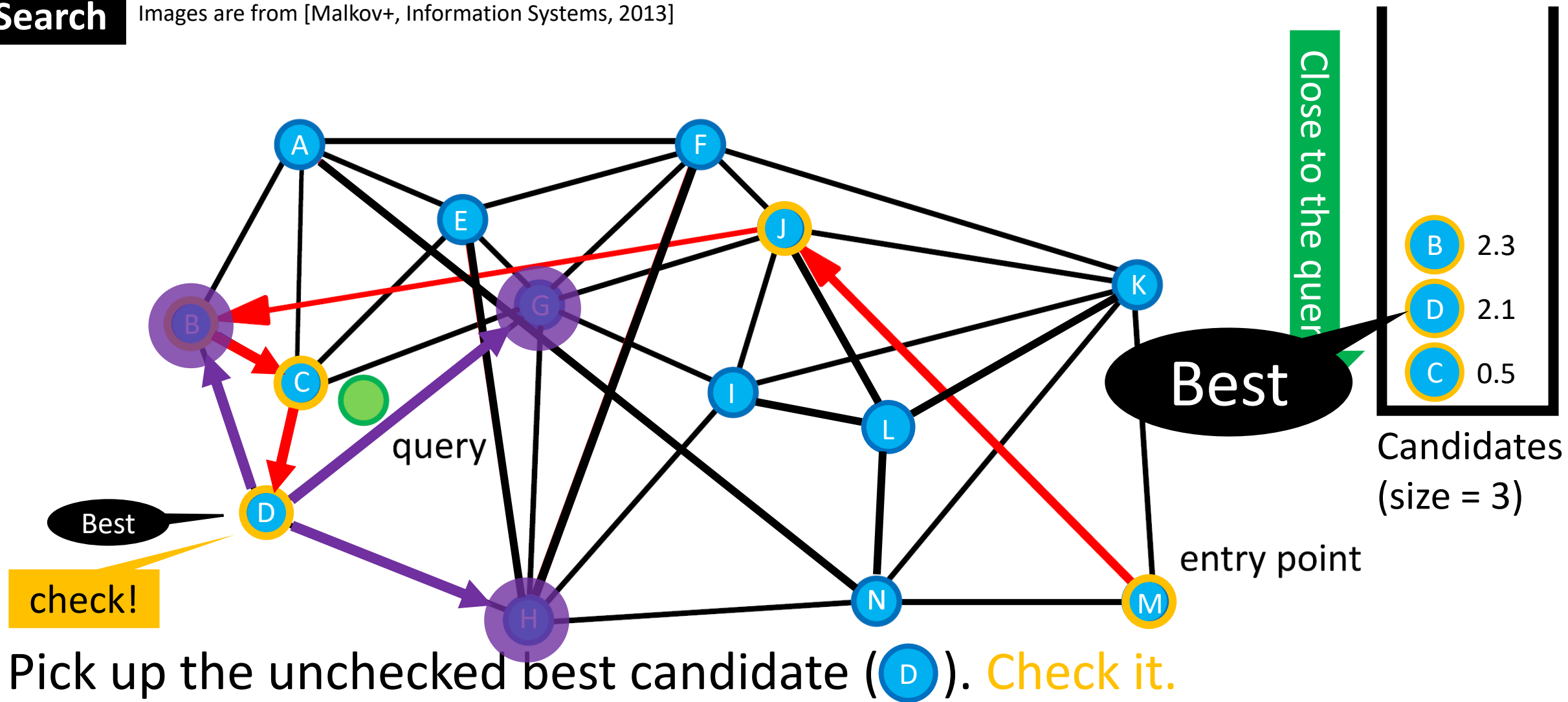
Candidates
(size = 3)

entry point

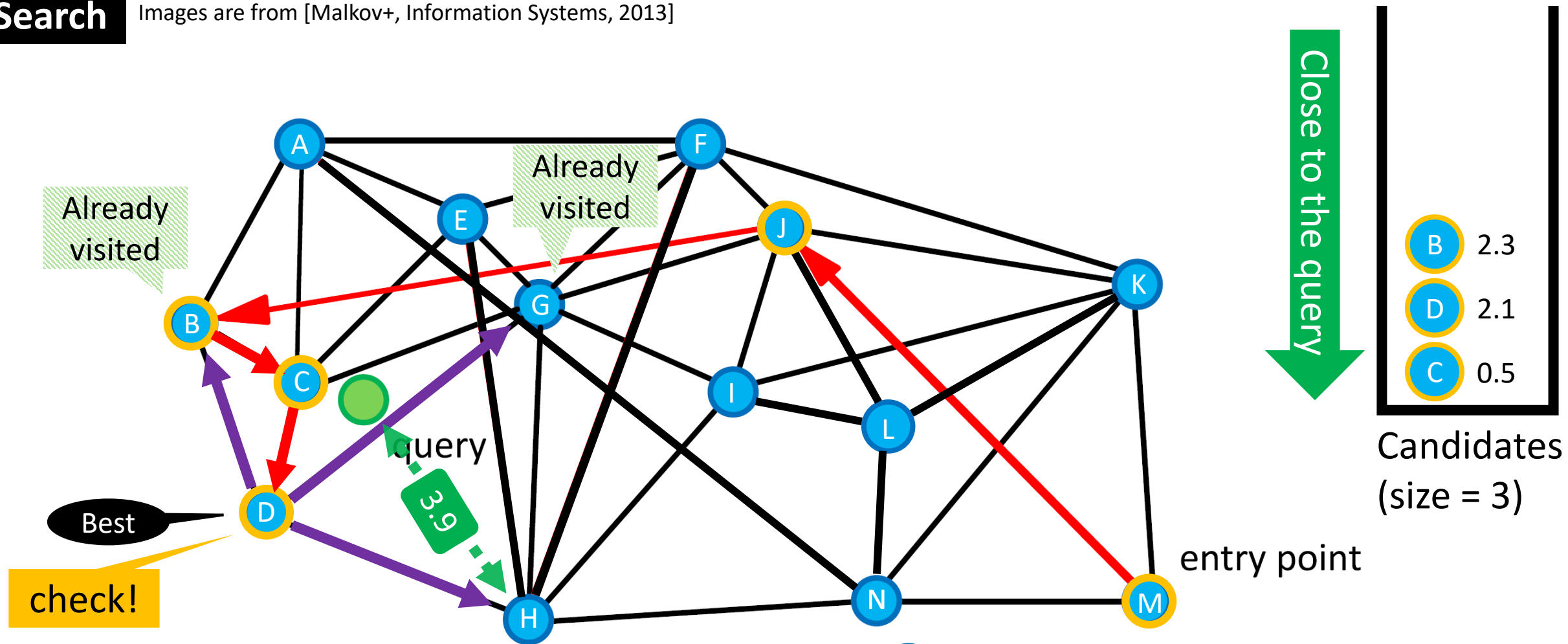




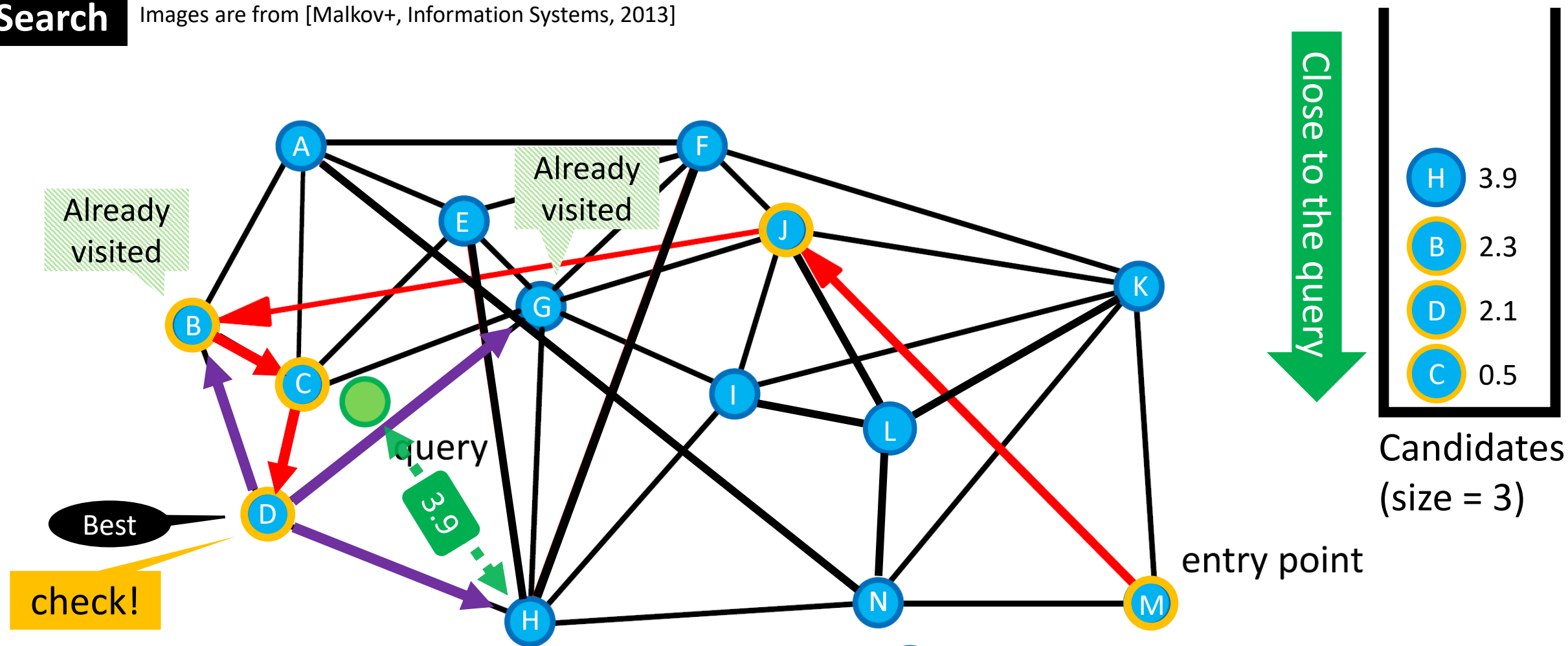




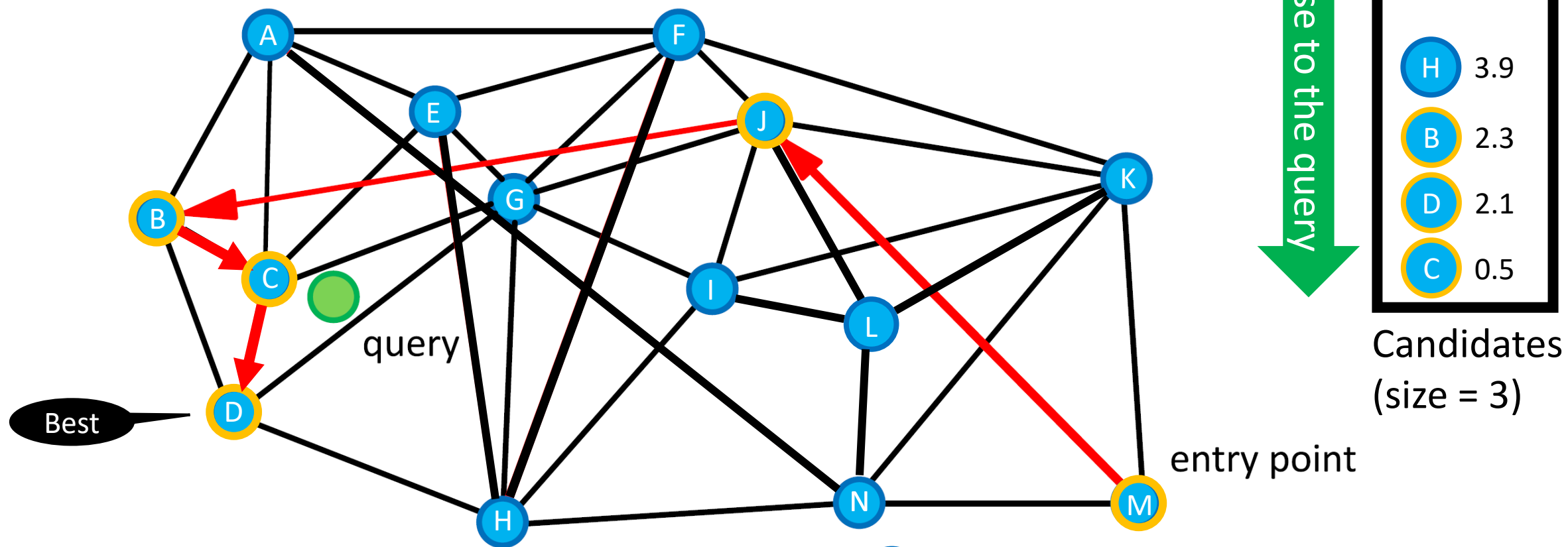
- Pick up the unchecked best candidate (D). Check it.
- Find the connected points.



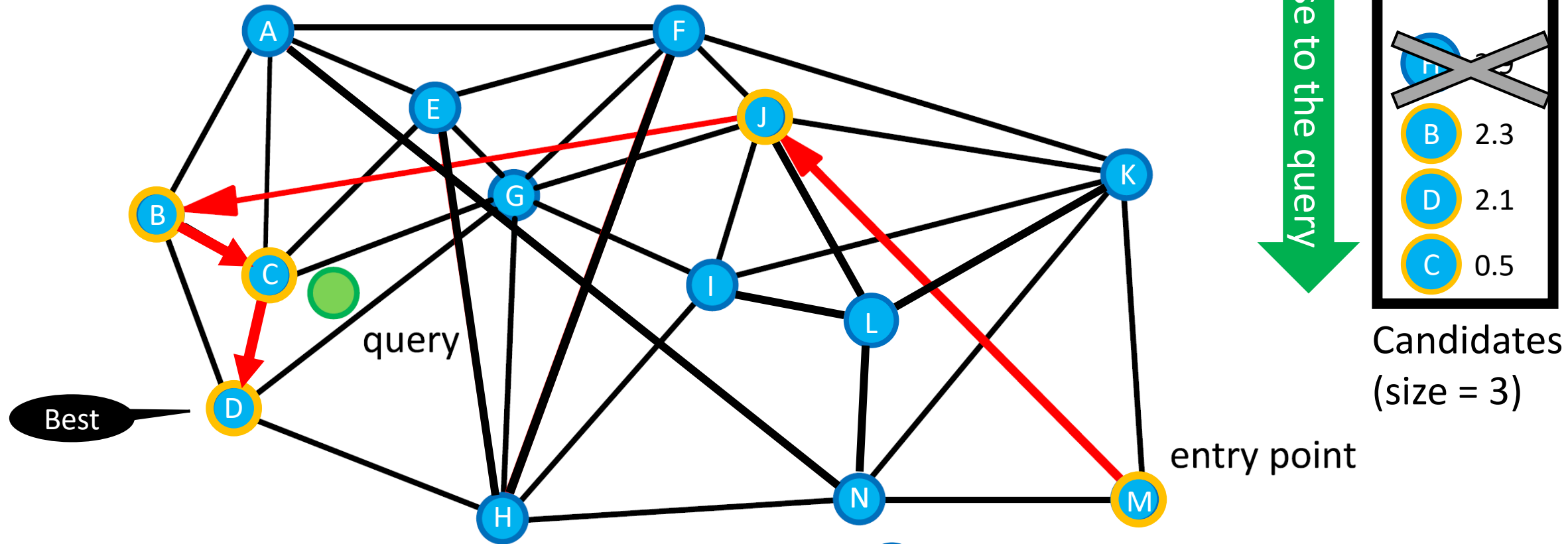
- Pick up the unchecked best candidate (D). Check it.
- Find the connected points.
- Record the distances to q.



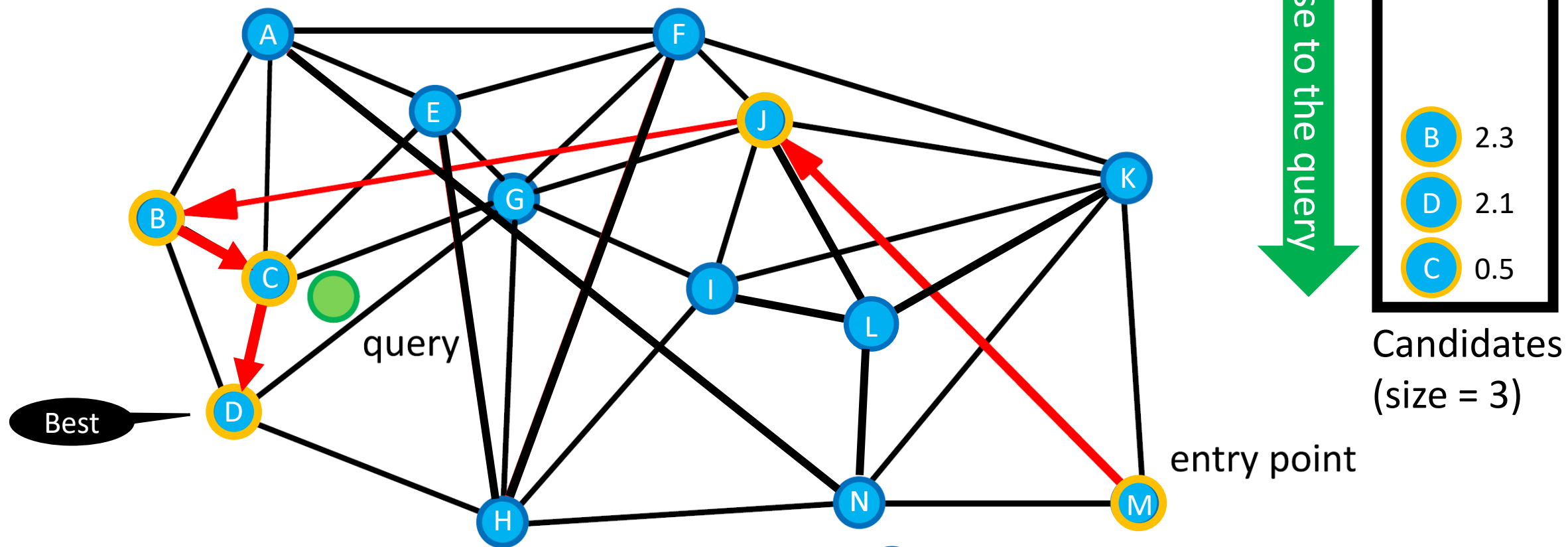
- Pick up the unchecked best candidate (D). Check it.
- Find the connected points.
- Record the distances to q.



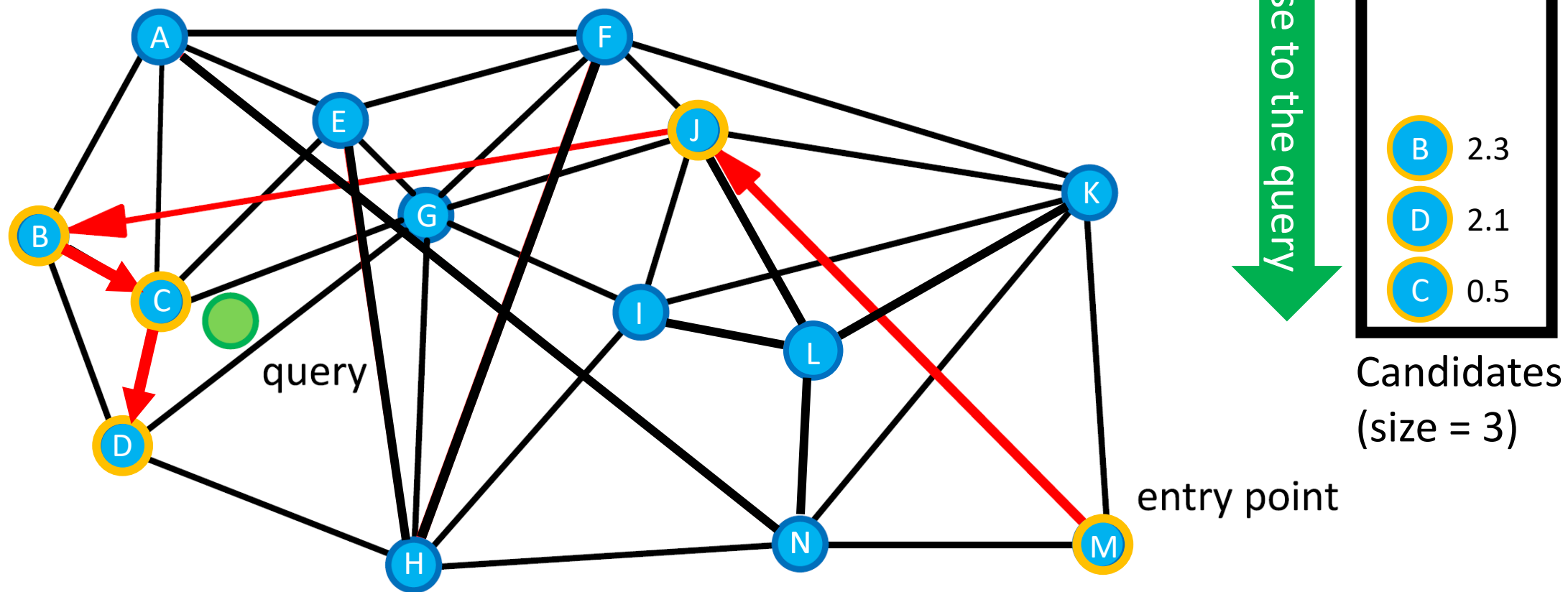
- Pick up the unchecked best candidate (D). Check it.
- Find the connected points.
- Record the distances to q.



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- Find the connected points.
- Record the distances to q.
- Maintain the candidates (size=3)



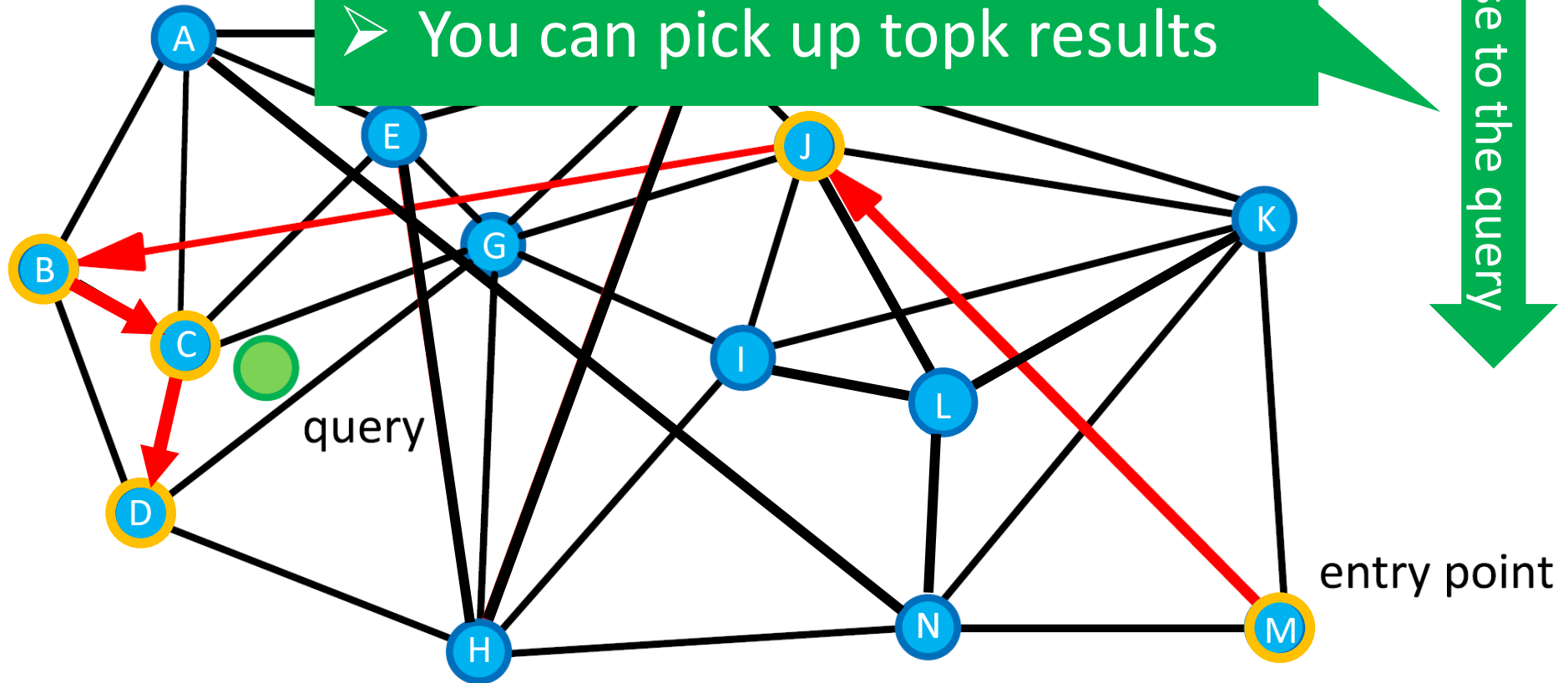
- Pick up the unchecked best candidate (D). Check it.
- Find the connected points.
- Record the distances to q.
- Maintain the candidates (size=3)



- All candidates are **checked**. Finish.
- Here, **c** is the closet to the query (●)

Final output 1: Candidates

➤ You can pick up topk results



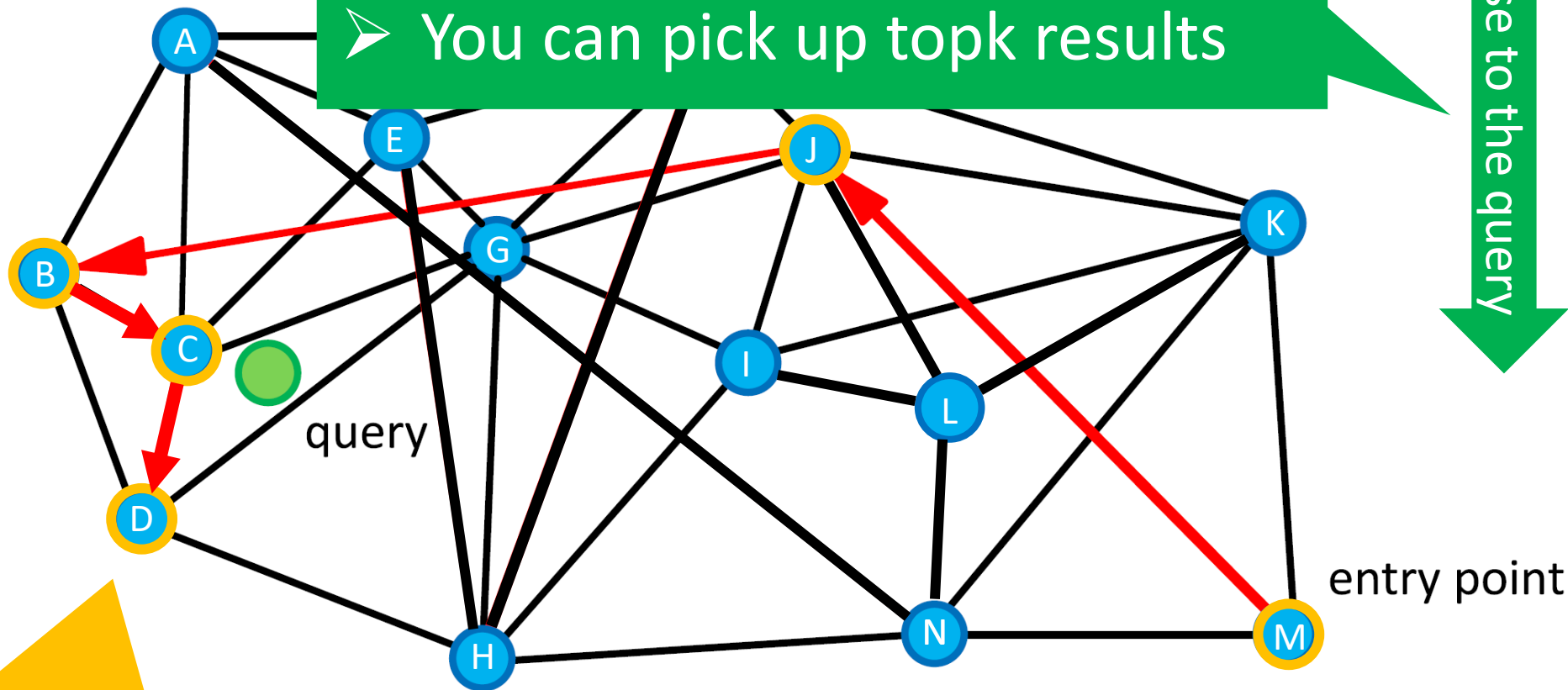
B	2.3
D	2.1
C	0.5

Candidates
(size = 3)

- All candidates are **checked**. Finish.
- Here, **C** is the closet to the query (●)

Final output 1: Candidates

➤ You can pick up topk results



B	2.3
D	2.1
C	0.5

Candidates
(size = 3)

➤ All candidates are checked. Finish.

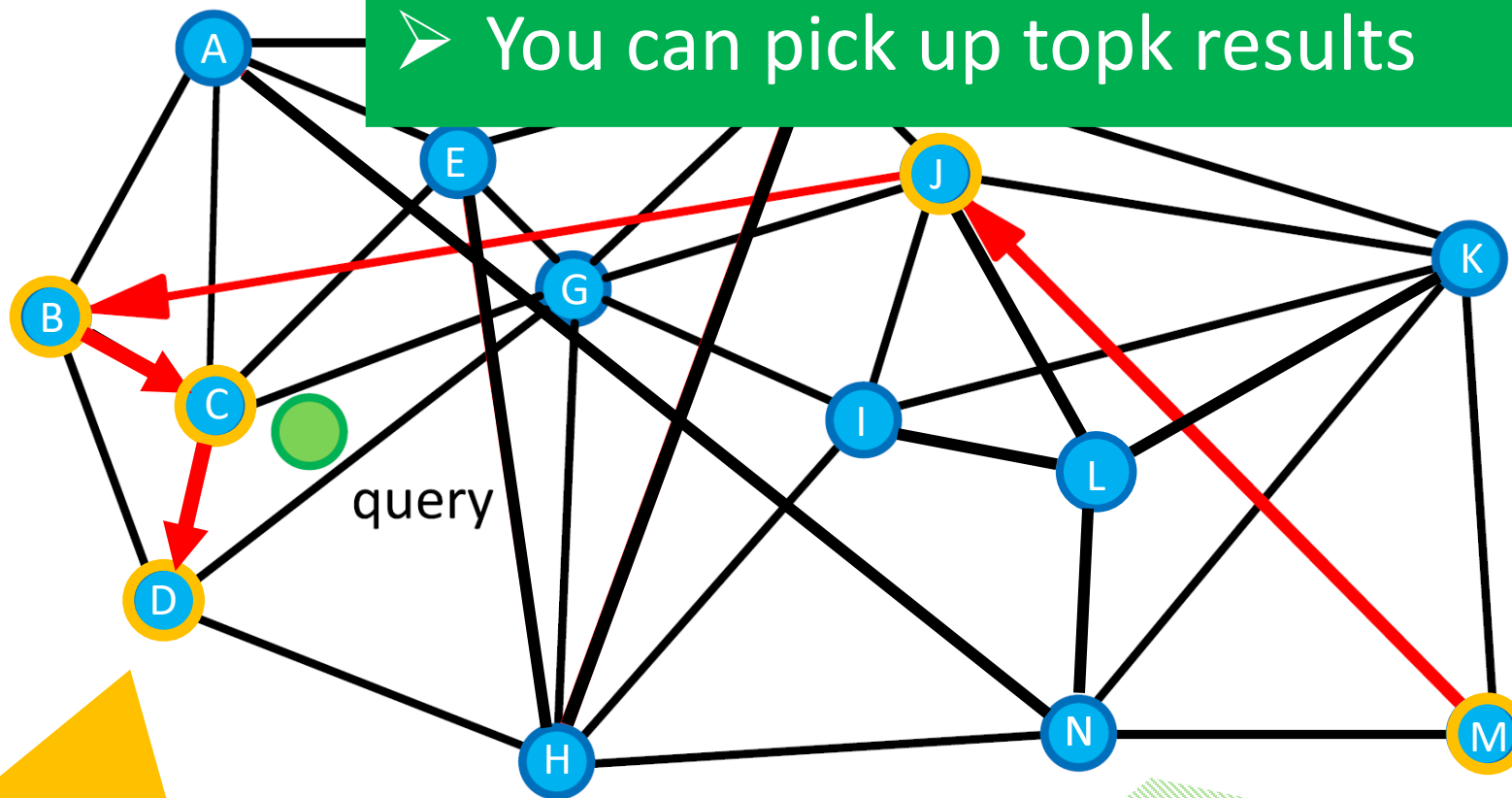
Final output 2: Checked items

➤ i.e., search path

query (●)

Final output 1: Candidates

➤ You can pick up topk results



Close to the query

B	2.3
D	2.1
C	0.5

Candidates
(size = 3)

entry point

➤ All candidates are checked. Finish

Final output 2: Checked items

➤ i.e., search path

Final output 3: Visit flag

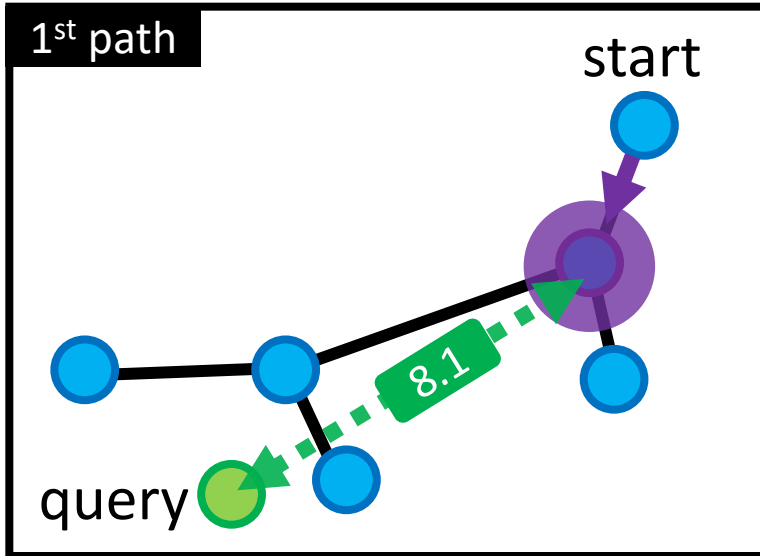
➤ For each item, visited or not

Observation: runtime

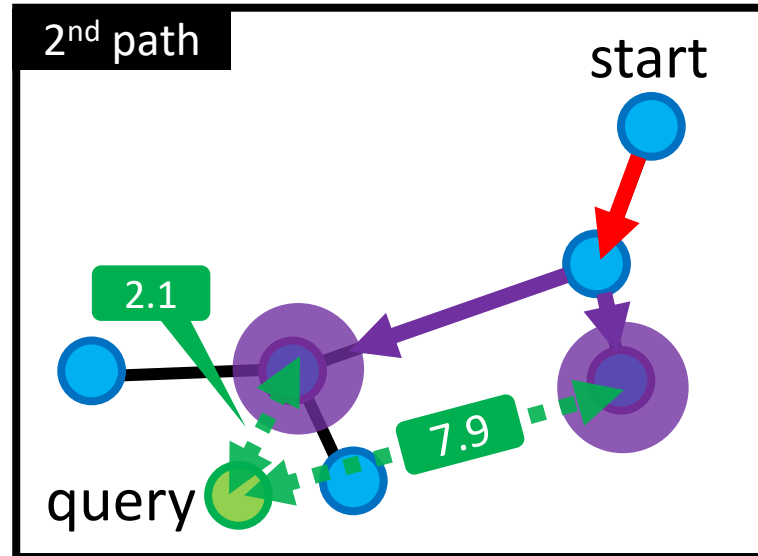
- Item comparison takes time; $O(D)$



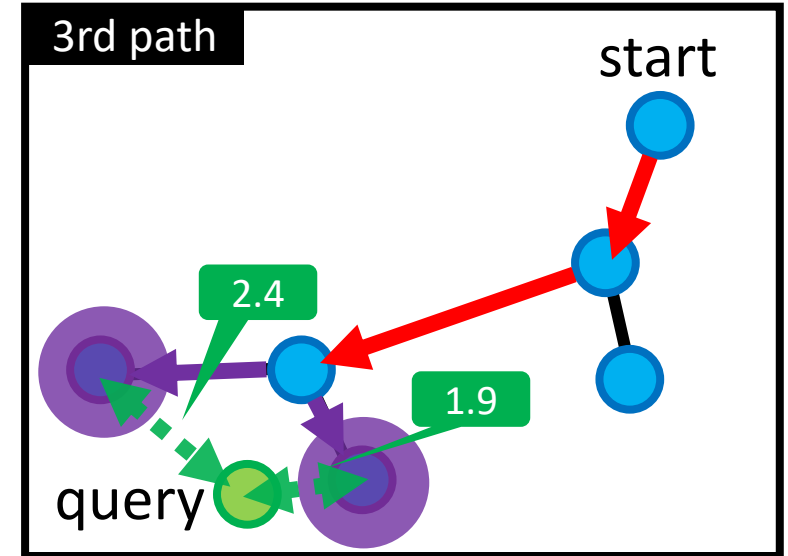
- The overall runtime $\sim \text{\#item_comparison}$
 $\sim \text{length_of_search_path} * \text{average_outdegree}$



outdegree = 1



outdegree = 2



outdegree = 2

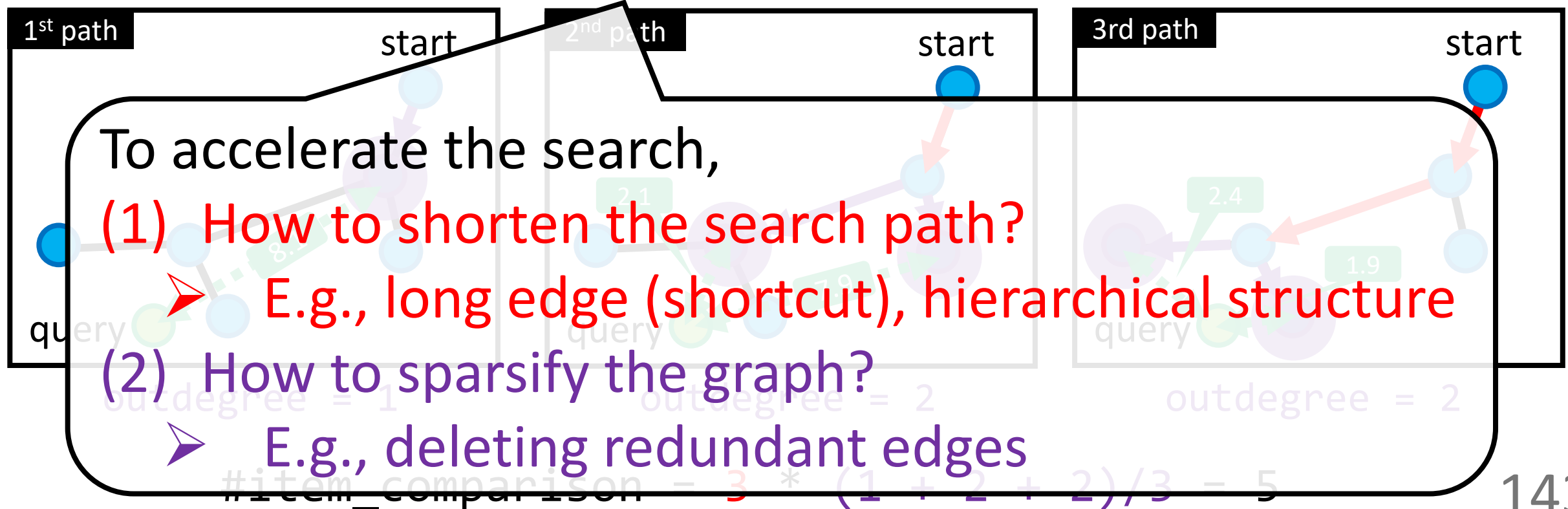
$$\#item_comparison = 3 * (1 + 2 + 2)/3 = 5$$

Observation: runtime

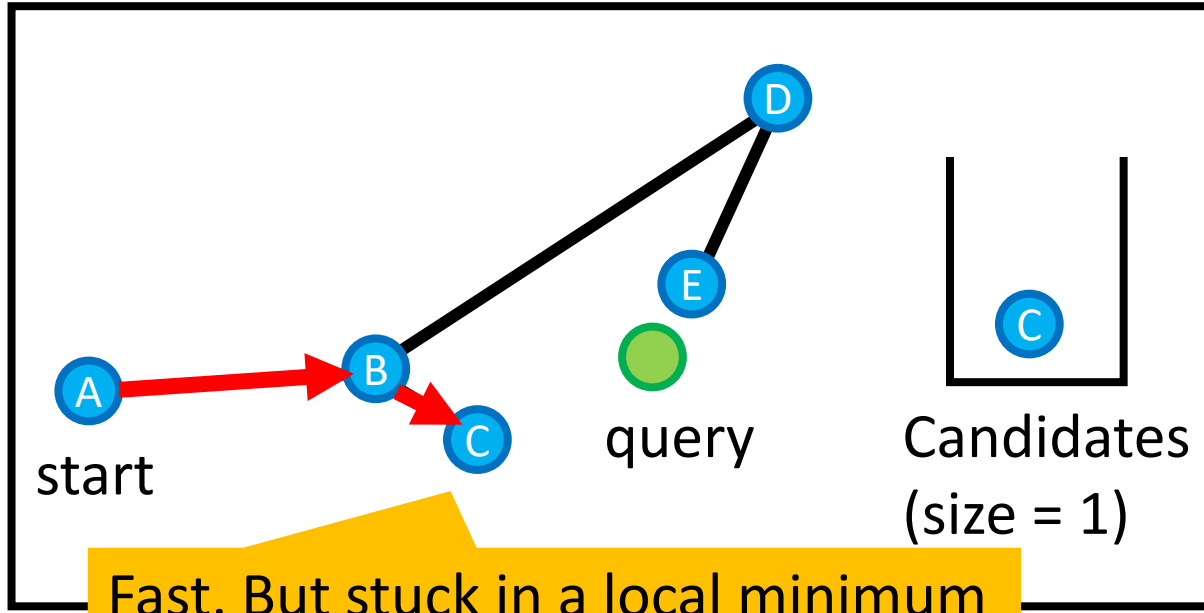
- Item comparison takes time; $O(D)$



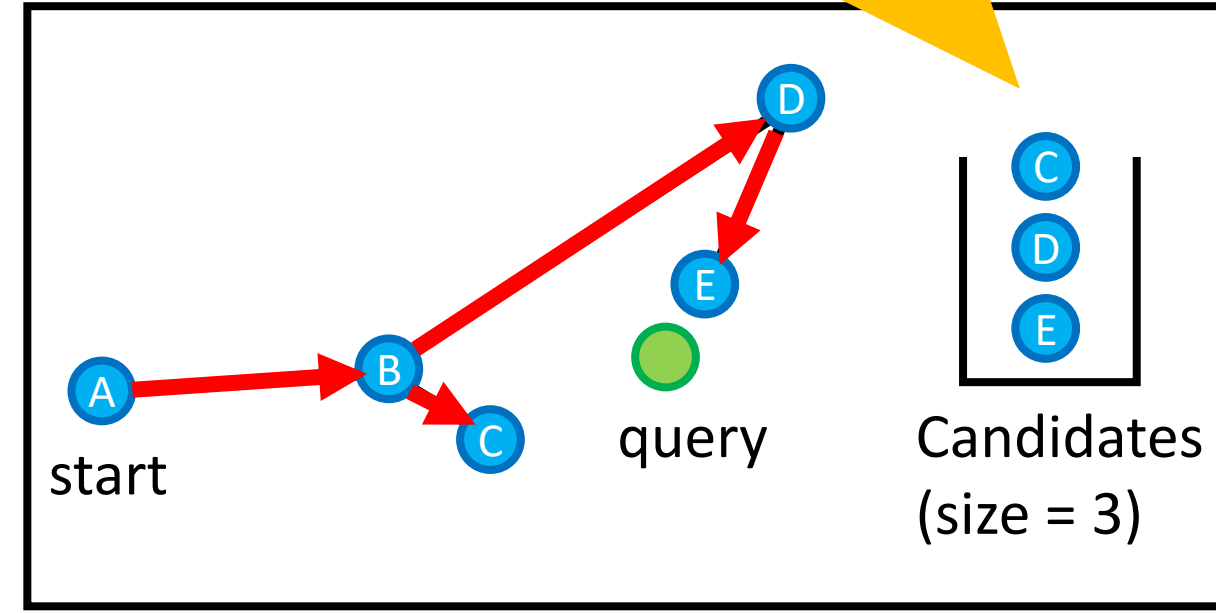
- The overall runtime $\sim \text{\#item_comparison}$
 $\sim \text{length_of_search_path} * \text{average_outdegree}$



Observation: candidate size



size = 1: Greedy search



size > 1: Beam search

- Larger candidate size, better but slower results
- Online parameter to control the trade-off
- Called “ef” in HNSW

Pseudo code

Algorithm 1. Search-on-Graph(G, p, q, l)

Require: graph G , start node p , query point q , candidate pool size l

Ensure: k nearest neighbors of q

```
1:  $i = 0$ , candidate pool  $S = \emptyset$ 
2:  $S.add(p)$ 
3: while  $i < l$  do
4:    $i =$  the id of the first unchecked node  $p_i$  in  $S$ 
5:   mark  $p_i$  as checked
6:   for all neighbor  $n$  of  $p_i$  in  $G$  do
7:     if  $n$  has not been visited then
8:        $S.add(n)$ 
9:     end if
10:  end for
11:  sort  $S$  in ascending order of the distance to  $q$ 
12:  if  $S.size() > l$  then
13:     $S.resize(l)$  // remove nodes from back of  $S$  to keep its
14:    size no larger than  $l$ 
15:  end if
16: end while
17: return the first  $k$  nodes in  $S$ 
```

NSG [Cong+, VLDB 19]

Algorithm 1: GreedySearch(s, x_q, k, L)

Data: Graph G with start node s , query x_q , result size k , search list size $L \geq k$

Result: Result set $\mathcal{L}$ containing k -approx NNs, and a set $\mathcal{V}$ containing all the visited nodes

```
begin
  initialize sets  $\mathcal{L} \leftarrow \{s\}$  and  $\mathcal{V} \leftarrow \emptyset$ 
  while  $\mathcal{L} \setminus \mathcal{V} \neq \emptyset$  do
    let  $p^* \leftarrow \arg \min_{p \in \mathcal{L} \setminus \mathcal{V}} \|x_p - x_q\|$ 
    update  $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$  and
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    if  $|\mathcal{L}| > L$  then
      update  $\mathcal{L}$  to retain closest  $L$ 
        points to  $x_q$ 
    end if
  end while
  return [closest  $k$  points from  $\mathcal{L}$ ;  $\mathcal{V}$ ]
```

DiskANN [Subramanya+, NeurIPS 19]

Algorithm 1 Beam search

Data: graph G , query q , initial vertex v_0 , output size k

Initialization:

$V = \{v_0\}$ // a set of visited vertices

$H = \{v_0 : d(v_0, q)\}$ // a heap of candidates

```
while has runtime budget do
   $v_i = \text{SelectNearest}(H, q)$ 
  for  $\hat{v} \in \text{Expand}(v_i, G)$  do
    if  $\hat{v} \notin V$  then
       $V := \text{Add}(V, \hat{v})$ 
       $H := \text{Insert}(H, \hat{v}, d(\hat{v}, q))$ 
    end if
  end for
end while
return TopK( $V, q, k$ )
```

Learning to route [Baranchuk+, ICML 19]

- All papers have totally different pseudo code 😞
- Principles are the same. But small details are different.
- Hint: Explicitly state the data structure or not

Pseudo code

Candidates are stored
in an array

Algorithm 1. Search-on-Graph(G, p, q, l)

Require: Graph G , start node p , query point q , candidate pool size l

Ensure: k nearest neighbors of q

```

1:  $i = 0$ , candidate pool  $S = \emptyset$ 
2:  $S.add(p)$ 
3: while  $i < l$  do
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15:  end if
16: end while
17: return the first  $k$  nodes in  $S$ 

```

NSG [Cong+, VLDB 19]

➤ Sort the array explicitly

➤ Principles are the same. But small details are different.

➤ Hint: Explicitly state the data structure or not

Candidates are stored
in a set

Algorithm 1: GreedySearch(s, x_q, k, L)

Data: Graph G with start node s , query x_q , result size k , search list size $L \geq k$

Result: Result set $\mathcal{L}$ containing k -approx NNs, and a set $\mathcal{V}$ containing all the visited nodes

```

begin
  initialize sets  $\mathcal{L} \leftarrow \{s\}$  and  $\mathcal{V} \leftarrow \emptyset$ 
  while  $\mathcal{L} \setminus \mathcal{V} \neq \emptyset$  do
    let  $p^* \leftarrow \arg \min_{p \in \mathcal{L} \setminus \mathcal{V}} \|x_p - x_q\|$ 
    update  $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$  and
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    if  $|\mathcal{L}| > L$  then
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    end if
  end while
  return [closest  $k$  points from  $\mathcal{L}$ ;  $\mathcal{V}$ ]

```

DiskANN [Subramanya+, NeurIPS 19]

When need to sort,
say “closest L points”

Algorithm 1 Beam search

Data: graph G , query q , initial vertex v_0 , output size k

Initialization:

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       $H := \text{Insert}(H, \hat{v}, d(\hat{v}, q))$ 
    end if
  end for
end while
return TopK( $V, q, k$ )

```

Learning to route [Baranchuk+, ICML 19]

Candidates are stored in a
heap; automatically sorted

Pseudo code

Algorithm 1. Search-on-Graph(G, p, q, l)

Require: graph G , start node p , query point q , candidate pool size l

Ensure: k nearest neighbors of q

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1:  $i = 0$ , candidate pool  $S = \emptyset$ 
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DiskANN [Subramanya+, NeurIPS 19]

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```

Learning to route [Baranchuk+, ICML 19]

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    end if
  end for
end while
return TopK( $V, q, k$ )
```

Learning to route [Baranchuk+, ICML 19]

Visited item are simply said to be “visited”; implying an additional hidden data structure (array)

Visited items are stored in a set

➤ Principles are the same. But small details are different.

➤ Hint: Explicitly state the data structure or not

Pseudo code

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Require: graph G , start node p , query point q , candidate pool size l

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Algorithm 1: GreedySearch(s, x_q, k, L)

Data: Graph G with start node s , query x_q , result size k , search list size $L \geq k$

Result: Result set $\mathcal{L}$ containing k -approx NNs, and a set $\mathcal{V}$ containing all the visited nodes

```
begin
  initialize sets  $\mathcal{L} \leftarrow \{s\}$  and  $\mathcal{V} \leftarrow \emptyset$ 
  while  $\mathcal{L} \setminus \mathcal{V} \neq \emptyset$  do
    let  $p^* \leftarrow \arg \min_{p \in \mathcal{L} \setminus \mathcal{V}} \|x_p - x_q\|$ 
    update  $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$  and
     $\mathcal{V} \leftarrow \mathcal{V} \cup \{p^*\}$ 
    if  $|\mathcal{L}| > L$  then
      update  $\mathcal{L}$  to retain closest  $L$ 
```

Termination condition??

Algorithm 1 Beam search

Data: graph G , query q , initial vertex v_0 , output size k

Initialization:

$V = \{v_0\}$ // a set of visited vertices

$H = \{v_0 : d(v_0, q)\}$ // a heap of candidates

```
while has runtime budget do
   $v_i = \text{SelectNearest}(H, q)$ 
  for  $\hat{v} \in \text{Expand}(v_i, G)$  do
    if  $\hat{v} \notin V$  then
       $V := \text{Add}(V, \hat{v})$ 
       $H := \text{Insert}(H, \hat{v}, d(\hat{v}, q))$ 
    end
  end
end
return TopK( $V, q, k$ )
```

Learning to route [Baranchuk+, ICML 19]

- All papers have totally different pseudo code 😞
- Principles are the same. But small details are different.
- Hint: Explicitly state the data structure or not

Pseudo code

Algorithm 1. Search-on-Graph(G, p, q, l)

Require: graph G , start node p , query point q , candidate pool size l

Ensure: k nearest neighbors of q

```
1:  $i = 0$ , candidate pool  $S = \emptyset$ 
2:  $S.add(p)$ 
3: while  $i < l$  do
4:    $i =$  the id of the first unchecked node  $p_i$  in  $S$ 
5:   mark  $p_i$  as checked
6:   for all neighbor  $n$  of  $p_i$  in  $G$  do
7:     if  $n$  has not been visited then
8:        $S.add(n)$ 
9:     end if
10:  end for
11:  sort  $S$  in ascending order of the distance to  $q$ 
12:  if  $S.size() > l$  then
13:     $S.resize(l)$  // remove nodes from back of  $S$  to keep its
14:    size no larger than  $l$ 
15:  end if
16: end while
17: return the first  $k$  nodes in  $S$ 
```

NSG [Cong+, VLDB 19]

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    if  $|\mathcal{L}| > L$  then
      update  $\mathcal{L}$  to retain closest  $L$ 
        points to  $x_q$ 
  end while
  return [closest  $k$  points from  $\mathcal{L}; \mathcal{V}$ ]
```

DiskANN [Subramanya+, NeurIPS 19]

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    end if
  end for
end while
return TopK( $V, q, k$ )
```

Learning to route [Baranchuk+, ICML 19]

➤ All papers have totally different pseudo code 😞

➤ My explanation was based on NSG, but with slight modifications for simplicity:

➤ ➤ Candidates are stored in an automatically-sorted array

➤ ➤ Termination condition is “all candidates are checked”

Pseudo code

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Ensure: k nearest neighbors of q

```
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update $\mathcal{L} \leftarrow \mathcal{L} \cup N_{out}(p^*)$ and

$\mathcal{V} \leftarrow \mathcal{V} \cup \{p^*\}$

update $\mathcal{L}$ to retain closest L

points to x_q

return [closest k points from $\mathcal{L}$; $\mathcal{V}$]

DiskANN [Subramanya+, NeurIPS 19]

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end

end

end

return TopK(V, q, k)

Learning to route [Baranchuk+, ICML 19]

Formal (?) definition would be helpful for everyone

- All papers have totally different pseudo code 😞
- Principles are the same. But small details are different.
- Hint: Explicitly state the data structure or not

Outline

1. History from an applications perspective
2. Importance of implementation
3. Basics of modern baseline

